

Chapter 1

Introduction to Fuzzy Logic

1.1 Fuzzy Logic: The word fuzzy means uncertainty. Any particular event which do not result any of the exact value (i.e. true or false) is fuzzy. Fuzzy Logic was introduced in 1965 by Lofti A. In other words, we can say that fuzzy logic is not logic that is fuzzy, but logic that is used to describe fuzziness. There can be numerous other examples like this with the help of which we can understand the concept of fuzzy logic.

The notion central to fuzzy systems is that truth values (in fuzzy logic) or membership values (in fuzzy sets) are indicated by a value on the range [0.0, 1.0], with 0.0 representing absolute Falseness and 1.0 representing absolute Truth. For example, let us take the statement:

"Rama is old."

If Rama's age was 75, we might assign the statement the truth value of 0.80. The statement could be translated into set terminology as follows:

"Rama is a member of the set of old people."

This statement would be rendered symbolically with fuzzy sets as:

$$m_{OLD}(Rama) = 0.80$$

where m is the membership function, operating in this case on the fuzzy set of old people, which returns a value between 0.0 and 1.0.

A **set** is an unordered collection of different elements. It can be written explicitly by listing its elements using the set bracket. If the order of the elements is changed or any element of a set is repeated, it does not make any changes in the set.

Example

- A set of all positive integers.
- A set of all the planets in the solar system.
- A set of all the states in India.
- A set of all the lowercase letters of the alphabet.

1.2 Differences between Fuzzy Logic and Neural Networks

Fuzzy logic allows making definite decisions based on imprecise or ambiguous data, whereas ANN tries to incorporate human thinking process to solve problems without mathematically modelling them. Even though both of these methods can be used to solve nonlinear problems, and problems that are not properly specified, they are not related. In contrast to Fuzzy logic, ANN tries to apply the thinking process in the human brain to solve problems. Further, ANN includes a learning process that involves learning algorithms and requires training data whereas

fuzzy logic includes development of membership functions and rules to relate them. Fuzzy logic basically deals with fixed and approximate (not exact) reasoning and the variables in fuzzy logic can take values from 0 to 1, this is contradicting to the traditional binary sets which takes value either 1 or 0 and since it can take any values in the range 0 to 1, it means that it is partially true and it is widely used for applications in control systems. Neural network on the other hand is based on biological neural network which is made up of artificial neurons interconnecting one another working in unison to produce outputs and it adapts to system with the data given to it by making adjustments to the synaptic connections that exist between the neurons.

Fuzzy logic makes decision based on the raw and ambiguous data given to it whereas Neural network tries to learn from the data, incorporating the same way involved in the biological neural network. Both of these system are used to solve non-linear and complex problems and are nowhere related to each other.

1.3 Applications of Fuzzy System:

Aerospace

In aerospace, fuzzy logic is used in the following areas –

- Altitude control of spacecraft
- Satellite altitude control
- Flow and mixture regulation in aircraft deicing vehicles

Automotive

In automotive, fuzzy logic is used in the following areas –

- Trainable fuzzy systems for idle speed control
- Shift scheduling method for automatic transmission
- Intelligent highway systems
- Traffic control
- Improving efficiency of automatic transmissions

Business

In business, fuzzy logic is used in the following areas –

- Decision-making support systems
- Personnel evaluation in a large company

Defense

In defense, fuzzy logic is used in the following areas –

- Underwater target recognition
- Automatic target recognition of thermal infrared images
- Naval decision support aids
- Control of a hypervelocity interceptor
- Fuzzy set modeling of NATO decision making

Electronics

In electronics, fuzzy logic is used in the following areas –

- Control of automatic exposure in video cameras
- Humidity in a clean room
- Air conditioning systems
- Washing machine timing
- Microwave ovens
- Vacuum cleaners

Finance

In the finance field, fuzzy logic is used in the following areas –

- Banknote transfer control
- Fund management
- Stock market predictions

Industrial Sector

In industrial, fuzzy logic is used in following areas –

- Cement kiln controls heat exchanger control
- Activated sludge wastewater treatment process control
- Water purification plant control
- Quantitative pattern analysis for industrial quality assurance
- Control of constraint satisfaction problems in structural design
- Control of water purification plants

Manufacturing

In the manufacturing industry, fuzzy logic is used in following areas –

- Optimization of cheese production
- Optimization of milk production

Marine

In the marine field, fuzzy logic is used in the following areas –

- Autopilot for ships
- Optimal route selection
- Control of autonomous underwater vehicles
- Ship steering

Medical

In the medical field, fuzzy logic is used in the following areas –

- Medical diagnostic support system
- Control of arterial pressure during anesthesia
- Multivariable control of anesthesia
- Modeling of neuropathological findings in Alzheimer's patients
- Radiology diagnoses
- Fuzzy inference diagnosis of diabetes and prostate cancer

Securities

In securities, fuzzy logic is used in following areas –

- Decision systems for securities trading
- Various security appliances

Transportation

In transportation, fuzzy logic is used in the following areas –

- Automatic underground train operation
- Train schedule control
- Railway acceleration
- Braking and stopping

Pattern Recognition and Classification

In Pattern Recognition and Classification, fuzzy logic is used in the following areas –

- Fuzzy logic based speech recognition
- Fuzzy logic based
- Handwriting recognition
- Fuzzy logic based facial characteristic analysis
- Command analysis
- Fuzzy image search

Psychology

In Psychology, fuzzy logic is used in following areas –

- Fuzzy logic based analysis of human behavior
- Criminal investigation and prevention based on fuzzy logic reasoning

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1.4 Historical Evolution Fuzzy Set:

Human Reasoning was dominated for centuries by the fundamental “Laws of Thought” (Korner, 1967), introduced by Aristotle (384-322 BC) and the philosophers that preceded him, which include:

- The principle of identity
- The law of the excluded middle
- The law of contradiction

In particular, the second of the above laws, stating that every proposition has to be either “True” or “False”, was the basis for the genesis of the Aristotle’s bi-valued Logic. The precision of the traditional mathematics owes undoubtedly a large part of its success to this Logic.

However, even when Parmenides proposed, around 400 BC, the first version of the law of the excluded middle, there were strong and immediate objections. For example, Heraclitus opposed that things could be simultaneously true and not true, whereas the Buddha Sidhartha Gautama, who lived in India a century earlier, had already indicated that almost every notion contains elements from its opposite one. The ancient Greek philosopher Plato (427-377 BC) laid the foundation of what it was later called FL by claiming that there exists a third area beyond “True” and “False”, where these two opposite notions can exist together. More modern philosophers like Hegel, Marx, Engels and others adopted and further cultivated the above Plato’s belief.

The Polish philosopher Jan Lukasiewicz (1878-1956) was the first to propose a systematic alternative of the bi-valued logic introducing in the early 1900’s a three valued logic by adding the term “Possible” between “True” and “False” (Lejewski, 1967). Eventually he developed an entire notation and axiomatic system from which he hoped to derive modern mathematics. Later he also proposed four and five valued

Logics and he finally arrived to the conclusion that axiomatically nothing could prevent the derivation of an infinite valued Logic.

But it was not until relatively recently that an infinite-valued Logic was introduced (Zadeh, 1973), called FL, because it is based on the notion of FS initiated in 1965 (Zadeh, 1965) by Lotfi (2018), Professor at the University of Berkeley, California. An important goal of FL is that through it algorithmic procedures can be devised which translate the “fuzzy” terminology into numerical values, perform reliable operations upon those values and then return natural language statements in a reliable manner.

Zadeh (1921–2017) (Wikipedia, retrieved from the Web on February, 2012) was born in Baku, Azerbaijan of USSR, to a Russian Jewish mother (FanyaKoriman), who was a pediatrician, and an Iranian Azeri father (Rahim Aleskerzade), who was a journalist on assignment from Iran.

At the age of 10, when Stalin introduced collectivization of farms in USSR, the Zadeh family moved to Iran. In 1942 Zadeh graduated from the University of Tehran with a degree in electrical engineering and moved to the USA in 1944. He received a MS from MIT in 1946 and a Ph.D. in electrical engineering from Columbia University in 1949. He taught for ten years in Columbia, promoted to a Full Professor in 1957, before moving to Berkeley in 1959. Among others he introduced jointly with J.R. Ragazzini in 1962 the pioneering *z-transform method* used today in the digital analysis (Brule, 2016) whereas his more recent works include computing with words and perceptions (Zadeh, 1984;2005a) and an outline towards a generalized theory of uncertainty (Zadeh, 2005b). It has been estimated that Zadeh, who died in Berkeley on 6 September 2017, aged 96, counted in 2011 more than 950 000 citations by other researchers!

As it was expected, the far-reaching theory of fuzzy systems aroused some objections to the scientific community. While there have been generic complaints about the fuzziness of assigning values to linguistic terms, the most cogent criticisms come from Haak (1979). She argued that there are only two areas – the nature of Truth and Falsity and the fuzzy systems’ utility – in which FL could be possibly needed, and then maintained that in both cases it can be shown that FL is unnecessary.

Fox (1981) responded to her objections indicating that FL is useful in three areas: To handle real-world relationships which are inherently fuzzy, to calculate the frequently existing in real world situations fuzzy data and to describe the operation of some inferential systems which are inherently fuzzy. His most powerful arguments were that traditional and FL need not be seen as competitive, but as complementary and that FL, despite the objections of classical logicians, has found its way into practical applications and has proved very successful there.

1.5 The concept of Fuzzy Set

Real life situations appear frequently where some definitions have not clear boundaries, like “the young people of a city”, “the good players of a team”, “the diligent students of

a class”, etc. The need to model mathematically such kind of situations was one of the main reasons that laid to the development of the FS theory.

Let U be the set of the discourse. Then, according to Zadeh (1965), a *fuzzy subset* A of U (or for brevity a FS in U) can be defined with the help of its *membership function* $m_A: U \rightarrow [0,1]$, which assigns to each element x of U a real value $m_A(x)$ in $[0, 1]$, called the *membership degree of x in A* . The closer is $m_A(x)$ to 1, the more x satisfies the characteristic property of A . Then one defines A as a set of ordered pairs of the form: $A = \{(x, m_A(x)) : x \in U\}$

Many authors, for reasons of simplicity, identify the FS A with its membership function m_A . A FS can be also denoted in the form of a symbolic sum, or a symbolic power series, or a symbolic integral, when U is a finite or numerable set or it has the power of the continuous respectively. For general facts on FS and the uncertainty connected to them we refer to the book of Klir and Folger (1988).

Example 1: The young human ages

Let U be the set of the non negative integers not exceeding 140 (considered as the upper bound of human life) representing the human ages. The set of all ages not exceeding a given integer in U , e.g. 20, is a crisp subset of U . On the contrary the set A of the young human ages, being not precisely defined, is a FS in U . The membership function of A can be defined by

$$m_A(x) = \begin{cases} \left[1 + (0.04x)^2\right]^{-1}, & x \leq 40 \\ 0 & , x > 40 \end{cases}$$

Therefore, the age of a recently born baby has membership degree $m_A(0) = 1$, the age of 25 years has membership degree:

$$m_A(25) = (1 + 1^2)^{-1} = 0.5, \text{ etc.}$$

Mathematical Representation of a Set

Sets can be represented in two ways –

Roster or Tabular Form

In this form, a set is represented by listing all the elements comprising it. The elements are enclosed within braces and separated by commas.

Following are the examples of set in Roster or Tabular Form –

- Set of vowels in English alphabet, $A = \{a, e, i, o, u\}$
- Set of odd numbers less than 10, $B = \{1, 3, 5, 7, 9\}$

Set Builder Notation

In this form, the set is defined by specifying a property that elements of the set have in common. The set is described as $A = \{x:p(x)\}$

Example 1 – The set $\{a,e,i,o,u\}$ is written as

$$A = \{x:x \text{ is a vowel in English alphabet}\}$$

Example 2 – The set $\{1,3,5,7,9\}$ is written as

$$B = \{x:1 \leq x < 10 \text{ and } (x\%2) \neq 0\}$$

If an element x is a member of any set S , it is denoted by $x \in S$ and if an element y is not a member of set S , it is denoted by $y \notin S$.

Example – If $S = \{1,1.2,1.7,2\}$, $1 \in S$ but $1.5 \notin S$

Cardinality of a Set

Cardinality of a set S , denoted by $|S|$, is the number of elements of the set. The number is also referred as the cardinal number. If a set has an infinite number of elements, its cardinality is ∞ .

Example – $|\{1,4,3,5\}| = 4, |\{1,2,3,4,5,\dots\}| = \infty$

If there are two sets X and Y , $|X| = |Y|$ denotes two sets X and Y having same cardinality. It occurs when the number of elements in X is exactly equal to the number of elements in Y . In this case, there exists a bijective function 'f' from X to Y .

$|X| \leq |Y|$ denotes that set X 's cardinality is less than or equal to set Y 's cardinality. It occurs when the number of elements in X is less than or equal to that of Y . Here, there exists an injective function 'f' from X to Y .

$|X| < |Y|$ denotes that set X 's cardinality is less than set Y 's cardinality. It occurs when the number of elements in X is less than that of Y . Here, the function 'f' from X to Y is injective function but not bijective.

If $|X| \leq |Y|$ and $|Y| \leq |X|$ then $|X| = |Y|$. The sets X and Y are commonly referred as **equivalent sets**.

Types of Sets

Sets can be classified into many types; some of which are finite, infinite, subset, universal, proper, singleton set, etc.

Finite Set

A set which contains a definite number of elements is called a finite set.

Example – $S = \{x | x \in \mathbb{N} \text{ and } 70 > x > 50\}$

Infinite Set

A set which contains infinite number of elements is called an infinite set.

Example – $S = \{x | x \in \mathbb{N} \text{ and } x > 10\}$

Subset

A set X is a subset of set Y (Written as $X \subseteq Y$) if every element of X is an element of set Y .

Example 1 – Let, $X = \{1,2,3,4,5,6\}$ and $Y = \{1,2\}$. Here set Y is a subset of set X as all the elements of set Y is in set X . Hence, we can write $Y \subseteq X$.

Example 2 – Let, $X = \{1,2,3\}$ and $Y = \{1,2,3\}$. Here set Y is a subset (not a proper subset) of set X as all the elements of set Y is in set X . Hence, we can write $Y \subseteq X$.

Proper Subset

The term “proper subset” can be defined as “subset of but not equal to”. A Set X is a proper subset of set Y (Written as $X \subset Y$) if every element of X is an element of set Y and $|X| < |Y|$.

Example – Let, $X = \{1,2,3,4,5,6\}$ and $Y = \{1,2\}$. Here set $Y \subset X$, since all elements in Y are contained in X too and X has at least one element which is more than set Y .

Universal Set

It is a collection of all elements in a particular context or application. All the sets in that context or application are essentially subsets of this universal set. Universal sets are represented as U .

Example – We may define U as the set of all animals on earth. In this case, a set of all mammals is a subset of U , a set of all fishes is a subset of U , a set of all insects is a subset of U , and so on.

Empty Set or Null Set

An empty set contains no elements. It is denoted by Φ . As the number of elements in an empty set is finite, empty set is a finite set. The cardinality of empty set or null set is zero.

Example – $S = \{x|x \in \mathbb{N} \text{ and } 7 < x < 8\} = \Phi$

Singleton Set or Unit Set

A Singleton set or Unit set contains only one element. A singleton set is denoted by $\{s\}$.

Example – $S = \{x|x \in \mathbb{N}, 7 < x < 9\} = \{8\}$

Equal Set

If two sets contain the same elements, they are said to be equal.

Example – If $A = \{1,2,6\}$ and $B = \{6,1,2\}$, they are equal as every element of set A is an element of set B and every element of set B is an element of set A.

Equivalent Set

If the cardinalities of two sets are same, they are called equivalent sets.

Example – If $A = \{1,2,6\}$ and $B = \{16,17,22\}$, they are equivalent as cardinality of A is equal to the cardinality of B. i.e. $|A| = |B| = 3$

Overlapping Set

Two sets that have at least one common element are called overlapping sets. In case of overlapping sets –

$$n(A \cup B) = n(A) + n(B) - n(A \cap B)$$

$$n(A \cup B) = n(A - B) + n(B - A) + n(A \cap B)$$

$$n(A) = n(A - B) + n(A \cap B)$$

$$n(B) = n(B - A) + n(A \cap B)$$

Example – Let, $A = \{1,2,6\}$ and $B = \{6,12,42\}$. There is a common element '6', hence these sets are overlapping sets.

Disjoint Set

Two sets A and B are called disjoint sets if they do not have even one element in common. Therefore, disjoint sets have the following properties –

$$n(A \cap B) = \phi$$

$$n(A \cup B) = n(A) + n(B)$$

Example – Let, $A = \{1, 2, 6\}$ and $B = \{7, 9, 14\}$, there is not a single common element, hence these sets are overlapping sets.

