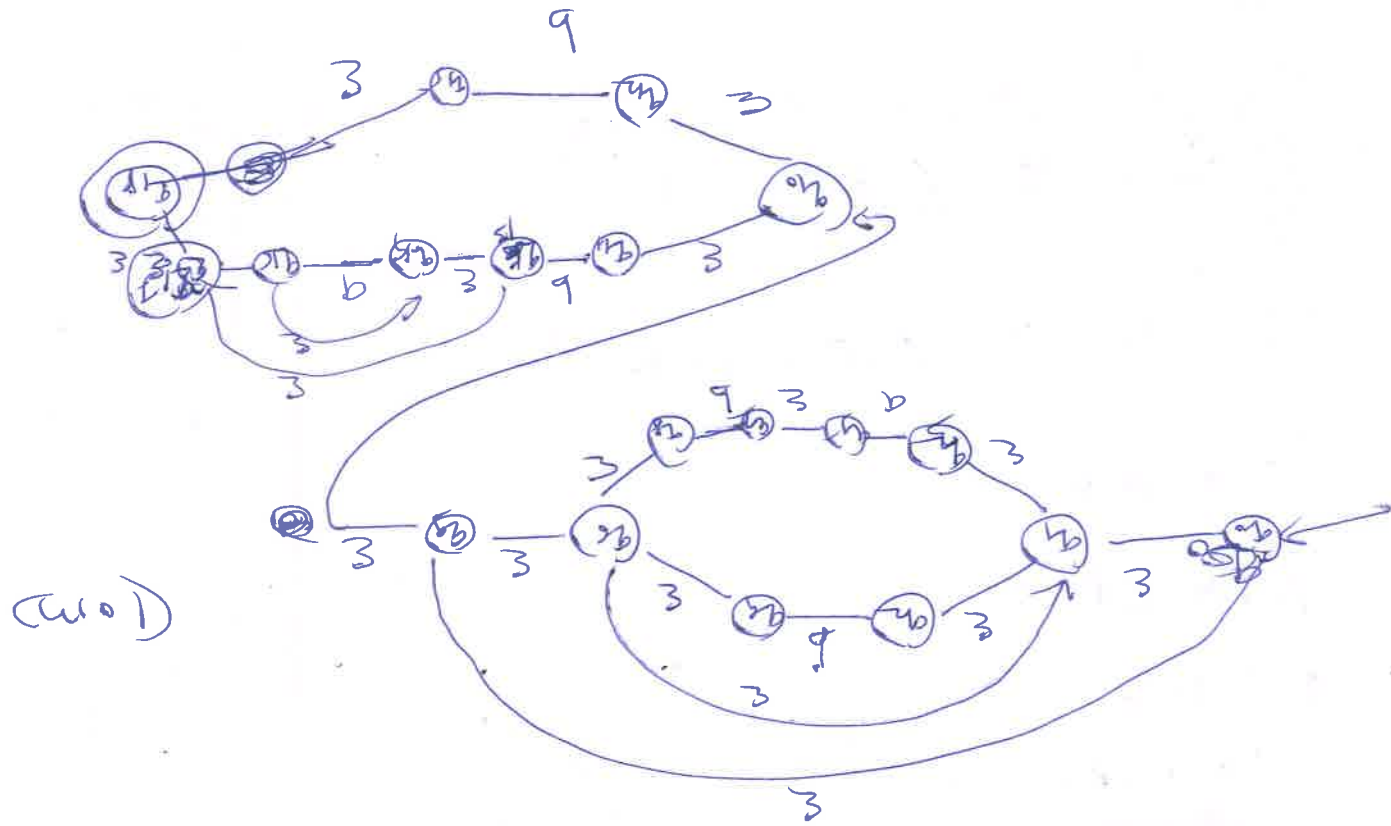


Questions should be —
WRITTEN IN LEGIBLE HANDWRITING IN BLACK INK.
SIGNS, SKETCHES OR FIGURES IF ANY BE DRAWN IN NEAT BLACK INK,
so as to avoid mistakes in the printed question papers.

Duration 03 Hours.