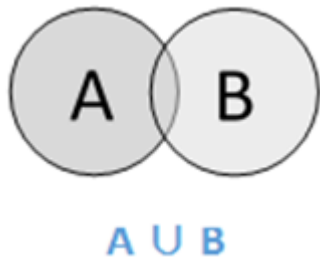
Operations on Classical Sets

Set Operations include Set Union, Set Intersection, Set Difference, Complement of Set, and Cartesian Product.

Union

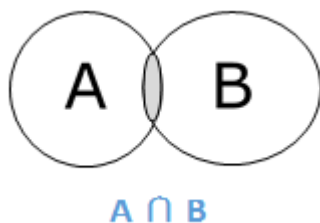
The union of sets A and B (denoted by $A \cup B$) is the set of elements which are in A, in B, or in both A and B. Hence, $A \cup B = \{x | x \in A \text{ OR } x \in B\}$.

Example – If $A = \{10, 11, 12, 13\}$ and $B = \{13, 14, 15\}$, then $A \cup B = \{10, 11, 12, 13, 14, 15\}$ – The common element occurs only once.



Intersection

The intersection of sets A and B (denoted by $A \cap B$) is the set of elements which are in both A and B. Hence, $A \cap B = \{x | x \in A \text{ AND } x \in B\}$.



Difference/ Relative Complement

The set difference of sets A and B (denoted by $A - B$) is the set of elements which are only in A but not in B. Hence, $A - B = \{x | x \in A \text{ AND } x \notin B\}$.

Example – If $A = \{10, 11, 12, 13\}$ and $B = \{13, 14, 15\}$, then $(A - B) = \{10, 11, 12\}$ and $(B - A) = \{14, 15\}$. Here, we can see $(A - B) \neq (B - A)$

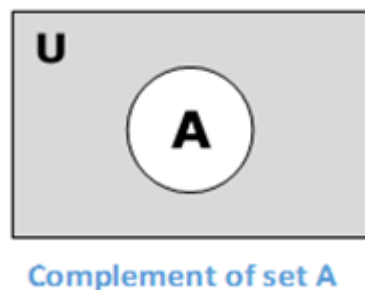


Complement of a Set

The complement of a set A (denoted by A') is the set of elements which are not in set A. Hence, $A' = \{x | x \notin A\}$.

More specifically, $A' = (U - A)$ where U is a universal set which contains all objects.

Example – If $A = \{x | x \text{ belongs to set of add integers}\}$ then $A' = \{y | y \text{ does not belong to set of odd integers}\}$



Cartesian Product / Cross Product

The Cartesian product of n number of sets A_1, A_2, \dots, A_n denoted as $A_1 \times A_2 \times \dots \times A_n$ can be defined as all possible ordered pairs (x_1, x_2, \dots, x_n) where $x_1 \in A_1, x_2 \in A_2, \dots, x_n \in A_n$

Example – If we take two sets $A = \{a, b\}$ and $B = \{1, 2\}$,

The Cartesian product of A and B is written as – $A \times B = \{(a, 1), (a, 2), (b, 1), (b, 2)\}$

And, the Cartesian product of B and A is written as – $B \times A = \{(1,a),(1,b),(2,a),(2,b)\}$

Properties of Classical Sets

Properties on sets play an important role for obtaining the solution. Following are the different properties of classical sets –

Commutative Property

Having two sets **A** and **B**, this property states –

$$A \cup B = B \cup A$$

$$A \cap B = B \cap A$$

Associative Property

Having three sets **A**, **B** and **C**, this property states –

$$A \cup (B \cap C) = (A \cup B) \cap C$$

$$A \cap (B \cup C) = (A \cap B) \cup C$$

Distributive Property

Having three sets **A**, **B** and **C**, this property states –

$$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

Idempotency Property

For any set **A**, this property states –

$$A \cup A = A$$

$$A \cap A = A$$

Identity Property

For set **A** and universal set **X**, this property states –

$$A \cup \varnothing = A$$

$$A \cap X = A$$

$$A \cap \varnothing = \varnothing$$

$$A \cup X = X$$

Transitive Property

Having three sets **A**, **B** and **C**, the property states –

$$\text{If } A \subseteq B \subseteq C$$

, then $A \subseteq C$

De Morgan's Law

It is a very important law and supports in proving tautologies and contradiction. This law states –

$$\overline{A \cap B} = \overline{A} \cup \overline{B}$$

$$\overline{A \cup B} = \overline{A} \cap \overline{B}$$

Exercise:

1. What is Fuzzy System? Give Example for the same.
2. Explain the concept of fuzziness with suitable example.
3. State the applications of fuzzy sets.
4. What is fuzzy set/
5. Give comparison between fuzzy system and neural networks.
6. Give the historical evolution of fuzzy system.

Chapter 2

Fuzzy Sets

Introduction:

Fuzzy logic starts with and builds on a set of user-supplied human language rules. The fuzzy systems convert these rules to their mathematical equivalents. This simplifies the job of the system designer and the computer, and results in much more accurate representations of the way systems behave in the real world.

Additional benefits of fuzzy logic include its simplicity and its flexibility. Fuzzy logic can handle problems with imprecise and incomplete data, and it can model nonlinear functions of arbitrary complexity. "If you don't have a good plant model, or if the system is changing, then fuzzy will produce a better solution than conventional control techniques," says Bob Varley, a Senior Systems Engineer at Harris Corp., an aerospace company in Palm Bay, Florida.

You can create a fuzzy system to match any set of input-output data. The Fuzzy Logic Toolbox makes this particularly easy by supplying adaptive techniques such as adaptive neuro-fuzzy inference systems (ANFIS) and fuzzy subtractive clustering.

Fuzzy logic models, called fuzzy inference systems, consist of a number of conditional "if-then" rules. For the designer who understands the system, these rules are easy to write, and as many rules as necessary can be supplied to describe the system adequately (although typically only a moderate number of rules are needed).

In fuzzy logic, unlike standard conditional logic, the truth of any statement is a matter of degree. (How cold is it? How high should we set the heat?) We are familiar with inference rules of the form $p \rightarrow q$ (p implies q). With fuzzy logic, it's possible to say $(.5 * p) \rightarrow (.5 * q)$. For example, for the rule if (weather is cold) then (heat is on), both variables, cold and on, map to ranges of values. Fuzzy inference systems rely on membership functions to explain to the computer how to calculate the correct value between 0 and 1. The degree to which any fuzzy statement is true is denoted by a value between 0 and 1.

Not only do the rule-based approach and flexible membership function scheme make fuzzy systems straightforward to create, but they also simplify the design of systems and ensure that you can easily update and maintain the system over time.

Fuzzy Set Theory was formalised by Professor LoftiZadeh at the University of California in 1965. What Zadeh proposed is very much a paradigm shift that first gained

acceptance in the Far East and its successful application has ensured its adoption around the world.

A paradigm is a set of rules and regulations which defines boundaries and tells us what to do to be successful in solving problems within these boundaries. For example the use of transistors instead of vacuum tubes is a paradigm shift - likewise the development of Fuzzy Set Theory from conventional bivalent set theory is a paradigm shift.

Bivalent Set Theory can be somewhat limiting if we wish to describe a 'humanistic' problem mathematically. For example, Fig 1 below illustrates bivalent sets to characterise the temperature of a room.

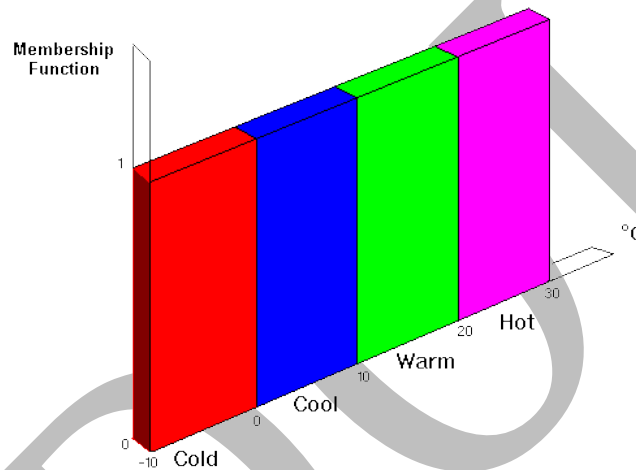


Fig. 1 : Bivalent Sets to Characterize the Temp. of a room.

The most obvious limiting feature of bivalent sets that can be seen clearly from the diagram is that they are mutually exclusive - it is not possible to have membership of more than one set (opinion would widely vary as to whether 50 degrees Fahrenheit is 'cold' or 'cool' hence the expert knowledge we need to define our system is mathematically at odds with the humanistic world). Clearly, it is not accurate to define a transition from a quantity such as 'warm' to 'hot' by the application of one degree Fahrenheit of heat. In the real world a smooth (unnoticeable) drift from warm to hot would occur.

This natural phenomenon can be described more accurately by Fuzzy Set Theory. Fig.2 below shows how fuzzy sets quantifying the same information can describe this natural drift.

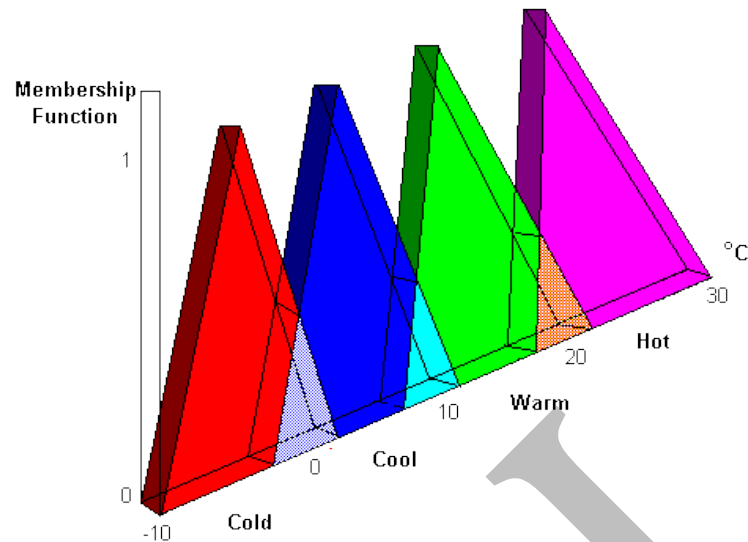


Fig. 2 - Fuzzy Sets to characterize the Temp. of a room.

Fuzzy Logic is a problem-solving control system methodology that lends itself to implementation in systems ranging from simple, small, embedded micro-controllers to large, networked, multi-channel PC or workstation-based data acquisition and control systems. It can be implemented in hardware, software, or a combination of both. FL provides a simple way to arrive at a definite conclusion based upon vague, ambiguous, imprecise, noisy, or missing input information. FL's approach to control problems.

FL requires some numerical parameters in order to operate such as what is considered significant error and significant rate-of-change-of-error, but exact values of these numbers are usually not critical unless very responsive performance is required in which case empirical tuning would determine them. For example, a simple temperature control system could use a single temperature feedback sensor whose data is subtracted from the command signal to compute "error" and then time-differentiated to yield the error slope or rate-of-change-of-error, hereafter called "error-dot".

2.1 History of Fuzzy Logic

Although, the concept of fuzzy logic had been studied since the 1920's. The term fuzzy logic was first used with 1965 by LotfiZadeh a professor of UC Berkeley in California. He observed that conventional computer logic was not capable of manipulating data representing subjective or unclear human ideas.

Fuzzy logic has been applied to various fields, from control theory to AI. It was designed to allow the computer to determine the distinctions among data which is neither true nor false. Something similar to the process of human reasoning. Like Little dark, Some brightness, etc.

2.2 Characteristics of Fuzzy Logic

Here, are some important characteristics of fuzzy logic:

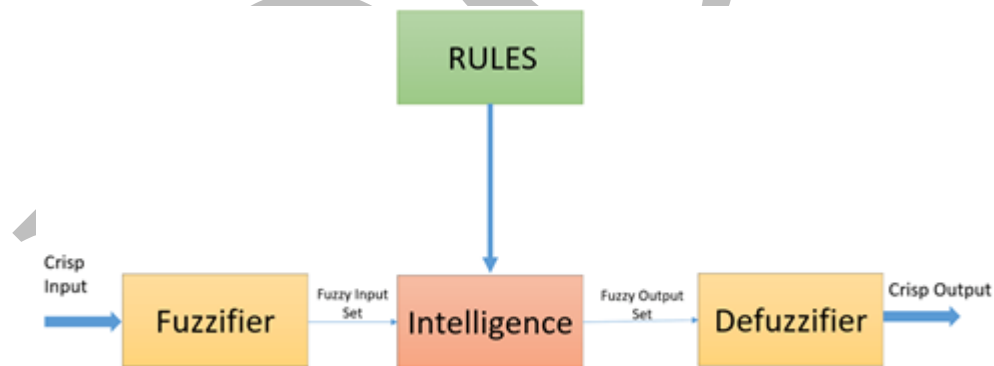
- Flexible and easy to implement machine learning technique
- Helps you to mimic the logic of human thought
- Logic may have two values which represent two possible solutions
- Highly suitable method for uncertain or approximate reasoning
- Fuzzy logic views inference as a process of propagating elastic constraints
- Fuzzy logic allows you to build nonlinear functions of arbitrary complexity.
- Fuzzy logic should be built with the complete guidance of experts

However, fuzzy logic is never a cure for all. Therefore, it is equally important to understand that where we should not use fuzzy logic.

Here, are certain situations when you better not use Fuzzy Logic:

- If you don't find it convenient to map an input space to an output space
- Fuzzy logic should not be used when you can use common sense
- Many controllers can do the fine job without the use of fuzzy logic

2.3 Fuzzy Logic Architecture



Fuzzy Logic architecture has four main parts as shown in the diagram:

Rule Base:

It contains all the rules and the if-then conditions offered by the experts to control the decision-making system. The recent update in fuzzy theory provides various methods for the design and tuning of fuzzy controllers. This updates significantly reduce the number of the fuzzy set of rules.

Fuzzification:

Fuzzification step helps to convert inputs. It allows you to convert, crisp numbers into fuzzy sets. Crisp inputs measured by sensors and passed into the control system for further processing. Like Room temperature, pressure, etc.

Inference Engine:

It helps you to determine the degree of match between fuzzy input and the rules. Based on the % match, it determines which rules need implement according to the given input field. After this, the applied rules are combined to develop the control actions.

Defuzzification:

At last the Defuzzification process is performed to convert the fuzzy sets into a crisp value. There are many types of techniques available, so you need to select it which is best suited when it is used with an expert system.

Fuzzy Logic vs. Probability

Fuzzy Logic	Probability
Fuzzy: Tom's degree of membership within the set of old people is 0.90.	Probability: There is a 90% chance that Tom is old.
Fuzzy logic takes truth degrees as a mathematical basis on the model of the vagueness phenomenon.	Probability is a mathematical model of ignorance.

Crisp vs. Fuzzy

Crisp	Fuzzy
It has strict boundary T or F	Fuzzy boundary with a degree of membership
Some crisp time set can be fuzzy	It can't be crisp
True/False {0,1}	Membership values on [0,1]
In Crisp logic law of Excluded Middle and Non- Contradiction may or may not hold	In the fuzzy logic law of Excluded Middle and Non- Contradiction hold

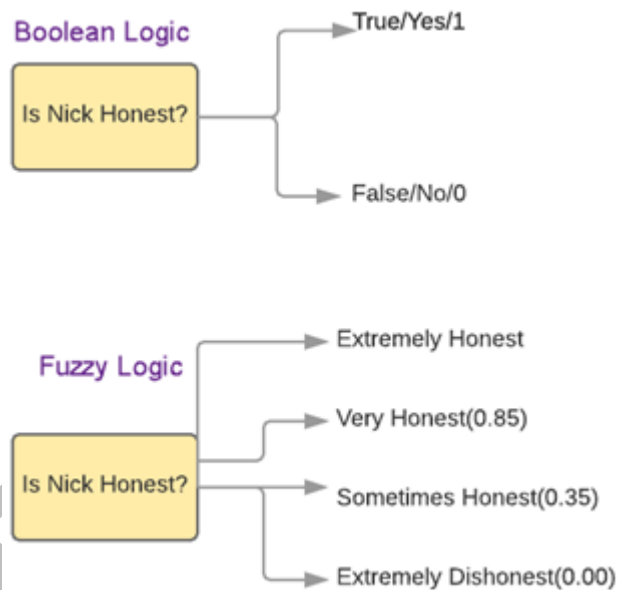
Classical Set vs. Fuzzy set Theory

Classical Set	Fuzzy Set Theory
Classes of objects with sharp boundaries.	Classes of objects do not have sharp boundaries.
A classical set is defined by crisp boundaries, i.e., there is clarity about the location of the set	A fuzzy set always has ambiguous boundaries, i.e., there may be uncertainty about the location

boundaries.	of the set boundaries.
Widely used in digital system design	Used only in fuzzy controllers.

Fuzzy Logic Examples

See the below-given diagram. It shows that in fuzzy systems, the values are denoted by a 0 to 1 number. In this example, 1.0 means absolute truth and 0.0 means absolute falseness.



2.4 Application Areas of Fuzzy Logic

The Below given table shows how famous companies using fuzzy logic in their products.

Product	Company	Fuzzy Logic
Anti-lock brakes	Nissan	Use fuzzy logic to controls brakes in hazardous cases depend on car speed, acceleration, wheel speed, and acceleration
Auto transmission	NOK/Nissan	Fuzzy logic is used to control the fuel injection and ignition based on throttle setting, cooling water temperature, RPM, etc.
Auto engine	Honda, Nissan	Use to select geat based on engine load, driving style, and road conditions.
Copy machine	Canon	Using for adjusting drum voltage based on picture

		density, humidity, and temperature.
Cruise control	Nissan, Isuzu, Mitsubishi	Use it to adjust throttle setting to set car speed and acceleration
Dishwasher	Matsushita	Use for adjusting the cleaning cycle, rinse and wash strategies based depend upon the number of dishes and the amount of food served on the dishes.
Elevator control	Fujitec, Mitsubishi Electric, Toshiba	Use it to reduce waiting for time-based on passenger traffic
Golf diagnostic system	Maruman Golf	Selects golf club based on golfer's swing and physique.
Fitness management	Omron	Fuzzy rules implied by them to check the fitness of their employees.
Kiln control	Nippon Steel	Mixes cement
Microwave oven	Mitsubishi Chemical	Sets power and cooking strategy
Palmtop computer	Hitachi, Sharp, Sanyo, Toshiba	Recognizes handwritten Kanji characters
Plasma etching	Mitsubishi Electric	Sets etch time and strategy

2.4.1 Advantages of Fuzzy Logic System

- The structure of Fuzzy Logic Systems is easy and understandable
- Fuzzy logic is widely used for commercial and practical purposes
- It helps you to control machines and consumer products
- It may not offer accurate reasoning, but the only acceptable reasoning
- It helps you to deal with the uncertainty in engineering
- Mostly robust as no precise inputs required
- It can be programmed to in the situation when feedback sensor stops working
- It can easily be modified to improve or alter system performance
- inexpensive sensors can be used which helps you to keep the overall system cost and complexity low
- It provides a most effective solution to complex issues

2.4.2 Disadvantages of Fuzzy Logic Systems

- Fuzzy logic is not always accurate, so The results are perceived based on assumption, so it may not be widely accepted.
- Fuzzy systems don't have the capability of machine learning as-well-as neural network type pattern recognition

- Validation and Verification of a fuzzy knowledge-based system needs extensive testing with hardware
- Setting exact, fuzzy rules and, membership functions is a difficult task
- Some fuzzy time logic is confused with probability theory and the terms

2.5 Features of Membership Functions

We will now discuss the different features of Membership Functions.

Core

For any fuzzy set \tilde{A} , the core of a membership function is that region of universe that is characterized by full membership in the set. Hence, core consists of all those elements y of the universe of information such that,

$$\mu_{\tilde{A}}(y)=1$$

Support

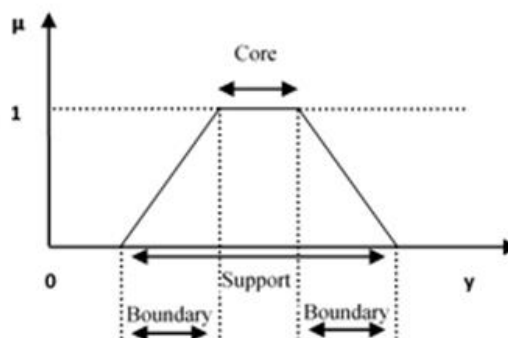
For any fuzzy set \tilde{A} , the support of a membership function is the region of universe that is characterized by a nonzero membership in the set. Hence core consists of all those elements y of the universe of information such that,

$$\mu_{\tilde{A}}(y)>0$$

Boundary

For any fuzzy set \tilde{A} , the boundary of a membership function is the region of universe that is characterized by a nonzero but incomplete membership in the set. Hence, core consists of all those elements y of the universe of information such that,

$$1>\mu_{\tilde{A}}(y)>0$$



Features of Membership Function

2.5.1 Fuzzification

It may be defined as the process of transforming a crisp set to a fuzzy set or a fuzzy set to fuzzier set. Basically, this operation translates accurate crisp input values into linguistic variables.

Following are the two important methods of fuzzification –

Support Fuzzification(s-fuzzification) Method

In this method, the fuzzified set can be expressed with the help of the following relation

$$\tilde{A} = \mu_1 Q(x_1) + \mu_2 Q(x_2) + \dots + \mu_n Q(x_n)$$

Here the fuzzy set $Q(x_i)$ is called as kernel of fuzzification. This method is implemented by keeping μ_i constant and x_i being transformed to a fuzzy set $Q(x_i)$.

Grade Fuzzification (g-fuzzification) Method

It is quite similar to the above method but the main difference is that it kept x_i constant and μ_i is expressed as a fuzzy set.

2.5.2 Defuzzification

It may be defined as the process of reducing a fuzzy set into a crisp set or to convert a fuzzy member into a crisp member.

We have already studied that the fuzzification process involves conversion from crisp quantities to fuzzy quantities. In a number of engineering applications, it is necessary to defuzzify the result or rather “fuzzy result” so that it must be converted to crisp result. Mathematically, the process of Defuzzification is also called “rounding it off”.

The different methods of Defuzzification are described below –

Max-Membership Method

This method is limited to peak output functions and also known as height method. Mathematically it can be represented as follows –

$$\mu_{\tilde{A}}(x^*) > \mu_{\tilde{A}}(x) \text{ for all } x \in X$$

Here, x^* is the defuzzified output.

Centroid Method

This method is also known as the center of area or the center of gravity method. Mathematically, the defuzzified output x^* will be represented as –

$$x^* = \frac{\int \mu_{\tilde{A}}(x) \cdot x dx}{\int \mu_{\tilde{A}}(x) \cdot dx}$$

Weighted Average Method

In this method, each membership function is weighted by its maximum membership value. Mathematically, the defuzzified output x^*

will be represented as –

$$x^* = \frac{\sum \mu_{\tilde{A}}(\bar{x}_i) \cdot \bar{x}_i}{\sum \mu_{\tilde{A}}(\bar{x}_i)}$$

Mean-Max Membership

This method is also known as the middle of the maxima. Mathematically, the defuzzified output x^*

will be represented as –

$$x^* = \sum_{i=1}^n \bar{x}_i \cdot n$$

2.6 Operations on Fuzzy Sets

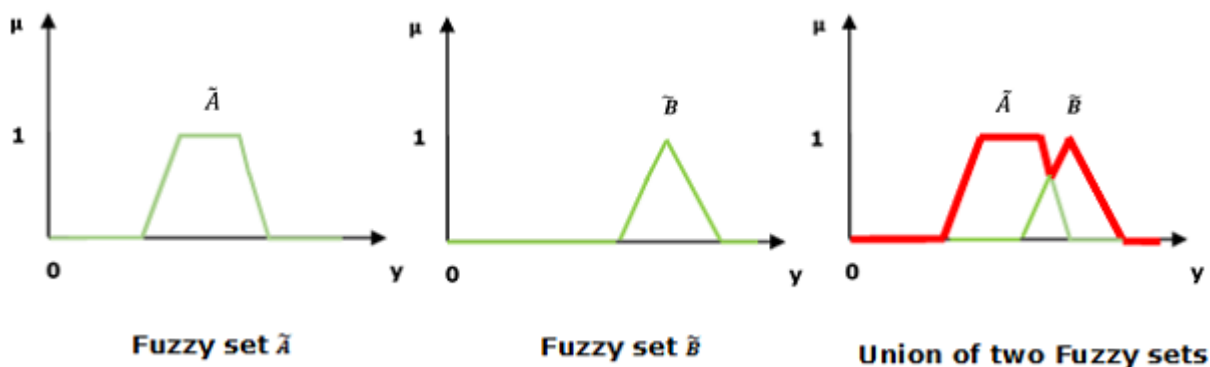
Having two fuzzy sets \tilde{A} and \tilde{B} , the universe of information U and an element y of the universe, the following relations express the union, intersection and complement operation on fuzzy sets.

Union/Fuzzy ‘OR’

Let us consider the following representation to understand how the **Union/Fuzzy ‘OR’** relation works –

$$\mu_{\tilde{A} \cup \tilde{B}}(y) = \mu_{\tilde{A}} \vee \mu_{\tilde{B}} \quad \forall y \in U$$

Here \vee represents the ‘max’ operation.

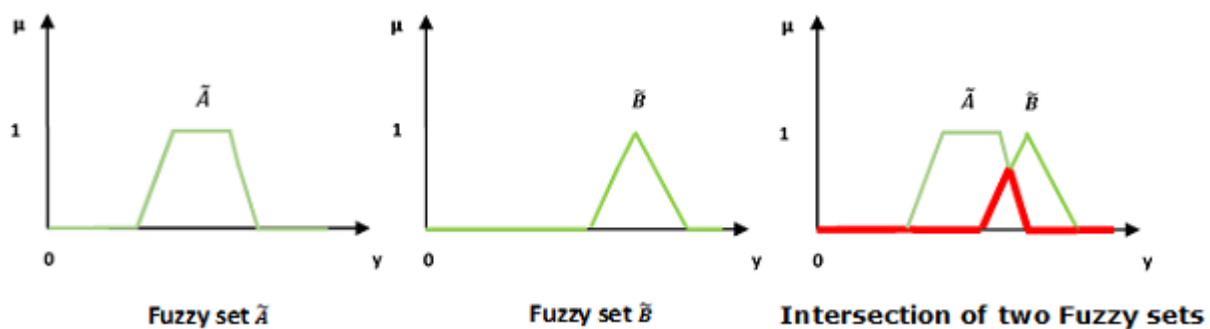


Intersection/Fuzzy ‘AND’

Let us consider the following representation to understand how the **Intersection/Fuzzy ‘AND’** relation works –

$$\mu_{\tilde{A} \cap \tilde{B}}(y) = \mu_{\tilde{A}} \wedge \mu_{\tilde{B}} \quad \forall y \in U$$

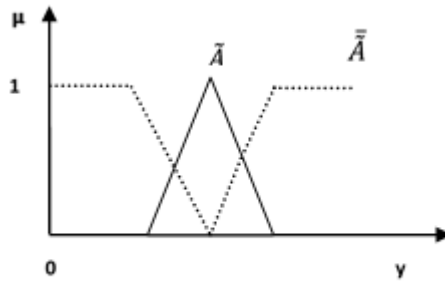
Here \wedge represents the ‘min’ operation.



Complement/Fuzzy ‘NOT’

Let us consider the following representation to understand how the **Complement/Fuzzy ‘NOT’** relation works –

$$\mu_{\tilde{A}^c} = 1 - \mu_{\tilde{A}}(y), y \in U$$



Complement of a fuzzy set

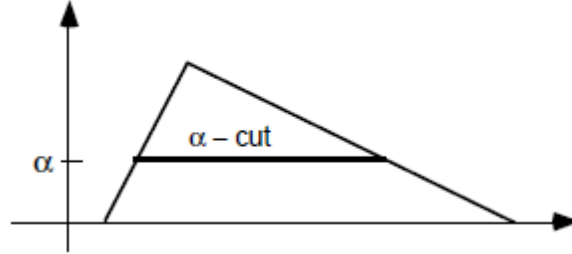
Definition. (support) Let A be a fuzzy subset of X ; the support of A , denoted $\text{supp}(A)$, is the crisp subset of X whose elements all have nonzero membership grades in A .

$$\text{supp}(A) = \{x \in X \mid A(x) > 0\}.$$

Definition. (normal fuzzy set) A fuzzy subset A of a classical set X is called normal if there exists an $x \in X$ such that $A(x) = 1$. Otherwise A is subnormal.

Definition. (α -cut) An α -level set of a fuzzy set A of X is a non-fuzzy set denoted by $[A]^\alpha$ and is defined by

$$[A]^\alpha = \begin{cases} \{t \in X \mid A(t) \geq \alpha\} & \text{if } \alpha > 0 \\ \text{cl}(\text{supp} A) & \text{if } \alpha = 0 \end{cases}$$



where $\text{cl}(\text{supp} A)$ denotes the closure of the support of A .

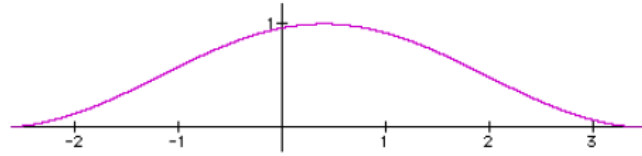
Definition. (convex fuzzy set) A fuzzy set A of X is called convex if $[A]^\alpha$ is a convex subset of $X \ \forall \alpha \in [0, 1]$. An α -cut of a triangular fuzzy number.

In many situations people are only able to characterize numeric information imprecisely. For example, people use terms such as, about 5000, near zero, or essentially bigger than 5000. These are examples of what are called *fuzzy numbers*. Using the theory of fuzzy subsets we can represent these fuzzy numbers as fuzzy subsets of the set of real numbers.

Definition. (fuzzy number) A fuzzy number A is a fuzzy set of the real line with a normal, (fuzzy) convex and continuous membership function of bounded support. The family of fuzzy numbers will be denoted by F .

Definition. (quasi fuzzy number) A quasi fuzzy number A is a fuzzy set of the real line with a normal, fuzzy convex and continuous membership functions satisfying the limit conditions

$$\lim_{t \rightarrow \infty} A(t) = 0, \quad \lim_{t \rightarrow -\infty} A(t) = 0.$$

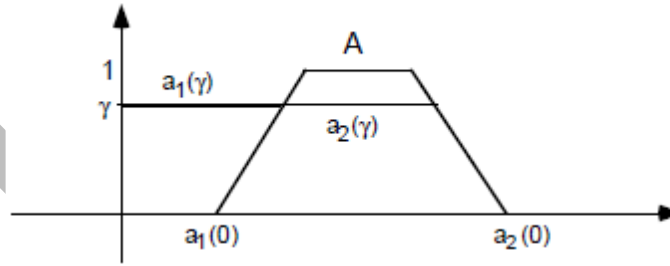


Fuzzy number.

Let A be a fuzzy number. Then $[A]_\gamma$ is a closed convex (compact) subset of \mathbb{R} for all $\gamma \in [0, 1]$. Let us introduce the notations $a_1(\gamma) = \min[A]_\gamma$, $a_2(\gamma) = \max[A]_\gamma$. In other words, $a_1(\gamma)$ denotes the left-hand side and $a_2(\gamma)$ denotes the right-hand side of the γ -cut. It is easy to see that if $\alpha \leq \beta$ then $[A]^\alpha \supset [A]^\beta$.

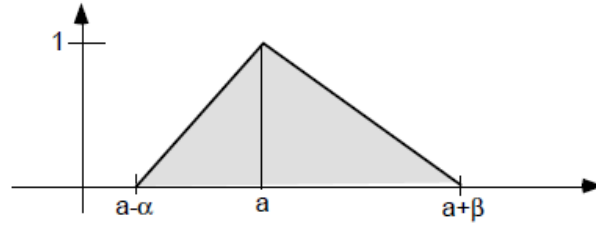
Furthermore, the left-hand side function $a_1: [0, 1] \rightarrow \mathbb{R}$ is monoton increasing and lower semi-continuous, and the right-hand side function $a_2: [0, 1] \rightarrow \mathbb{R}$ is monoton decreasing and upper semi-continuous.

We shall use the notation $[A]_\gamma = [a_1(\gamma), a_2(\gamma)]$. The support of A is the open interval $(a_1(0), a_2(0))$.



If A is not a fuzzy number then there exists a $\gamma \in [0, 1]$ such that $[A]_\gamma$ is not a convex subset of \mathbb{R} .

Definition. (Triangular fuzzy number) A fuzzy set A is called triangular fuzzy number with peak (or center) a , left width $\alpha > 0$ and right width $\beta > 0$ if its membership function has the following form



Triangular fuzzy number.