Total Marks assigned to the paper ...80...

| Q. No. | Answer | Marks |
|----------|---------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------|
| Q. 1 (a) | <p>N.B. :</p> | (5m) |
| Q. 2 (b) | <p>CFG for $(11)^* (010 + 01)^*$</p> <p> $S \rightarrow AB$ $A \rightarrow CA$ $B \rightarrow DB$ $C \rightarrow A_1 A_1$ $D \Rightarrow A_0 A_1 A_0 \mid A_0 A_1$ $A_0 \rightarrow 0$ $A_1 \rightarrow 1$ </p> | (5m) |
| Q. 3 | <p>Chomsky hierarchy</p> | (5m) |



(a)

Q.2 (a)

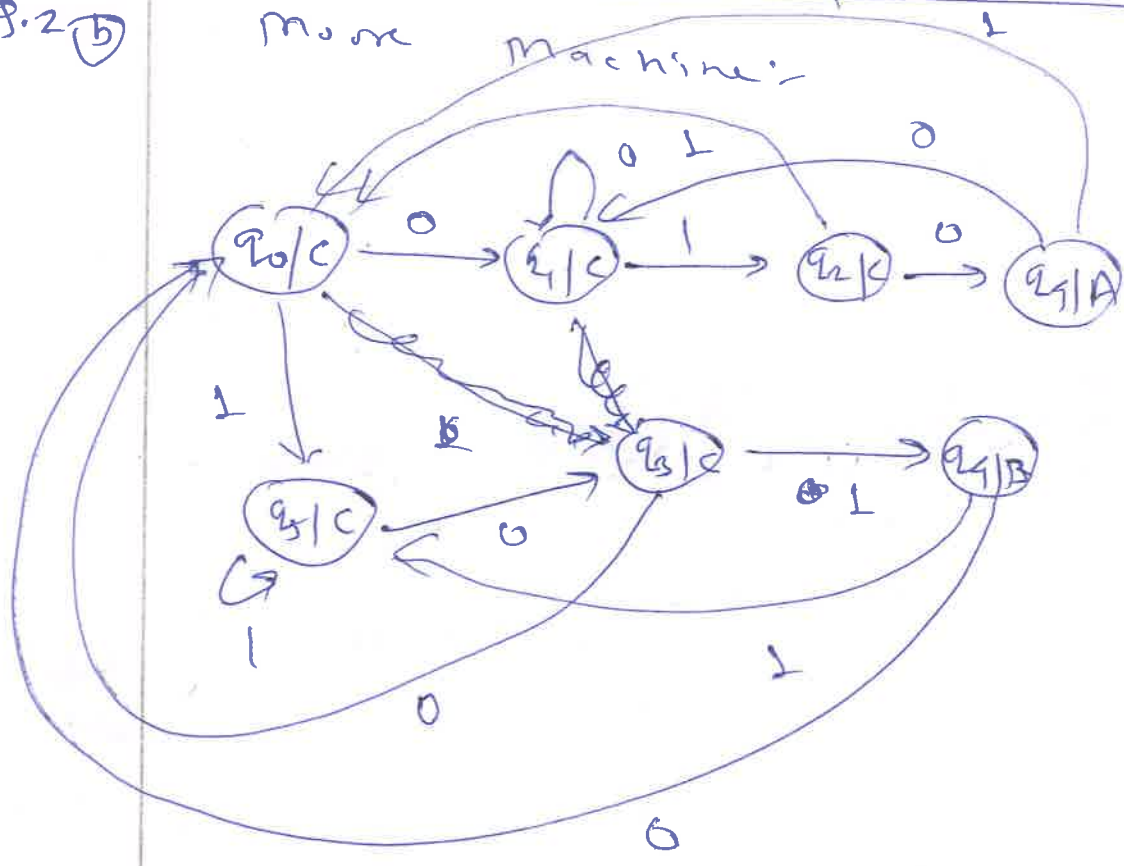
① NFA with epilogue -
 Σ = set of input symbols
 Q = set of states
 q_0 → initial state
 f → final state
 δ - transition function

(b)

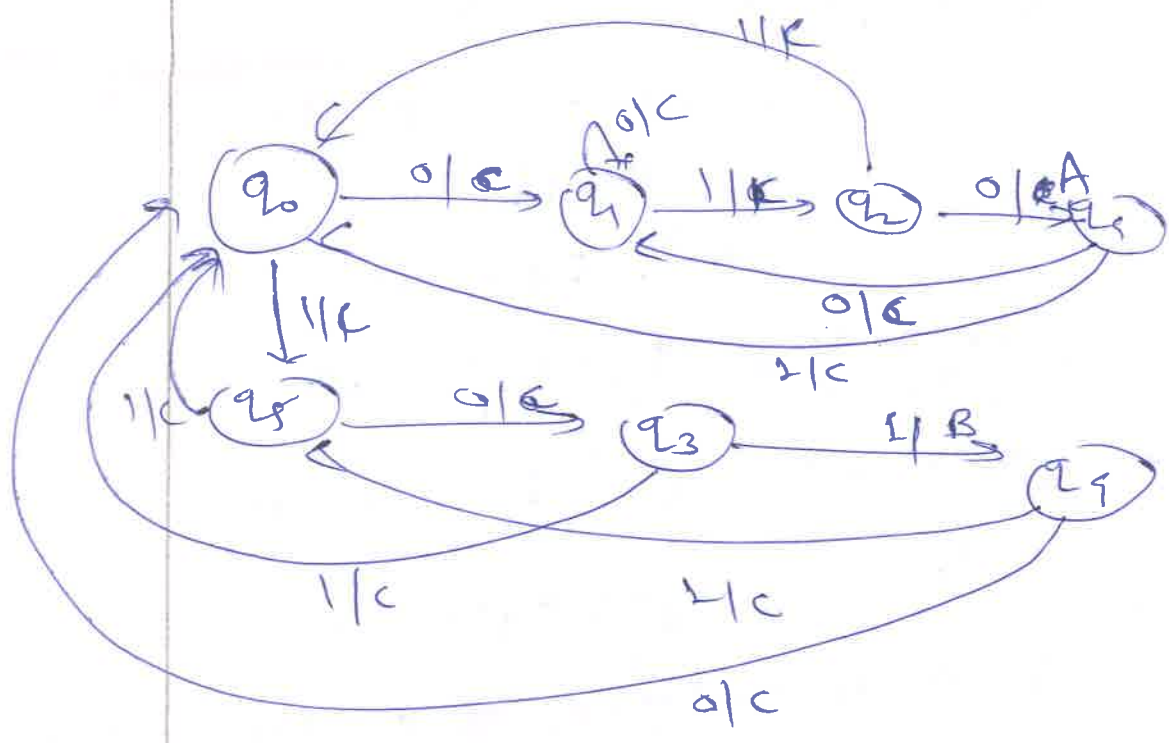
Q. No.
9.2 (b)

only consider last character of o/p as o/p

Moore machine:-



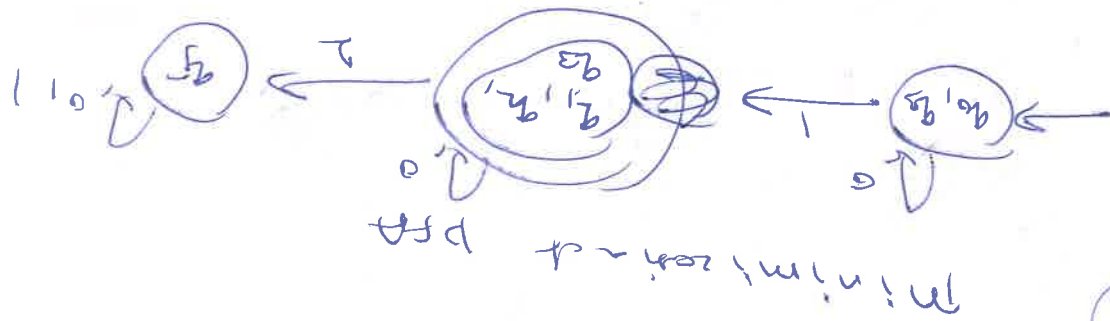
mealy machine :-



Marks

(10 m)

Q.3 P



(10m)

$S = \{a^n b^n \mid n \geq 1\}$ not regular

let $|xy| < p$
 $|y| > 0$
 If $x = a^i b^j$ and $z = a^k b^l$
 $\Rightarrow S = a^i b^j a^k b^l = a^i b^j a^k b^l$
 $x = a^i, y = a^j, z = a^k, l = b^l$

Q.3

Q.3 (a) let P be pumping length f

$S = a^f b^f a^f$

Q. No.

Marks

Q.4 (a)

rules for simplification of CFG.

(10m)

(i) useless production removal - The production that can never take part in derivation of any string, are called useless production.

(ii) λ production: The production of type $A \rightarrow \lambda$ called λ production (null production). These production cannot be removed from those grammar that don't generate λ .

(iii) unit production: The productions of type $A \rightarrow B$ are called unit production.

Q.4 (b)

CFG \Rightarrow

$$S \rightarrow ASB$$

$$S \rightarrow \epsilon$$

$$B \rightarrow S b S \mid A \mid L b$$

$$A \rightarrow a A S \mid a$$

Remove ϵ -production :-

(10m)

~~→~~

$S \rightarrow ASB / AB$
 $B = \cancel{SB} / b5 / b / A1 b2.$

Remove unit production, $B \rightarrow A$

$A \rightarrow aAs / aA / a$

$S \rightarrow ASB$

$S \rightarrow AB$

~~$S \rightarrow BS$~~

$B \rightarrow sb / b5 / b / bb$

$B \rightarrow aAs / aA / a.$

$A \rightarrow aAs / aA / a$

$CNF \rightarrow$

$S \rightarrow aSB$
 $S \rightarrow aB$
 $\therefore A \rightarrow a$

$B \rightarrow aBb / b5 / b / bb$

$A \rightarrow aAs / aA / a$

Q. No.

Q. 5 (a)

PDA:

$$L = \{ a^m b^m c^n \mid m, n \geq 1 \}$$

$$(q_0, a, \epsilon) \rightarrow (q_0, a)$$

$$(q_0, a, a) \rightarrow (q_0, a)$$

$$(q_0, b, a) \rightarrow (q_1, \epsilon)$$

$$(q_1, b, a) \rightarrow (q_1, \epsilon)$$

$$(q_1, c, \epsilon) \rightarrow (q_1, \epsilon)$$

$$(q_1, c, \epsilon) \rightarrow (q_1, \epsilon)$$

$$(q_1, \epsilon, \epsilon) \rightarrow (q_2, \epsilon) \Rightarrow \text{accept}$$