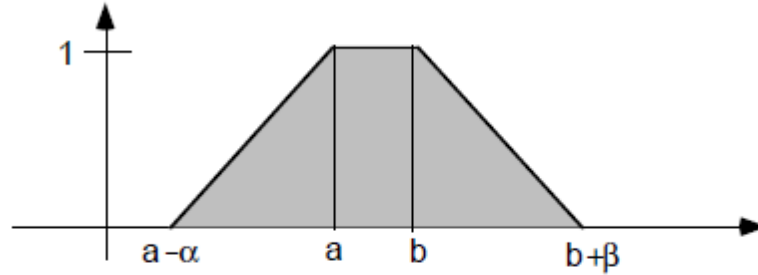
and we use the notation $A = (a, \alpha, \beta)$. It can easily be verified that

$[A]\gamma = [a - (1 - \gamma)\alpha, a + (1 - \gamma)\beta]$, $\forall \gamma \in [0, 1]$. The support of A is $(a - \alpha, b + \beta)$.

A triangular fuzzy number with center a may be seen as a fuzzy quantity “ x is approximately equal to a ”.

Definition .(trapezoidal fuzzy number) A fuzzy set A is called trapezoidal fuzzy number with tolerance interval $[a, b]$, left width α and right width β if its membership function has the following form

$$A(t) = \begin{cases} 1 - (a - t)/\alpha & \text{if } a - \alpha \leq t \leq a \\ 1 & \text{if } a \leq t \leq b \\ 1 - (t - b)/\beta & \text{if } a \leq t \leq b + \beta \\ 0 & \text{otherwise} \end{cases}$$



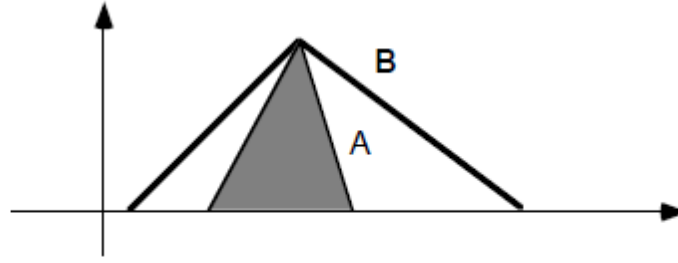
Trapezoidal fuzzy number.

and we use the notation $A = (a, b, \alpha, \beta)$. It can easily be shown that

$[A]\gamma = [a - (1 - \gamma)\alpha, b + (1 - \gamma)\beta]$, $\forall \gamma \in [0, 1]$. The support of A is $(a - \alpha, b + \beta)$.

A trapezoidal fuzzy number may be seen as a fuzzy quantity “ x is approximately in the interval $[a, b]$ ”.

Definition. (subsethood) Let A and B are fuzzy subsets of a classical set X . We say that A is a subset of B if $A(t) \leq B(t)$, $\forall t \in X$.



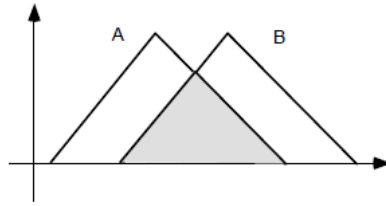
A is a subset of B.

Operations on fuzzy sets

We extend the classical set theoretic operations from ordinary set theory to fuzzy sets. We note that all those operations which are extensions of crisp concepts reduce to their usual meaning when the fuzzy subsets have membership degrees that are drawn from $\{0,1\}$. For this reason, when extending operations to fuzzy sets we use the same symbol as in set theory. Let A and B be fuzzy subsets of a nonempty (crisp) set X .

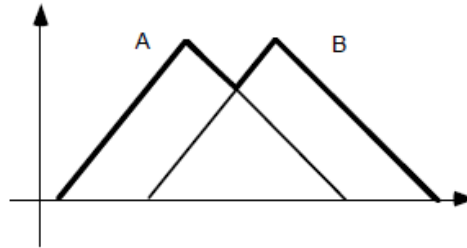
Definition. (intersection) The intersection of A and B is defined as

$$(A \cap B)(t) = \min\{A(t), B(t)\} = A(t) \wedge B(t), \text{ for all } t \in X.$$



Intersection of two triangular fuzzy numbers.

Definition. (union) The union of A and B is defined as $(A \cup B)(t) = \max\{A(t), B(t)\} = A(t) \vee B(t)$, for all $t \in X$.



Union of two triangular fuzzy numbers.

Definition. (complement) The complement of a fuzzy set A is defined as

$$(\neg A)(t) = 1 - A(t)$$

A closely related pair of properties which hold in ordinary set theory are *the law of excluded middle* $A \vee \neg A = X$ and *the law of non-contradiction principle* $A \wedge \neg A = \emptyset$. It is clear that $\neg 1X = \emptyset$ and $\neg \emptyset = 1X$, however, the laws of excluded middle and non-contradiction are not satisfied in fuzzy logic.

2.7 FUZZY LOGIC OBJECTIONS

It would be remarkable if a theory as far-reaching as fuzzy systems did not arouse some objections in the professional community. While there have been generic complaints about the "fuzziness" of the process of assigning values to linguistic terms, perhaps the most cogent criticisms come from Haack. A formal logician, Haack argues that there are only two areas in which fuzzy logic could possibly be demonstrated to be "needed," and then maintains that in each case it can be shown that fuzzy logic is not necessary.

The first area Haack defines is that of the nature of Truth and Falsity: if it could be shown, she maintains, that these are fuzzy values and not discrete ones, then a need for fuzzy logic would have been demonstrated. The other area she identifies is that of fuzzy systems' utility: if it could be demonstrated that generalizing classic logic to encompass fuzzy logic would aid in calculations of a given sort, then again a need for fuzzy logic would exist.

In regards to the first statement, Haack argues that True and False are discrete terms. For example, "The sky is blue" is either true or false; any fuzziness to the statement arises from an imprecise definition of terms, not out of the nature of Truth. As far as fuzzy systems' utility is concerned, she maintains that no area of data manipulation is made easier through the introduction of fuzzy calculus; if anything, she says, the calculations become more complex. Therefore, she asserts, fuzzy logic is unnecessary.

Fox has responded to her objections, indicating that there are three areas in which fuzzy logic can be of benefit: as a "requisite" apparatus (to describe real-world relationships which are inherently fuzzy); as a "prescriptive" apparatus (because some data is fuzzy, and therefore requires a fuzzy calculus); and as a "descriptive" apparatus (because some inferencing systems are inherently fuzzy).

His most powerful arguments come, however, from the notion that fuzzy and classic logics need not be seen as competitive, but complementary. He argues that many of Haack's objections stem from a lack of semantic clarity, and that ultimately fuzzy statements may be translatable into phrases which classical logicians would find palatable.

Exercise:

1. Explain in brief about fuzzy logic.
2. Discuss the history of fuzzy system.

3. Explain fuzzy operations with suitable example.
4. What is subethood? Explain in brief with suitable example.
5. Explain trapezoidal and union operation.
6. Explain intersection and triangular operations.
7. Explain in brief quasi fuzzy number.

DDOL

Chapter 3

Fuzzy Relations and Implications

3.1 Crisp and Fuzzy Relation

A fuzzy relation generalizes these degrees to membership grades. So, a crisp relation is a restricted case of a fuzzy relation.

Crisp relation:

+ degrees or strengths of relation

Fuzzy relation

◦ Cartesian product :

$$\times_{i \in N_n} X_i = \{(x_1, \dots, x_n) \mid x_i \in X_i, \forall i \in N_n\}$$

$$N_n = \{1, 2, \dots, n\}$$

◦ n-ary relation: a subset of $\times_{i \in N_n} X_i$

$$R(X_1, X_2, \dots, X_n) \subseteq X_1 \times X_2 \times \dots \times X_n$$

i.e., $\hat{\quad} \quad \quad \hat{\quad}$

a set the universal set

Characteristic function:

$$\mu_R(x_1, x_1, \dots, x_n) = \begin{cases} 1 & \text{if } (x_1, x_1, \dots, x_n) \in R \\ 0 & \text{otherwise} \end{cases}$$

◦ Binary, Ternary, Quaternary, Quinary, n-ary

Relations

Definition of Relation

A relation among crisp sets X_1, \dots, X_n is a subset of $X_1 \times \dots \times X_n$ denoted as $R(X_1, \dots, X_n)$ or $R(X_i \mid 1 \leq i \leq n)$. So, the relation $R(X_1, \dots, X_n) \subseteq X_1 \times \dots \times X_n$ is set, too. The basic concept of sets can be also applied to relations: • containment, subset, union, intersection, complement Each crisp

relation can be defined by its characteristic function $R(x_1, \dots, x_n) = (1, \text{ if and only if } (x_1, \dots, x_n) \in R, 0, \text{ otherwise.}$

The membership of (x_1, \dots, x_n) in R signifies that the elements of (x_1, \dots, x_n) are related to each other.

Relation as Ordered Set of Tuples

A relation can be written as a set of ordered tuples. Thus $R(X_1, \dots, X_n)$ represents n-dim. membership array $R = [r_{i1}, \dots, r_{in}]$. • Each element of i_1 of R corresponds to exactly one member of X_1 . • Each element of i_2 of R corresponds to exactly one member of X_2 . • And so on...

If $(x_1, \dots, x_n) \in X_1 \times \dots \times X_n$ corresponds to $r_{i1}, \dots, r_{in} \in R$, then $r_{i1}, \dots, r_{in} = (1, \text{ if and only if } (x_1, \dots, x_n) \in R, 0, \text{ otherwise.}$

Fuzzy Relations

The characteristic function of a crisp relation can be generalized to allow tuples to have degrees of membership. • Recall the generalization of the characteristic function of a crisp set!

Then a fuzzy relation is a fuzzy set defined on tuples (x_1, \dots, x_n) that may have varying degrees of membership within the relation. The membership grade indicates strength of the present relation between elements of the tuple. The fuzzy relation can also be represented by an n-dimensional membership array.

Cartesian Product of Fuzzy Sets: n Dimensions

Let $n \geq 2$ fuzzy sets A_1, \dots, A_n be defined in the universes of discourse X_1, \dots, X_n , respectively. The Cartesian product of A_1, \dots, A_n denoted by $A_1 \times \dots \times A_n$ is a fuzzy relation in the product space $X_1 \times \dots \times X_n$. It is defined by its membership function

$$\mu_{A_1 \times \dots \times A_n}(x_1, \dots, x_n) = T(\mu_{A_1}(x_1), \dots, \mu_{A_n}(x_n))$$

whereas $x_i \in X_i, 1 \leq i \leq n$. Usually T is the minimum (sometimes also the product).

Cartesian Product of Fuzzy Sets: 2 Dimensions

A special case of the Cartesian product is when $n = 2$. Then the Cartesian product of fuzzy sets $A \in F(X)$ and $B \in F(Y)$ is a fuzzy relation $A \times B \in F(X \times Y)$ defined by

$$\mu_{A \times B}(x, y) = T[\mu_A(x), \mu_B(y)], \forall x \in X, \forall y \in Y.$$

Subsequences

Consider the Cartesian product of all sets in the family

$$X = \{X_i \mid i \in \mathbb{N}_n = \{1, 2, \dots, n\}\}.$$

For each sequence (n-tuple) $x = (x_1, \dots, x_n) \in \times_{i \in I} N_n X_i$ and each sequence (r-tuple, $r \leq n$) $y = (y_1, \dots, y_r) \in \times_{j \in J} X_j$ where $J \subseteq I$ and $|J| = r$ y is called subsequence of x if and only if $y_j = x_j, \forall j \in J$. $y < x$ denotes that y is subsequence of x .

Projection

Given a relation $R(x_1, \dots, x_n)$. Let $[R \downarrow Y]$ denote the projection of R on Y . It disregards all sets in X except those in the family

$$Y = \{X_j \mid j \in J \subseteq I, n\}.$$

Then $[R \downarrow Y]$ is a fuzzy relation whose membership function is defined on the Cartesian product of the sets in Y

$$[R \downarrow Y](y) = \max_{x > y} R(x).$$

$$R(x).$$

Under special circumstances, this projection can be generalized by replacing the max operator by another t-conorm.

Cylindric Extension

Another operation on relations is called cylindric extension. Let X and Y denote the same families of sets as used for projection. Let R be a relation defined on Cartesian product of sets in family Y . Let $[R \uparrow X \setminus Y]$ denote the cylindric extension of R into sets $X_1, (i \in I, n)$ which are in X but not in Y . It follows that for each x with $x > y$

$$[R \uparrow X \setminus Y](x) = R(y).$$

The cylindric extension • produces largest fuzzy relation that is compatible with projection, • is the least specific of all relations compatible with projection, • guarantees that no information not included in projection is used to determine extended relation.

Example

Consider again the example for the projection. The membership functions of the cylindric extensions of all projections are already shown in the table under the assumption that their arguments are extended to (x_1, x_2, x_3) e.g.

$$[R_{23} \uparrow \{X_1\}](0,0,2) = [R_{23} \uparrow \{X_1\}](1,0,2) = R_{23}(0,2) = 0.2.$$

In this example none of the cylindric extensions are equal to the original fuzzy relation. This is identical with the respective projections. Some information was lost when the given relation was replaced by any one of its projections.

Cylindric Closure

Relations that can be reconstructed from one of their projections by cylindric extension exist. However, they are rather rare. It is more common that relation can be exactly reconstructed • from several of its projections (max), • by taking set intersection of their cylindric extensions (min). The resulting relation is usually called cylindric closure. Let the set of projections $\{P_i \mid i \in I\}$ of a relation on X be given. Then the cylindric closure $\text{cyl}\{P_i\}$ is defined for each $x \in X$ as

$$\text{cyl}\{P_i\}(x) = \min_{i \in I} [P_i \uparrow X \setminus Y_i](x).$$

Y_i denotes the family of sets on which P_i is defined.

Example

Consider again the example for the projection. The cylindric closures of three families of the projections are shown below:

(x_1, x_2, x_3)	$R(x_1, x_2, x_3)$	$\text{cyl}(R_{12}, R_{13}, R_{23})$	$\text{cyl}(R_1, R_2, R_3)$	$\text{cyl}(R_{12}, R_3)$
0 0 0	0.4	0.5	0.9	0.9
0 0 1	0.9	0.9	0.9	0.9
0 0 2	0.2	0.2	0.9	0.9
0 1 0	1.0	1.0	1.0	1.0
0 1 1	0.0	0.5	0.9	0.9
0 1 2	0.8	0.8	1.0	1.0
1 0 0	0.5	0.5	0.9	0.5
1 0 1	0.3	0.5	0.9	0.5
1 0 2	0.1	0.2	0.9	0.5
1 1 0	0.0	0.5	1.0	1.0
1 1 1	0.5	0.5	0.9	0.9
1 1 2	1.0	1.0	1.0	1.0

Motivation and Domain

Binary relations are significant among n -dimensional relations. They are (in some sense) generalized mathematical functions. On the contrary to functions from X to Y , binary relations $R(X, Y)$ may assign to each element of X two or more elements of Y . Some basic operations on functions, e.g. inverse and composition, are applicable to binary relations as well. Given a fuzzy relation $R(X, Y)$. Its domain $\text{dom}R$ is the fuzzy set on X whose membership function is defined for each $x \in X$ as

$$\text{dom}R(x) = \max_{y \in Y}$$

$$R(x, y),$$

i.e. each element of X belongs to the domain of R to a degree equal to the strength of its strongest relation to any $y \in Y$.

Range and Height

The range ran of $R(X, Y)$ is a fuzzy relation on Y whose membership function is defined for each $y \in Y$ as

$$\text{ran}R(y) = \max_{x \in X}$$

$R(x,y)$,

i.e. the strength of the strongest relation which each $y \in Y$ has to an $x \in X$ equals to the degree of membership of y in the range of R . The height h of $R(X,Y)$ is a number defined by

$$h(R) = \max_{y \in Y}$$

$$\max_{x \in X} R(x,y)$$

$R(x,y)$.

$h(R)$ is the largest membership grade obtained by any pair $(x,y) \in R$.

Representation and Inverse

Consider e.g. the membership matrix $R = [r_{xy}]$ with $r_{xy} = R(x,y)$.

Its inverse $R^{-1}(Y,X)$ of $R(X,Y)$ is a relation on $Y \times X$ defined by

$$R^{-1}(y,x) = R(x,y), \forall x \in X, \forall y \in Y.$$

$R^{-1} = [r^{-1}_{xy}]$ representing $R^{-1}(y,x)$ is the transpose of R for $R(X,Y)$

$$(R^{-1})^{-1} = R, \forall R.$$

Standard Composition

Consider the binary relations $P(X,Y)$, $Q(Y,Z)$ with common set Y . The standard composition of P and Q is defined as

$$(x,z) \in P \circ Q \iff \exists y \in Y : \{(x,y) \in P \wedge (y,z) \in Q\}.$$

In the fuzzy case this is generalized by

$$[P \circ Q](x,z) = \sup_{y \in Y}$$

$$\min\{P(x,y), Q(y,z)\}, \forall x \in X, \forall z \in Z.$$

If Y is finite, sup operator is replaced by max. Then the standard composition is also called max-min composition.

Inverse of Standard Composition

The inverse of the max-min composition follows from its definition:

$$[P(X,Y) \circ Q(Y,Z)]^{-1} = Q^{-1}(Z,Y) \circ P^{-1}(Y,X).$$

Its associativity also comes directly from its definition:

$$[P(X,Y) \circ Q(Y,Z)] \circ R(Z,W) = P(X,Y) \circ [Q(Y,Z) \circ R(Z,W)].$$

Note that the standard composition is not commutative. Matrix notation: $[r_{ij}] = [p_{ik}] \circ [q_{kj}]$ with $r_{ij} = \max_k \min(p_{ik}, q_{kj})$.

Example

$$P \circ Q = R$$

$$\begin{bmatrix} .3 & .5 & .8 \\ 0 & .7 & 1 \\ .4 & .6 & .5 \end{bmatrix} \circ \begin{bmatrix} .9 & .5 & .7 & .7 \\ .3 & .2 & 0 & .9 \\ 1 & 0 & .5 & .5 \end{bmatrix} = \begin{bmatrix} .8 & .3 & .5 & .5 \\ 1 & .2 & .5 & .7 \\ .5 & .4 & .5 & .5 \end{bmatrix}$$

For instance:

$$r_{11} = \max\{\min(p_{11}, q_{11}), \min(p_{12}, q_{21}), \min(p_{13}, q_{31})\} = \max\{\min(.3, .9), \min(.5, .3), \min(.8, 1)\} = .8$$

$$r_{32} = \max\{\min(p_{31}, q_{12}), \min(p_{32}, q_{22}), \min(p_{33}, q_{32})\} = \max\{\min(.4, .5), \min(.6, .2), \min(.5, 0)\} = .4$$

Example: Types of Airplanes (Speed, Height, Type)

Consider the following fuzzy relations for airplanes: • relation A between maximal speed and maximal height, • relation B between maximal height and the type.

A	h_1	h_2	h_3
s_1	1	.2	0
s_2	.1	1	0
s_3	0	1	1
s_4	0	.3	1

B	t_1	t_2
h_1	1	0
h_2	.9	1
h_3	0	.9

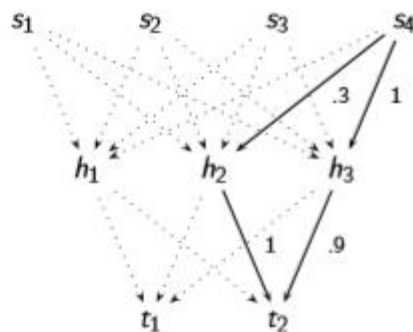
Example (cont.)

matrix multiplication scheme

A	o	B		
			1	0
			.9	1
			0	.9
1	.2	0	1	.2
.1	1	0	.9	1
0	1	1	.9	1
0	.3	1	.3	.9

$A \circ B$ speed-type relation

flow scheme



$$(A \circ B)(s_4, t_2) = \max\{\min\{.3, 1\}, \min\{.1, .9\}\} = .9$$

Relational Join

A similar operation on two binary relations is the relational join. It yields triples (whereas composition returned pairs). For $P(X, Y)$ and $Q(Y, Z)$, the relational join $P * Q$ is defined by

$$[P * Q](x,y,z) = \min\{P(x,y), Q(y,z)\}, \forall x \in X, \forall y \in Y, \forall z \in Z.$$

Then the max-min composition is obtained by aggregating the join by the maximum:

$$[P \circ Q](x,z) = \max_{y \in Y}$$

$$[P * Q](x,y,z), \forall x \in X, \forall z \in Z.$$

Example

The join $S = P * Q$ of the relations P and Q has the following membership function (shown below on left-hand side). To convert this join into its corresponding composition $R = P \circ Q$ (shown on right-hand side), the two indicated pairs of $S(x,y,z)$ are aggregated using max.

x	y	z	$\mu_{S(x,y,z)}$
1	a	α	.6
1	a	β	.7*
1	b	β	.5*
2	a	α	.6
2	a	β	.8
3	b	ρ	.1
4	b	β	.4*
4	c	β	.3*

x	z	$\mu_R(x,z)$
1	α	.6
1	β	.7
2	α	.6
2	β	.8
3	ρ	.1
4	β	.4

For instance,

$$R(1,\beta) = \max\{S(1,a,\beta), S(1,b,\beta)\} = \max\{.7, .5\} = .7$$

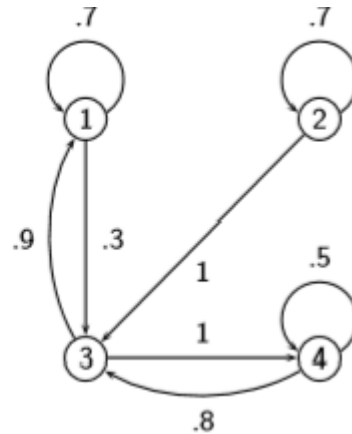
Binary Relations on a Single Set

It is also possible to define crisp or fuzzy binary relations among elements of a single set X . Such a binary relation can be denoted by $R(X,X)$ or $R(X^2)$ which is a subset of $X \times X = X^2$. These relations are often referred to as directed graphs which is also an representation of them. • Each element of X is represented as node. • Directed connections between nodes indicate pairs of $x \in X$ for which the grade of the membership is nonzero. • Each connection is labeled by its actual membership grade of the corresponding pair in R .

Example

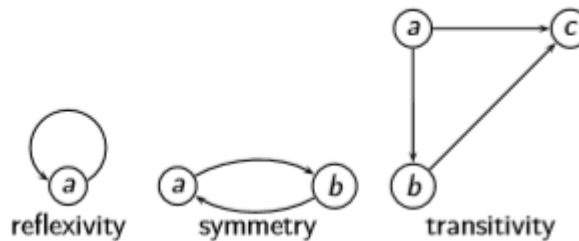
An example of $R(X,X)$ defined on $X = 1,2,3,4$. Two different representation are shown below.

	1	2	3	4
1	.7	0	.3	0
2	0	.7	1	0
3	.9	0	0	1
3	0	0	.8	.5



Properties of Crisp Relations

A crisp relation $R(X,X)$ is called • reflexive if and only if $\forall x \in X : (x,x) \in R$, • symmetric if and only if $\forall x,y \in X : (x,y) \in R \leftrightarrow (y,x) \in R$, • transitive if and only if $(x,z) \in R$ whenever both $(x,y) \in R$ and $(y,z) \in R$ for at least one $y \in X$.



Properties of Fuzzy Relations

These properties can be extended for fuzzy relations. So one can define them in terms of the membership function of the relation. A fuzzy relation $R(X,X)$ is called • reflexive if and only if $\forall x \in X : R(x,x) = 1$, • symmetric if and only if $\forall x,y \in X : R(x,y) = R(y,x)$, • transitive if it satisfies

$$R(x,z) \geq \max_{y \in Y}$$

$$\min\{R(x,y), R(y,z)\}, \forall (x,z) \in X^2.$$

Note that a fuzzy binary relation that is reflexive, symmetric and transitive is called fuzzy equivalence relation.

· Representation of a relation

$$R(\bar{X}_1, \dots, \bar{X}_n)$$

□ $(r_{i_1, i_2, \dots, i_n})$: n-D membership array

$$r_{i_1, i_2, \dots, i_n} = \begin{cases} 1 & \text{iff } (x_1, \dots, x_n) \in R \\ 0 & \text{Otherwise} \end{cases}$$

○ Example 3.1 :

$$R(\bar{X} \bar{Y}) Z Y_1 Y_2 Y_3 Y_4 Z_2 Z_1 Z_3 Z_4 Z_5$$

$$R(\{\bar{X}_1, \bar{X}_3\} \dots, \bar{X}_i \mid J \subseteq N_n, Y \leq X)$$

$$\{\bar{X}_i \mid j \in J \subseteq N_n\} \Rightarrow R_2 = \frac{1}{X, a, *} + \frac{1}{X, a, \$} + \frac{1}{Y, b, *} + \frac{1}{Y, a, \$} + \frac{0.8}{X, b, *} + \frac{0.8}{X, b, \$} + \frac{0.8}{Y, b, *} + \frac{0.8}{Y, b, \$}$$

$$\Rightarrow y = \{\bar{X}_1, \bar{X}_3\}$$

$$Y_j = X_j \mid \forall j \in J$$

$$[R \uparrow X - Y] : [R_2 \uparrow \{\bar{X}_1, \bar{X}_3\}]$$

$$\Rightarrow y = \{\bar{X}_2\}, X - Y = \{\bar{X}_1, \bar{X}_3\} = \{(*, *), (x, \$), (Y, *), (Y, \$)\}$$

$$(x) = R(y) = [R \downarrow \{\bar{X}_i\}](Y) = \max_{x \succ y} R(x) R_{ij}$$

$$\langle Y_j \mid j \in J \rangle \in \bigcup_{j \in J} \bar{X}_j$$

$$R_{2,3} = \frac{1}{a, *} + \frac{0.7}{a, \$} + \frac{0.4}{b, *} + \frac{0.8}{b, \$}.$$

$$\frac{0.9}{X, a, *} + \frac{0.4}{X, b, *} + \frac{1}{Y, a, *} + \frac{0.7}{Y, a, \$} + \frac{0.8}{Y, b, \$}$$

$$\Rightarrow R_{1,2} = \frac{0.9}{X, a} + \frac{0.4}{X, b} + \frac{1}{Y, a} + \frac{0.8}{Y, b} \quad R_{1,3} = \frac{0.9}{X, *} + \frac{0}{X, \$} + \frac{1}{Y, *} + \frac{0.8}{Y, \$}$$

$$0 \leq R(x_1, x_2, \dots, x_n) \leq 1 \frac{1}{(NY, Beijing)} + \frac{0}{(NY, NY)} + \frac{0.6}{(NY, London)} + \frac{0.9}{(Paris, Beijing)} + \frac{0.7}{(Paris, NY)} + \frac{0.8}{(Paris, London)}$$

$$\bar{X} = \{\text{English, French}\}, \bar{Y} = \{\text{dollar, pound, franc, mark}\}$$

$Z=\{US, France, Canada, Britain, Germany\}$

$R(\bar{X} \bar{Y} Z) = \{(English, dollar, US), (French, franc, France)$

$(English, dollar, Canada), (French, dollar, Canada)$

$(English, pound, Britain)\}$

Y_1 Dollar 1 0 1 0 0 Dollar 0 0 1 0 0

Y_2 Pound 0 0 0 1 0 Pound 0 0 0 0 0

Y_3 Franc 0 0 0 0 0 Franc 0 1 0 0 0

Y_4 Mark 0 0 0 0 0 Mark 0 0 0 0 0

US Fran Can Brit Ger

US Fran Can Brit Ger

$Z_1 Z_2 Z_3 Z_4 Z_5$ $Z_1 Z_2 Z_3 Z_4 Z_5$

English

Franch

$X_1 X_2$

• Fuzzy Relations

Cartesian Product : $\bar{X}_1 \times \bar{X}_2 \times \dots \times \bar{X}_n$

tuples : (x_1, x_2, \dots, x_n)

membership grade :

$$0 \leq R(x_1, x_2, \dots, x_n) \leq 1$$

Example 3.2: Binary relation R : represents the concept “very far”

$\bar{X} = \{ \text{New York, Paris} \}$

$\bar{Y} = \{ \text{Beijing, New York, London} \}$

Relation in list notation

$$R(\bar{X}, \bar{Y}) = \frac{1}{(NY, Beijing)} + \frac{0}{(NY, NY)} + \frac{0.6}{(NY, London)} + \frac{0.9}{(Paris, Beijing)} + \frac{0.7}{(Paris, NY)} + \frac{0.3}{(Paris, London)}$$

Relation in membership array

	NY	Paris
Beijing	1	0.9
NY	0	0.7
London	0.6	0.3

· Ordinary fuzzy relation

with valuation set [0,1]

L-fuzzy relation

With ordered valuation set L

3.2 Projection and Cyclindric Extensions

· set family $X = \{\bar{X}_i \mid i \in N_n\}$

Let $X = \langle x_i \mid i \in N_n \rangle \in \prod_{i \in J} \bar{X}_i$

Let $Y = \langle Y_j \mid j \in J \rangle \in \prod_{j \in J} \bar{X}_j$

Where $J \subseteq N_n, |J| = r$

Y a subsequence of X, $Y \leq X$

iff $Y_j = X_j \forall j \in J$

⊙ Projection : $[R \downarrow y]$ the projection of R on Y

$R(\bar{X}_1, \bar{X}_2, \dots, \bar{X}_n)$: a relation

$Y = \{\bar{X}_i \mid i \in J \subseteq N_n\}$

$[R \downarrow y]$: a fuzzy relation (set)

$[R \downarrow y](Y) = \max_{x \succ y} R(x)$

※ max can be generalized by other t-conorms

· Example 3.3 $\bar{X}_1 = \{X, Y\}$, $\bar{X}_2 = \{a, b\}$, $\bar{X}_3 = \{*, \$\}$

$$R(\bar{X}_1, \bar{X}_2, \bar{X}_3) = \frac{0.9}{X, a, *} + \frac{0.4}{X, b, *} + \frac{1}{Y, a, *} + \frac{0.7}{Y, a, \$} + \frac{0.8}{Y, b, \$}$$

Let $R_{ij} = [R \downarrow \{\bar{X}_i, \bar{X}_j\}]$, $R_i = [R \downarrow \{\bar{X}_i\}]$

$$\Rightarrow R_{1,2} = \frac{0.9}{X, a} + \frac{0.4}{X, b} + \frac{1}{Y, a} + \frac{0.8}{Y, b}$$

$$R_{1,3} = \frac{0.9}{X, *} + \frac{0}{X, \$} + \frac{1}{Y, *} + \frac{0.8}{Y, \$}$$

$$R_{2,3} = \frac{1}{a, *} + \frac{0.7}{a, \$} + \frac{0.4}{b, *} + \frac{0.8}{b, \$}$$

$$R_1 = \frac{0.9}{*} + \frac{1}{y}$$

$$R_2 = \frac{1}{a} + \frac{0.8}{b}$$

$$R_3 = \frac{1}{*} + \frac{0.8}{\$}$$

⊙Cylindric Extension $[R \downarrow X - Y]$ the CE of R into X-Y

X-Y : sets \bar{X}_i that are in X but are not in Y

$$[R \downarrow X - Y](x) = R(y)$$

R: a relation defined on Y

·Example 3.4 (Refer to example 3.3)

$$\text{Let } X = \{\bar{X}_1, \bar{X}_2, \bar{X}_3\}$$

$$\text{And } R = R_{1,2} \Rightarrow y = \{\bar{X}_1, \bar{X}_3\}$$

$$\therefore X - Y = \bar{X}_3 = \{*, \$\}$$

$$\text{From example 3.3 } R_{1,2} = \frac{0.9}{X,a} + \frac{0.4}{X,b} + \frac{1}{Y,a} + \frac{0.8}{Y,b}$$

$$\therefore [R \uparrow X - Y] = [R_{1,2} \uparrow \{\bar{X}_3\}] =$$

$$\frac{0.9}{X,a,*} + \frac{0.9}{X,a,\$} + \frac{0.4}{X,b,*} + \frac{0.4}{X,b,\$} + \frac{1}{Y,a,\$} + \frac{1}{Y,a,\$} + \frac{0.8}{Y,b,*} + \frac{0.8}{Y,b,\$}$$

$$[R \uparrow X - Y] : [R_{12} \uparrow \{\bar{X}_3\}] [R_{13} \uparrow \{\bar{X}_2\}] [R_{23} \uparrow \{\bar{X}_1\}] [R_1 \uparrow \{\bar{X}_2, \bar{X}_2\}] [R_2 \uparrow \{\bar{X}_1, \bar{X}_3\}]$$

$$[R_3 \uparrow \{\bar{X}_1, \bar{X}_2\}]$$

$$\text{Consider } [R \uparrow X - Y] = [R_2 \uparrow \{\bar{X}_1, \bar{X}_3\}]$$

$$\Rightarrow y = \{\bar{X}_2\}, X - Y = \{\bar{X}_1, \bar{X}_3\} = \{(*, *), (x, \$), (Y, *), (Y, \$)\}$$

$$\{x, y\} \{x, \$\}$$

$$R = R_2 = \frac{1}{a} + \frac{0.8}{b}$$

$$\therefore [R_2 \uparrow \{\bar{X}_1, \bar{X}_3\}] = \frac{1}{X,a,*} + \frac{1}{X,a,\$} + \frac{1}{Y,b,*} + \frac{1}{Y,a,\$} + \frac{0.8}{X,b,*} + \frac{0.8}{X,b,\$} + \frac{0.8}{Y,b,*} + \frac{0.8}{Y,b,\$}$$

3.4 Cyclindric closure

-A relation may be exactly reconstructed from several of its projections by taking the set intersection of their cylindric extensions $\{P_i \mid i \in I\}$: a set of projections of a relation on X

$$\Rightarrow cyl\{P_i\}(X) = \min_{i \in I} [P_i \uparrow X - Y_i](X) \geq R$$

Y_i : The family of sets on which P_i is defined.