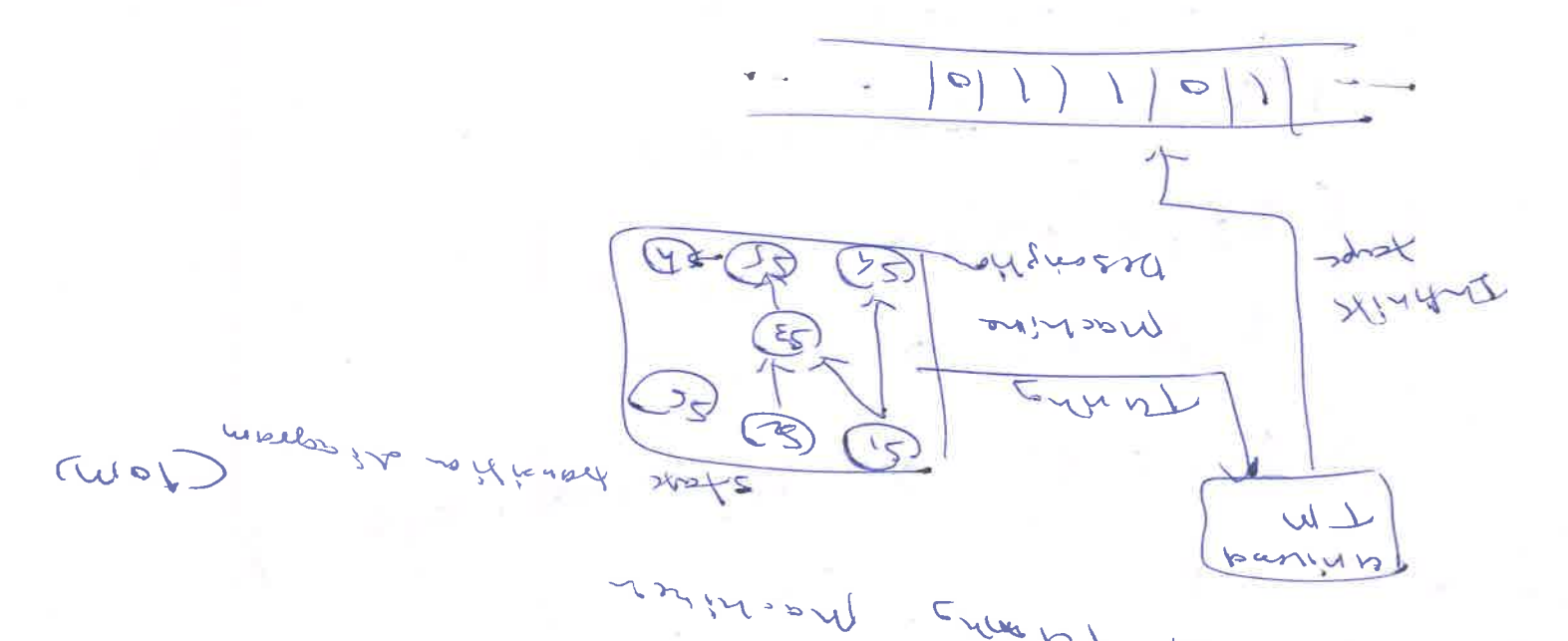
(10 marks)

Q. 5 (b)



(10 marks)

Q.6(a) pumping lemma for regular lang:-
 In the theory of formal lang, the pumping lemma for regular lang is a lemma that describe an essential property of all regular lang. Specifically, pumping lemma says that for any regular lang L , there exists a constant p such that any word w in L with length at least p can be split into three substring $w = xyz$ where $|x| \leq p$ and $|y| > 0$ and $xy^i z$ is in L for all $i \geq 0$.
 Note: x and y must not be empty. z can be empty.



Q.6(c) Unsolvability problem:-
 Def:- A computational problem that cannot be solved by a Turing Machine. The associated function is called unsolvable function.
 A problem in computer science is considered unsolvable when an expert in the field considers it unsolvable or when several experts in the field disagree about solution or problem.