• Example :

$$Cyl\{R_{1,2}, R_{1,3}, R_{2,3}\} = \frac{0.9}{x, a, *} + \frac{0.4}{x, b, *} + \frac{1}{y, a, *} + \frac{0.7}{y, a, \$} + \frac{0.4}{y, b, *} + \frac{0.8}{y, b, *}$$

Refer to the original relation $R(\bar{X}_{-1}, \bar{X}_{-2}, \bar{X}_{-3})$ in example 3.3.

It is not fully reconstructable from its projections because of ignorance of R_1 , R_2 , and R_3 .

3.3. Binary Relations $R(\bar{X}, \bar{Y})$

$\bar{X} \neq \bar{Y}$: bipartite graph

$\bar{X} = \bar{Y}$: directed graph

- Representations

i, matrices $R = [r_{ij}]$, where $r_{ij} = R(x_i, y_j)$

ii, sagittal diagrams

Examples :

i) $y_1 \quad y_2 \quad y_3 \quad y_4 \quad y_5$

x_1	.9	1	0	0	0
x_2	0	.4	0	0	0
x_3	0	0	1	.2	0
x_4	0	0	0	0	.4
x_5	0	0	0	0	.5
x_6	0	0	0	0	.2

ii)

3.9

- Domain : $\text{dom } R$

Crisp – $\text{dom } R = \{x \in X \mid (x, y) \in R, \exists y \in Y\}$

Fuzzy – $\text{dom } R(x) = \max_{y \in Y} R(x, y)$

The domain of a fuzzy relation $R(x, y)$ is a fuzzy set on X ; $\text{dom } R(x)$ is its membership function.

e.g. $\text{dom } R(X_1) = \max(0.9, 1) = 1$

- Range : $\text{ran } R$

Crisp – $\text{ran } R = \{y \in Y \mid (x, y) \in R, \exists x \in X\}$

Fuzzy – $\text{ran } R(y) = \max_{x \in X} R(x, y)$

e.g. $\text{ran } R(y_5) = \max(0.4, 0.5, 0.2) = 0.5$

• Height : $h(R) = \max_{y \in Y} \max_{x \in X} R(x, y)$

e.g., $h(R) = 1$ normal fuzzy relation

• Inverse : $R^{-1}(Y, X)$

$$R^{-1}(y, x) = R(x, y)$$

$$\therefore R^{-1} = R^T, (R^{-1})^{-1} = R$$

$$\text{e.g. } R = \begin{bmatrix} 0.3 & 0.2 \\ 0 & 1 \\ 0.6 & 0.4 \end{bmatrix}$$

$$R^{-1} = R^T = \begin{bmatrix} 0.3 & 0 & 0.6 \\ 0.2 & 1 & 0.4 \end{bmatrix}$$

• Composition : $R(X, Z) = P(X, Y) \circ Q(Y, Z)$

$$R(x, z) = [P \circ Q](x, z) = \max_{y \in Y} \min[P(x, y), Q(y, z)]$$

Max-min composition

$$\text{Properties : } \begin{cases} P \circ Q \neq Q \circ P \\ (P \circ Q)^{-1} \neq Q^{-1} \circ P^{-1} \\ (P \circ Q) \circ R = P \circ (Q \circ R) \end{cases}$$

$$\text{Matric form : } [r_{ij}] = [p_{ik}] \circ [q_{kj}]$$

$$\text{Where } [r_{ij}] = \max_k \min(p_{ik}, q_{kj})$$

3.11

$$R(x, z) = [P \circ Q](x, z) = \max_{y \in Y} [P(x, y) \cdot Q(y, z)]$$

max-product composition

$$\text{matrix form } [r_{ij}] = [p_{ik}] \circ [q_{kj}]$$

$$\text{Where } [r_{ij}] = \max(p_{ik}, q_{kj})$$

• Example

$$\begin{bmatrix} 0.3 & 0.5 & 0.8 \\ 0.0 & 0.7 & 1.0 \\ 0.4 & 0.6 & 0.5 \end{bmatrix} \circ \begin{bmatrix} 0.9 & 0.5 & 0.7 & 0.7 \\ 0.3 & 0.2 & 0.0 & 0.9 \\ 1.0 & 0.0 & 0.5 & 0.5 \end{bmatrix}$$

$$\text{Max-min} = \begin{bmatrix} 0.8 & 0.3 & 0.5 & 0.5 \\ 1.0 & 0.2 & 0.5 & 0.7 \\ 0.5 & 0.4 & 0.5 & 0.6 \end{bmatrix}$$

$$\text{Max-prod} = \begin{bmatrix} 0.8 & 0.15 & 0.4 & 0.45 \\ 1.0 & 0.14 & 0.5 & 0.63 \\ 0.5 & 0.2 & 0.28 & 0.54 \end{bmatrix}$$

• Relational join : $R(X, Y, Z) = P(X, Y) * Q(Y, Z)$

$$R(x, y, z) = [P * Q](x, y, z) = \min[P(x, y), Q(y, z)]$$

✖The max-min composition can be obtained by aggregating appropriate elements of the corresponding join.

3.12 • Example

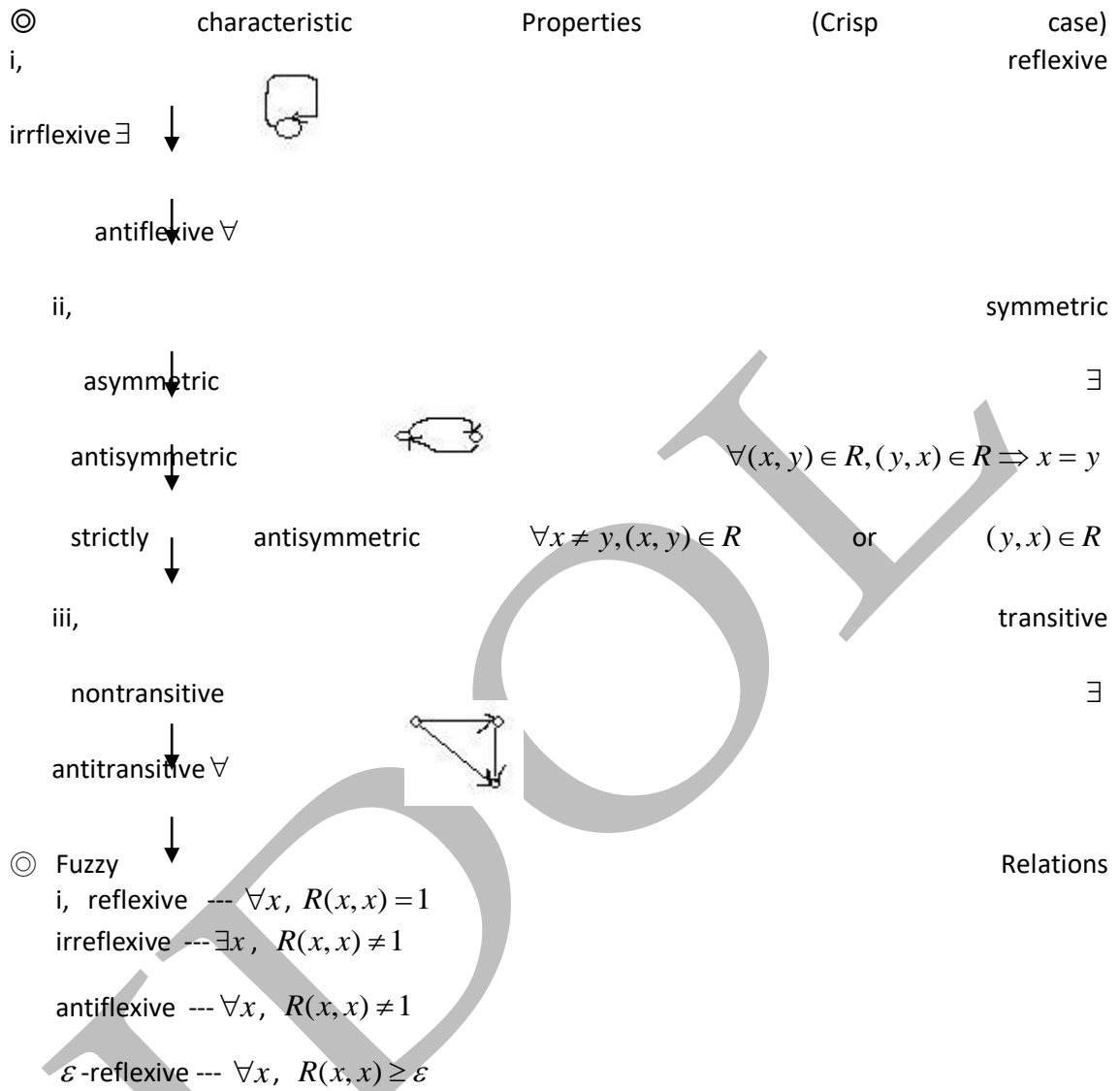
$$\circledast [P \circ Q](x, z) = \max_{y \in Y} [P * Q](x, y, z)$$

5.4 Binary Relation on a Simple Set

• Representations

DRAFT

3-13



3-14

ii, symmetric -- $\forall x, y, R(x, x) = R(y, x)$

asymmetric -- $\exists x, y, R(x, x) \neq R(y, x)$

antisymmetric

$$\text{--} \begin{cases} R(x, x) > 0 \\ R(y, x) > 0 \end{cases} \Rightarrow x = y$$

iii, max-min transitive -- $\forall x, z$

$$R(x, z) \geq \max_{y \in \bar{Y}} \min [R(x, y), R(y, z)]$$

max-product transitive -- $\exists x, z$

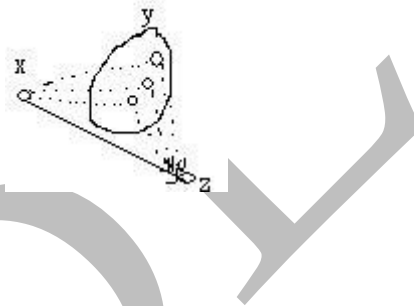
$$R(x, z) \geq \max_{y \in \bar{Y}} \min [R(x, y) \bullet R(y, z)]$$

nontransitive -- $\exists x, z$

$$R(x, z) < \max_{y \in \bar{Y}} \min [R(x, y), R(y, z)]$$

antitransitive -- $\forall x, z$

$$R(x, z) < \max_{y \in \bar{Y}} \min [R(x, y), R(y, z)]$$



© Example 3.7 R:very near
 \Rightarrow reflexive, symmetric, nontransitive

Summary

	Reflexive	Antireflexive	Symmetric	Antisymmetric	Transitive
Crisp: equivalence; Fuzzy: similarity					
Quasi-equivalence					
Compatibility or Tolerance					
Partial ordering					
Preordering or Quasi-ordering					
Strict ordering					

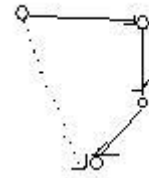
Figure3.6 Some important types of binary relation $R(X,X)$

transitive closure: $R_T(\bar{X})$

Algorithm for computing R_T

1. $R' = R \cup (R \circ R)$
2. If $R' \neq R$, Let $R' = R$, go to step1
3. Stop, $R_T = R'$

Where \cup : component-wise max



©Example 3.8

$$R = \begin{bmatrix} 0.7 & 0.5 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.0 & 1.0 \\ 0.0 & 0.4 & 0.0 & 0.0 \\ 0.0 & 0.0 & 0.8 & 0.0 \end{bmatrix}$$

$$\text{Step1: } R \circ R = \begin{bmatrix} 0.7 & 0.5 & 0.0 & 0.5 \\ 0.0 & 0.0 & 0.8 & 0.0 \\ 0.0 & 0.0 & 0.0 & 0.4 \\ 0.0 & 0.4 & 0.0 & 0.0 \end{bmatrix}$$

$$R \cup (R \circ R) = \begin{bmatrix} 0.7 & 0.5 & 0.0 & 0.5 \\ 0.0 & 0.0 & 0.8 & 1.0 \\ 0.0 & 0.4 & 0.0 & 0.4 \\ 0.0 & 0.4 & 0.8 & 0.0 \end{bmatrix} = R'$$

Step2: $\because R' \neq R$, Let $R = R'$

repeat step1

$$R \circ R = \begin{bmatrix} 0.7 & 0.5 & 0.5 & 0.5 \\ 0.0 & 0.4 & 0.8 & 0.4 \\ 0.0 & 0.4 & 0.4 & 0.4 \\ 0.0 & 0.4 & 0.4 & 0.4 \end{bmatrix}$$

$$R \cup (R \circ R) = \begin{bmatrix} 0.7 & 0.5 & 0.5 & 0.5 \\ 0.0 & 0.4 & 0.8 & 1.0 \\ 0.0 & 0.4 & 0.4 & 0.4 \\ 0.0 & 0.4 & 0.8 & 0.0 \end{bmatrix} = R'$$

Step3: $\because R' \neq R$, Let $R = R'$

repeat step1

$$R' = \begin{bmatrix} 0.7 & 0.5 & 0.5 & 0.5 \\ 0.0 & 0.4 & 0.8 & 1.0 \\ 0.0 & 0.4 & 0.4 & 0.4 \\ 0.0 & 0.4 & 0.8 & 0.4 \end{bmatrix} = R$$

Step4: Stop

$$R_T = R'$$

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3.5 Fuzzy Equivalence Relation

⊙ Crisp binary relation

equivalence: reflexive, symmetric, and transitive

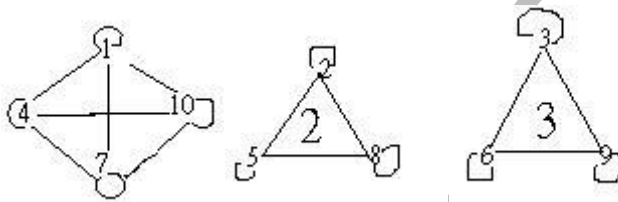
equivalence classes

partition: X/R

⊙ Example 3.9:

$$\bar{X} = \{1, 2, \dots, 10\}$$

$$R(\bar{X} \times \bar{X}) = \{(x, y) \mid x, y \text{ have the same remainder when divided by } 3\}$$



R : reflexive, symmetric, transitive

⇒ equivalence

$$\text{partition } \bar{X}/R = \{(1, 4, 7, 10), (2, 5, 8), (3, 6, 9)\}$$

©Fuzzy Binary Relation

◦ Fuzzy

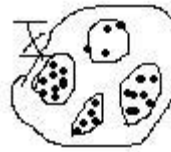
Similarity relation equivalence relation

Similarity classes equivalence classes

2 Interpretations of a similarity relation:

1. Group similar elements into crisp classes

whose members are similar to each other
to some specified degree.



2. $\forall x \in \bar{X}$, associate a fuzzy set A_x

defined on \bar{X} .



◦ a fuzzy relation $R = \bigcup_{\alpha \in [0,1]} \alpha \square^{\alpha} R$ (Theorem 2.5, Eqs.(2.1)(2.2))

If R : Similarity relation,

$\Rightarrow {}^{\alpha}R$: equivalence relation

◦ Let $\Pi({}^{\alpha}R)$: the partition of \bar{X} w.r.t. ${}^{\alpha}R$

$\Rightarrow \Pi(R) = \{\Pi({}^{\alpha}R) \mid \alpha \in (0,1]\}$

$\Pi({}^{\alpha}R)$: nested, i.e.,

$\Pi({}^{\alpha}R)$: are refinement of $\Pi({}^{\beta}R)$ iff $\alpha \geq \beta$

Prove that :A fuzzy relation $R: X \rightarrow X$ is a similarity relation, then ${}^{\alpha}R$ is a equivalent relation

Pf :: R : a similar relation

$$\begin{aligned} \therefore R : \text{ reflexive, i.e., } & \left\{ \begin{aligned} \forall x \in X, R(x, x) &= 1 \\ \text{symmetric, i.e., } \forall x, y \in X, R(x, y) &= R(y, x) \end{aligned} \right. \\ \text{transitive, i.e., } \forall x, z \in X^2, R(x, z) &\geq \max_{y \in Y} [R(x, y), R(y, z)] \end{aligned}$$

i, ${}^{\alpha}R$: reflexive

$$\therefore \forall x \in X, R(x, x) = 1 \geq \alpha \in [0, 1], (x, x) \in {}^{\alpha}R$$

${}^{\alpha}R$: reflexive

ii, ${}^{\alpha}R$: symmetric

$\therefore R$: symmetric

$$\forall x, y \in Z, R(x, y) = R(y, x)$$

$$\text{Let } R(x, y) = R(y, x) = \beta$$

$$\text{Then } \beta \geq \alpha \text{ or } \beta < \alpha$$

$$\text{a, if } \beta \geq \alpha \Rightarrow (x, y), (y, x) \in {}^{\alpha}R$$

$$\text{b, if } \beta < \alpha \Rightarrow (x, y), (y, x) \notin {}^{\alpha}R$$

iii, ${}^{\alpha}R$: transitive

$\therefore R$: transitive

$$\forall x, z \in X^2, R(x, z) \geq \max_{y \in Y} [R(x, y), R(y, z)]$$

$$\text{Let } R(x, y) = \beta_1, R(y, z) = \beta_2$$

$$\text{Assume } \beta_1 < \beta_2$$

$$\text{Then } \alpha \leq \beta_1 < \beta_2, \beta_1 < \alpha < \beta_2, \text{ or } \beta_1 < \beta_2 \leq \alpha$$

$$\text{a. if } \alpha \leq \beta_1 < \beta_2 \Rightarrow (x, y) {}^{\alpha}R, (y, z) \in {}^{\alpha}R \text{ --- (A)}$$

$$\because R(x, z) \geq \max_{y \in Y} [R(x, y), R(y, z)] \geq \min[\beta_1, \beta_2] = \beta_1 \geq \alpha$$

$$\therefore R(x, z) \geq \alpha, (x, z) \in {}^\alpha R \text{ --- (B)}$$

$$(A), (B) \Rightarrow {}^\alpha R : \text{transitive}$$

$$\text{b. if } \beta_1 < \alpha < \beta_2$$

$$\Rightarrow (x, y) \notin {}^\alpha R, (y, z) \in {}^\alpha R, \text{ don't care } (x, z)$$

$$\text{c. if } \beta_1 < \beta_2 \leq \alpha$$

$$\Rightarrow (x, y) \notin {}^\alpha R, (y, z) \notin {}^\alpha R, \text{ don't care } (x, z)$$

Example 3.10 : $R(X, X)$: a fuzzy relation

R : reflexive, symmetric, transitive ($\because R' = R \cup (R \circ R) = R$)

\therefore level set : $A_R = \{0.0, 0.4, 0.5, 0.8, 0.9, 1.0\}$

There are five nested partition ${}^\alpha \pi$'s

- The similarity class for each element is a fuzzy set defined by the row of the membership matrix corresponds to that element

Example : see Example 3.10

$$\text{For c : } \frac{0}{a} + \frac{0}{b} + \frac{1}{c} + \frac{0}{d} + \frac{1}{e} + \frac{0.9}{f} + \frac{0.5}{g}$$

$$\text{For e : } \frac{0}{a} + \frac{0}{b} + \frac{1}{c} + \frac{0}{d} + \frac{1}{e} + \frac{0.9}{f} + \frac{0.5}{g}$$

$\therefore c$ and e are similar at any level α

3.6 Compatibility Relations ---- reflexive , symmetric

compatibility
 Alternatives : tolerance
 proximity

} relation

- Crisp case :
 Maximal compatibility classes – not properly contained within any other compatibility class

Complete cover – all the maximal compatibility classes

- Fuzzy case :
 α -compatibility class ---- a subset A of X ,
 s.t. $\forall x, y \in A, \text{ if } \exists R(x, y) \Rightarrow R(x, y) \geq \alpha, R : \text{fuzzy compatibility relation}$

maximal α ---- compatibility classes

complete α -cover

Example 3.11 : $R(X, X)$: a fuzzy relation

$\therefore R$: reflexive , symmetric

$\therefore a$ compatibility relation

$\therefore A_R = \{0.0, 0.4, 0.5, 0.7, 0.8, 1.0\}$

\Rightarrow the complete α -covers

$$\partial = 0.5$$

$$\frac{0.7}{b} + \frac{0.9}{d}$$

$$\begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} \begin{bmatrix} 1 & 0.7 & 0 & 1 & 0.7 \\ 0 & 1 & 0 & 0.9 & 0 \\ 0.5 & 0.7 & 1 & 1 & 0.8 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0.1 & 0 & 0.9 & 1 \end{bmatrix}$$

$${}^{0.5}R = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

e.q. \therefore $\partial = 0.4$

$$\Rightarrow x \leq y \Rightarrow xx = y \forall y \neq xx \in X(x, y) \rightarrow \leq \in XS = \{x_1, x_2\} \subseteq Xy \in Ax \geq y$$

$$(x, y)y \in X \exists (x \leq y, \text{or} \Leftrightarrow y \leq x) Ax \in XA \subseteq Xx \neq yR_{\leq[x]}(y) = R(y, x)$$

$$x \in U(R, A)(x) = \bigcap_{x \in A} R_{\geq[x]} \forall y \in$$

$${}^{0.4}R = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Look for complete subgraphs

$(1,2), (3,4,5), (4,5,6,7), (5,8), (9)$

$(3,4), (4,5,6), (4,5,7), (3,5), (5,6)$

$(4,5), (5,6,7), (4,6,7), (4,6), (6,7)$

\therefore maximal compatible classes (the complete 0.4-cover):

$(1,2), (3,4,5), (4,5,6,7), (5,8), (9)$

These do not partition X.

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e.g. $\partial=0.5$

$${}^{0.5}R = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

\Rightarrow maximal compatible classes

(The complete 0.5-cover)

$(1,2),(3,4,5),(4,5,6),(4,5,7),(5,8),(9)$

3.7. Ordering Relations

partial ordering: reflexive , antisymmetric , transitive

$X \leq Y$: X : predecessor

precedes Y : successor

if exist First member : if $x \leq y \forall y \in X$ (minimum)

unique Last member : if $y \leq x \forall y \in X$ (maximum)

may not Minimal member : if $y \leq x \Rightarrow x = y$

be unique Maximal member : if $x \leq y \Rightarrow x = y$

properties :

1, if \exists , at most one first member

if \exists , at most one last member

2, There may be several maximal and minimal member

3, if \exists a first member X , \Rightarrow only one minimal member Y exists and $x=y$

4, if \exists a last member x , \Rightarrow only one maximal member Y exists and $x=y$.

5, partial the first member \Leftrightarrow the last member inverse

ordering the last member \Leftrightarrow the first member partial ordering

✱ In a partial ordering, it does not guarantee that $\forall(x,y), \exists(x \leq y, \text{ or } y \leq x)$.

If $\exists, \Rightarrow (x,y)$: comparable (total ordering)

Otherwise (x,y) : non comparable

· $A \subseteq X$

If $x \in X$, and $\forall y \in A, x \leq y$,

$\Rightarrow x$: lower bound of A on X

If , $x \geq y$

$\Rightarrow x$: upper bound of A on X

· greatest lower bound (or infimum) GLB

- a lower bound which succeeds every other lower bound

Least upper bound (or supremum) LUB

- a upper bound which preceeds every other upper bound

· Lattice – A partial ordering on X contains GLB and LUB, $\forall S = \{x_1, x_2\} \subseteq X$

· Connected – a partial ordering is said to be connected

iff $\forall x, y \in X, x \neq y \Rightarrow x < y \text{ or } y < x$

· Linear ordering (total ordering, simple ordering, complete ordering)

- when a partial ordering is connected, then $\forall (x, y) : \text{comparable}$

· Hasse diagrams – representing partial orderings in which \rightarrow indicates \leq

· Example 3.12 : Crisp partial orderings

· Fuzzy partial ordering

- reflexive , antisymmetric , and transitive under some form of transitivity.

※ any fuzzy partial ordering can be resolved into a series of crisp partial ordering .

i.e. taking a series of α cut that produce increasing levels of refinement

· In a fuzzy partial ordering , R

$\forall x \in X$, two fuzzy sets are associated with

$R_{\geq[x]}$: dominating class

$$R_{\geq[x]}(y) = R(x, y)$$

$R_{\leq[x]}$: dominated class

$$R_{\leq[x]}(y) = R(y, x)$$

· xundominatediff $R(x,y) = 0 \quad \forall y \neq x$

X undominatingiff $R(y,x) = 0 \quad \forall y \neq x$

· Fuzzy upper bound for $A \subseteq X$ is a fuzzy set

$$U(R, A) = \bigcap_{x \in A} R_{\geq [x]}$$

※ If a least upper bound of A exists , it is the unique element $x \in U(R, A)$

s.t. 1, 2,

$$U(R, A)(x) > 0 \quad R(x,y) > 0 ,$$

$$\forall y \in \text{support} [U(R,A)]$$

· Example 3.13

a b c d e

Fuzzy partial ordering R:

$$\begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} \begin{bmatrix} 1 & 0.7 & 0 & 1 & 0.7 \\ 0 & 1 & 0 & 0.9 & 0 \\ 0.5 & 0.7 & 1 & 1 & 0.8 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0.1 & 0 & 0.9 & 1 \end{bmatrix}$$

1. row : dominating class for each element
column : dominated class for each element

2. d : undominated , C : undominating

3. For $A = \{a,b\}$, $U(R,A)$ = the intersection of

$$\text{The dominating classes of a and b} = \frac{0.7}{b} + \frac{0.9}{d}$$

- 4, LUB(A) =b

3-32

3. Crisp ordering captured by the fuzzy ordering

e.g. $\alpha = 0.5$

$${}^{\alpha}R \Rightarrow \begin{bmatrix} 1 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$

is $\rightarrow 2\ 3\ 1\ 5\ 3$

✕The ordering become weaken with the increasing α

5-33

Fuzzy preordering – reflexive and transitive

Fuzzy weak ordering –

i, an ordering satisfying the proportion of a fuzzy total ordering except antisymmetry.

ii, a fuzzy preordering in which $\forall x \neq y$, either $R(x,y) > 0$ or $R(y,x) > 0$

Fuzzy strict ordering –

Antireflexive

Antisymmetric

Transitive

3.8. Morphisms

• Crisp homomorphism h from (X,R) to (Y,Q)

Where $R(X,X), Q(Y,Y)$: binary relations

$$\forall (x_1, x_2) \in R \Rightarrow (h(x_1), h(x_2)) \in Q$$

• Fuzzy homomorphism h

If $R(X,X), Q(Y,Y)$: Fuzzy binary relations

$$\text{And } R(x_1, x_2) \leq Q[h(x_1), h(x_2)]$$

※ It's possible that a relation $(h(x_1), h(x_2)) \in Q$ which $(x_1, x_2) \notin R$.

※ If this is never the case h is called a strong homomorphism.

3-34

• Crisp strong homomorphism h

If $(x_1, x_2) \in R \Rightarrow (h(x_1), h(x_2)) \in Q$

And $(y_1, y_2) \in Q \Rightarrow (h^{-1}(y_1), h^{-1}(y_2)) \in R$

※ where h : many to one $\rightarrow h^{-1}(y)$ contains a set of X s

• Fuzzy strong homomorphism h

H imposes a partition Π_h on X

Let

$$A = \{a_1, a_2, \dots, a_n\}$$

$$B = \{b_1, b_2, \dots, b_n\} \in \Pi_h$$

R, Q : fuzzy relations

h : strong homomorphism

iff $\max_{i,j} (R(a_i, b_j)) = Q(y_1, y_2)$

where $\begin{cases} y_1 = h(a_i) \forall a_i \in A \\ y_2 = h(b_j) \forall b_j \in B \end{cases}$

3-33

• Example 3.14

$R(X,X)$

$$R = \begin{bmatrix} 0 & 0.5 & 0 & 0 \\ 0 & 0 & 0.9 & 0 \\ 1 & 0 & 0 & 0.5 \\ 0 & 0.6 & 0 & 0 \end{bmatrix}$$

$Q(Y,Y)$

$$Q = \begin{bmatrix} 0.5 & 0.9 & 0 \\ 1 & 0 & 0.9 \\ 1 & 0.9 & 0 \end{bmatrix}$$

$\rightarrow h$: ordinary fuzzy homomorphism (one way)

$\therefore \forall R(x_1, x_2) \leq Q(h(x_1), h(x_2)) \leftarrow \text{strong}$

But $Q(\gamma, \beta) = 0.9$ $R(d, c) = 0$

i.e, $(\gamma, \beta) \in Q$ while $(d, c) \notin R$ where $h(d) = \gamma, h(c) = \beta$

3-36

$R(X,X)$

$$\begin{bmatrix} 0.8 & 0.4 & 0 & 0 & 0 & 0 \\ 0 & 0.5 & 0 & 0.7 & 0 & 0 \\ 0 & 0 & 0.3 & 0 & 0 & 0 \\ 0 & 0.5 & 0 & 0 & 0.9 & 0.5 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0.8 \end{bmatrix}$$

$Q(Y,Y)$

$$\begin{bmatrix} 0.7 & 0 & 0.9 \\ 0.4 & 0.8 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

$\rightarrow h$: strong fuzzy homomorphism (two way)

3-37

※Q represents a simplification of R

• Isomorphism : (congruence)

$h:1-1, \text{ onto} \Rightarrow X \neq Y$

Endomorphism : (subgraph)

$h:X \rightarrow Y, Y \subseteq X$

Automorphism :

Isomorphism and End Endomorphism

i.e. $m X=Y$ nad $R=Q$

3.9 SUP-i Compositions of Fuzzy Relations Generalize max-min Composition

i : t-norm

sup : t-conorm

- $P(X,Y), Q(Y,Z)$: fuzzy relations

 $P \overset{i}{o} Q(X \times Z)$: sup-i composition

$$[P \overset{i}{o} Q](X, Z) = \sup_{y \in Y} i[P(x, y), Q(y, z)]$$

- Properties

1. $(P \overset{i}{o} Q) \overset{i}{o} R = P \overset{i}{o} (Q \overset{i}{o} R)$
2. $P \overset{i}{o} (\cup_j Q_j) = \cup_j (P \overset{i}{o} Q_j)$
3. $P \overset{i}{o} (\cap_j Q_j) = \cap_j (P \overset{i}{o} Q_j)$
4. $(\cup_j P_j) \overset{i}{o} Q = \cup_j (P_j \overset{i}{o} Q)$
5. $(\cap_j P_j) \overset{i}{o} Q = \cap_j (P_j \overset{i}{o} Q)$
6. $(P \overset{i}{o} Q)^{-1} = Q^{-1} \overset{i}{o} P^{-1}$

3. Show Eq.(3.16), i.e., $P \circ (\bigcap_{j \in J} Q_j) \subseteq \bigcap_{j \in J} (P \circ Q_j)$,

Where $P(\bar{X}, \bar{Y})$ and $Q(\bar{Y}, \bar{Z})$ are fuzzy relations.

pf. From Eq.(3.13), i.e., $\left[P \circ Q \right](x, z) = \sup_{y \in \bar{Y}} i[P(x, y), Q(y, z)]$

$$\therefore \left[P \circ \left(\bigcap_{j \in J} Q_j \right) \right](x, z) = \sup_{y \in \bar{Y}} i \left[P(x, y), \bigcap_{j \in J} Q_j(y, z) \right]$$

$$\text{Let } Q = \bigcap_{j \in J} Q_j$$

$$\Rightarrow Q \subseteq Q_1, Q \subseteq Q_2, \dots, Q \subseteq Q_{|J|}$$

$$\text{i.e., } \forall(y, z), Q(y, z) \leq Q_1(y, z), \dots, Q(y, z) \leq Q_{|J|}(y, z)$$

$\therefore i$ is monotonically increasing

$$\therefore \begin{cases} i[P(x, y), \bigcap_{j \in J} Q_j(y, z)] \leq i[P(x, y), Q_1(y, z)] \\ \dots\dots\dots \\ i[P(x, y), \bigcap_{j \in J} Q_j(y, z)] \leq i[P(x, y), Q_{|J|}(y, z)] \end{cases} \quad \forall(x, y)$$

$$i[P(x, y), \left(\bigcap_{j \in J} Q_j \right)(y, z)] \leq \bigcap_{j \in J} i[P(x, y), Q_j(y, z)]$$

$$\therefore \sup_{y \in \bar{Y}} i[P(x, y), \left(\bigcap_{j \in J} Q_j \right)(y, z)] \leq \sup_{y \in \bar{Y}} \bigcap_{j \in J} i[P(x, y), Q_j(y, z)]$$

$$= \bigcap_{j \in J} \sup_{y \in \bar{Y}} i[P(x, y), Q_j(y, z)], \forall(x, y), (y, z)$$

$$\Rightarrow \left[P \circ \left(\bigcap_{j \in J} Q_j \right) \right](x, z) \leq \left[\bigcap_{j \in J} (P \circ Q_j) \right](x, z), \forall(x, z)$$

$$\text{i.e., } P \circ \left(\bigcap_{j \in J} Q_j \right) \subseteq \bigcap_{j \in J} (P \circ Q_j)$$

◦ Sup- i composition monotonically increases

$$\text{i.e., } P \circ^i Q_1 \subseteq P \circ^i Q_2 \text{ -----(5.20) if } Q_1 \subseteq Q_2$$

$$Q_1 \circ^i P \subseteq Q_2 \circ^i P \text{ -----(5.21)}$$

◦ Identity of \circ^i

$$E(x, y) = \begin{cases} 1, & x = y \\ 0, & x \neq y \end{cases} \begin{pmatrix} 1 & & 0 \\ & \ddots & \\ 0 & & 1 \end{pmatrix}$$

$$\text{i.e., } E \circ^i P = P \circ^i E = P$$

◦ Relation R on \bar{X} : i -transitive

$$\text{iff } R(x, z) \geq i[R(x, y), R(y, z)], \forall x, y, z \in \bar{X}$$

$$\Leftrightarrow R \circ^i R \subseteq R$$

◦ i -transitive closure $R_{T(i)}$

--- The smallest i -transitive relation containing R

◦ Theorem 3.1: R : any fuzzy relation

$$\Rightarrow R_{T(i)} = \bigcup_{n=1}^{\infty} R^{(n)}, \text{ where}$$

$$R^{(n)} = R \circ^i R^{(n-1)}$$

3-41

proof:

By (5.15) (5.17)

$$\begin{aligned} i, R_{T(i)} \circ^i R_{T(i)} &= \left(\bigcup_{n=1}^{\infty} R^{(n)} \right) \circ^i \left(\bigcup_{m=1}^{\infty} R^{(m)} \right) \stackrel{\downarrow}{=} \bigcup_{n=1}^{\infty} \bigcup_{m=1}^{\infty} \left(R^{(n)} \circ^i R^{(m)} \right) = \bigcup_{n,m=1}^{\infty} R^{(n+m)} \\ &= \bigcup_{k=2}^{\infty} R^{(k)} \subseteq \bigcup_{k=1}^{\infty} R^{(k)} = R_{T(i)} \end{aligned}$$

i.e., $R_{T(i)} : i$ -transitive ($\because R_{T(i)} \circ^i R_{T(i)} \subseteq R_{T(i)}$)

ii, Let $S : i$ -transitive, $R \subseteq S$ (5.20)(5.21) monotonically increasing

$$\begin{aligned} \Rightarrow R^{(2)} &= R \circ^i R \subseteq S \circ^i S \subseteq S \\ \text{If } R^{(n)} &\subseteq S, \quad \left. \begin{array}{l} \text{i-transitive} \\ \text{mathematical} \\ \text{induction} \end{array} \right\} \\ \Rightarrow R^{(n+1)} &= R \circ^i R^{(n)} \subseteq S \circ^i S \subseteq S \\ \therefore R^{(k)} &\subseteq S, \forall k \\ \therefore R_{T(i)} &= \bigcup_{k=1}^{\infty} R^{(k)} \subseteq S \end{aligned}$$

i.e., $R_{T(i)} : \text{smallest}$

◦ Theorem 3.2: R : reflexive fuzzy relation on $\bar{X}^2, \left| \bar{X} \right| = n$

$$\Rightarrow R_{T(i)} = R^{(n-1)} \Leftrightarrow \left(\begin{array}{l} R^{(m)} \subseteq R^{m+1} \\ R^n = R^{n-1} \end{array} \right) \forall m$$

3-42

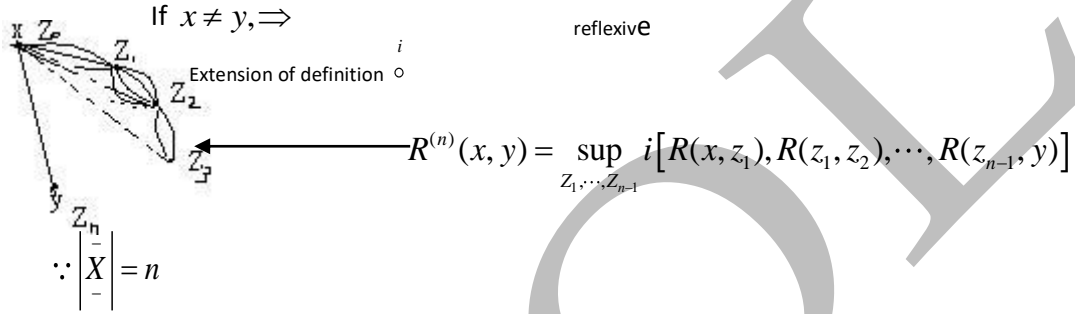
proof :i, $\because R : reflexive$,

$$\therefore E \subseteq R, R = E \circ R \subseteq R \circ R = R^{(2)}$$

$$\therefore R^{(n-1)} \subseteq R^{(n)} \quad (\text{By repetition})$$

$$\text{ii, show } R^{(n-1)} \supseteq R^{(n)}$$

proof: If $x = y, \Rightarrow R^{(n-1)}(x, x) = 1$



$\Rightarrow X = Z_0, Z_1, \dots, Z_n = y$ contains

at least 2 identical element.

$$\text{Say } Z_r = Z_s (r < s)$$

$$\Rightarrow i[R(x, z_1), \dots, R(z_{r-1}, z_r), \dots, R(z_s, z_{s+1}), \dots, R(z_{n-1}, y)] \leq R^{(k)}(x, y), (k \leq n-1)$$

$$\therefore \forall x, y \in X, R^{(n)}(x, y) \leq R^{(n-1)}(x, y),$$

$$\therefore R^{(n)} \subseteq R^{(n-1)} \text{ --- } (B) \therefore R^{(n)} = R^{(n-1)} \leftarrow (A, B)$$

$$\therefore R_{T(i)} = R^{(n-1)}$$

3.10 INF- W_i Compositions of Fuzzy Relations

◦ W_i operation:

$$\begin{aligned} a > b &\rightarrow b \\ a \leq b &\rightarrow 1 \end{aligned}$$

$$w_i(a, b) = \sup \{x \in [0, 1] \mid i(a, x) \leq b\} \rightarrow$$

where $a, b \in [0, 1]$, i : continuous t-norm

※ If i : logical conjunction (i.e., \wedge , and)

$\Rightarrow w_i$: logical implication (i.e., \Rightarrow , if then)

a	b	$a \Rightarrow b$
0	0	1
0	1	1
1	0	0
1	1	1

◦ Theorem 3.3

$$1, i(a, b) \leq d \text{ iff } w_i(a, b) \geq b$$

$$2, w_i(w_i(a, b), b) \geq a \quad \begin{array}{|c|c|} \hline a & b \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline b & a \\ \hline \end{array}$$

$$3, w_i(i(a, b), d) = w_i(a, w_i(b, d)) \quad \begin{array}{|c|c|} \hline a & b \\ \hline \end{array} \quad \begin{array}{|c|c|} \hline b & a \\ \hline \end{array}$$

$$4, a \leq b \Rightarrow w_i(a, d) \geq w_i(b, d) \text{ --- i}$$

$$w_i(d, a) \leq w_i(d, b) \text{ --- ii}$$

$$3, i(w_i(a, b), w_i(b, d)) \leq w_i(a, d)$$

$$6, w_i(\inf_j a_j, b) \geq \sup_j w_i(a_j, b)$$

$$7, w_i(\sup_j a_j, b) = \inf_j w_i(a_j, b)$$

$$8, w_i(b, \sup_j a_j) \geq \sup_j w_i(b, a_j)$$

$$9, w_i(b, \inf_j a_j) = \inf_j w_i(b, a_j)$$

$$10, i(a, w_i(a, b)) \leq b$$

proof: (1) i, If $i(a, b) \leq d, \Rightarrow b \in \{x \mid i(a, x) \leq d\}$

$$(\Rightarrow) \Rightarrow b \leq \sup\{x \mid i(a, x) \leq d\} = w_i(a, d)$$

ii, If $b \leq w_i(a, d)$

$$(3) \quad (\Leftarrow) \Rightarrow i(a, b) \leq i(a, w_i(a, d)) = i(a, \sup\{x \mid i(a, x) \leq d\}) \stackrel{\text{i: continuous monotone}}{=} \sup\{i(a, x) \mid i(a, x) \leq d\} \leq d$$

$$\because i(a, x) \leq w_i(b, d) \Leftrightarrow i(b, i(a, x)) \leq d$$

Associativity
commutation

$$\Leftrightarrow i(i(a, b), x) \leq d \Leftrightarrow x \leq w_i(i(a, b), d)$$

$$w_i(i(a, b), d)$$

$$\therefore w_i(a, w_i(b, d)) = \sup\{x \mid i(a, x) \leq w_i(b, d)\} \stackrel{\text{By (A)}}{=} \sup\{x \mid x \leq w_i(i(a, b), d)\} = w_i(i(a, b), d)$$

$$(7) \text{ Let } S = \sup_{j_0} a_j \text{ --- (B)} \Rightarrow a_j \leq s, \forall j$$

$$\Rightarrow w_i(s, b) \leq w_i(a_j, b), \forall j$$

$$\therefore w_i(s, b) \leq \inf_j w_i(a_j, b) \text{ --- (C)}$$

$$\because \inf_j w_i(a_j, b) \leq w_i(a_{j_0}, b), \forall j_0 \in J$$

$$\Rightarrow i(a_{j_0}, \inf_j w_i(a_j, b)) \leq b, \forall j_0$$

$$\therefore i(s, \inf_j w_i(a_j, b)) = \sup_{j_0} i(a_{j_0}, \inf_j w_i(a_j, b)) \leq b$$

$$\Rightarrow w_i(s, b) \geq \inf_j w_i(a_j, b) \text{ --- (D)}$$

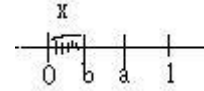
$$\Rightarrow w_i(\sup_j a_j, b) = w_i(s, b) = \inf_j w_i(a_j, b)$$

(2) Show $w_i(w_i(a,b),b) \geq a$ (Theorem 3.3 (2))

proof $\because w_i(a,b) = \text{Sup}\{x \in [0,1] \mid i(a,x) \leq b\}$ and by Theorem 3.10

$$i_{\min}(a,b) \leq i(a,b) \leq \min(a,b)$$

i, If $a > b$



$$w_i(a,b) = \text{Sup}\{x \in [0,1] \mid i(a,x) \leq b\} \leq \text{Sup}\{x \in [0,1] \mid \min(a,x) \leq b\} = b$$

$$i(w_i(a,b),a)$$

$$\leq i(b,a)$$

$$\leq i(b,1)$$

$$= b$$

$$\Rightarrow i(w_i(a,b),a) \leq b$$

i, If $a \leq b$

$$w_i(a,b) = \text{Sup}\{x \in [0,1] \mid i(a,x) \leq b\} \leq \text{Sup}\{x \in [0,1] \mid \min(a,x) \leq b\} = 1$$

$$\therefore i(w_i(a,b),a) \leq i(w_i(a,b),b)$$

$$\leq i(1,a)$$

$$\leq i(b,1)$$

$$= b$$

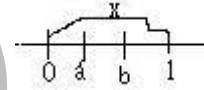
$$\Rightarrow i(w_i(a,b),a) \leq b$$

By Theorem 3.3 property 1 (i.e., $i(a,b) \leq d$ iff $w_i(a,d) \geq b$)

$$i(w_i(a,b),a) \leq b \Rightarrow w_i(w_i(a,b),b) \geq a$$

By Axiom i2 $w_i(a,b) \leq b$

By Axiom i2 ($a \leq 1$)



(4) prove Theorem 3.3 (4) : $a \leq b \Rightarrow i, W_i(a, d) \geq W_i(b, d)$

ii, $W_i(d, a) \leq W_i(d, b)$

proof :i, $W_i(a, d) \geq W_i(b, d)$

$$W_i(a, d) = \sup\{x \mid i(a, x) \leq d\} \text{ ---- (A)}$$

$$W_i(b, d) = \sup\{x \mid i(b, x) \leq d\} \text{ ---- (B)}$$

a, if $d \leq a \leq b \Rightarrow (A)=d, (B)=d,$

$$\therefore (A)=(B) \text{ ---- (1)}$$

b, if $a \leq d \leq b \Rightarrow (A)=1, (B)=d$

$$\therefore (A) \geq (B) \text{ ---- (2)}$$

c, if $a \leq b \leq d \Rightarrow (A)=1, (B)=1$

$$\therefore (A)=(B) \text{ ---- (3)}$$

$$(1),(2),(3) \Rightarrow (A) \geq (B)$$

$$\text{i.e., } W_i(a, d) \geq W_i(b, d)$$

ii, see i

5. show $i(W_i(a, d), W_i(b, d)) \leq W_i(a, d)$

Proof : \because if $a > b \Rightarrow W_i(a, b) = b$

if $a \leq b \Rightarrow W_i(a, b) = 1$

$$\text{A, if } a \leq b \Rightarrow i(W_i(a, b), W_i(b, d)) = i(b, W_i(b, d))$$

$$\Rightarrow i(b, W_i(b, d)) = i(b, d) \leq \min(b, d) = d$$

$$W_i(a, d) = d$$

$$\Rightarrow i(b, W_i(b, d)) = i(b, 1) = b$$

$$W_i(a, d) = 1 \quad \left\{ \right.$$

$$\Rightarrow i(b, W_i(b, d)) = i(b, 1) = b$$

$$W_i(a, d) = 1$$

$$\text{B, if } a \leq b \Rightarrow i(W_i(a, b), W_i(b, d)) = i(1, W_i(b, d)) = W_i(b, d)$$

$$\Rightarrow W_i(b, d) = d$$

$$W_i(a, d) = d$$

$$W_i(b, d) = d$$

$$W_i(a, d) = 1$$

$$W_i(b, d) = 1$$

$$W_i(a, d) = 1$$

$$10. \text{ show } i(a, W_i(a, b)) \leq b$$

$$\text{Proof } \because a > b \Rightarrow W_i(a, b) = b$$

$$a \leq b \Rightarrow W_i(a, b) = 1$$

$$\text{A, if } a > b$$

$$\Rightarrow i(a, W_i(a, b)) = i(a, b) \leq \min(a, b) = b$$

$$\text{B, if } a \leq b$$

$$\Rightarrow i(a, W_i(a, b)) = i(a, 1) = a \leq b$$

- $\inf - W_i$ composition

$$(P \circ Q)(x, z) = \inf_{y \in Y} W_i(P(x, y), Q(y, z))$$

- Theorem 3.4 :

$$(1) (P \circ Q \subseteq R) \Leftrightarrow (Q \subseteq P^{-1} \circ R) \Leftrightarrow (P \subseteq (Q \circ R^{-1})^{-1})$$

$$(2) (P \circ (Q \circ S) = (P \circ Q) \circ S$$

- Theorem 3.5 :

$$(\bigcup_j P_j) \circ Q = \bigcap_j (P_j \circ Q)$$

$$(\bigcap_j P_j) \circ Q \supseteq \bigcup_j (P_j \circ Q)$$

$$P \circ (\bigcap_j Q_j) = \bigcap_j (P \circ Q_j)$$

$$P \circ (\bigcup_j Q_j) \supseteq \bigcup_j (P \circ Q_j)$$

- Theorem 3.6 : if $Q_1 \subseteq Q_2 \Rightarrow P \circ Q_1 \subseteq P \circ Q_2$

$$Q_1 \circ R \supseteq Q_2 \circ R$$

$$\text{Proof: } Q_1 \subseteq Q_2 \Rightarrow Q_1 \cap Q_2 = Q_1, Q_1 \cup Q_2 = Q_2$$

$$\therefore (P \circ Q_1) \cap (P \circ Q_2) = P \circ (Q_1 \cap Q_2) = P \circ Q_1$$

$$\Rightarrow P \circ Q_1 \subseteq P \circ Q_2$$

$$\therefore (Q_1 \circ R) \cap (Q_2 \circ R) = (Q_1 \cup Q_2) \circ R = Q_2 \circ R$$

$$\Rightarrow Q_1 \circ R \supseteq Q_2 \circ R$$

● Theorem 3.7 :

$$1. P^{-1} \circ (P \circ Q) \subseteq Q$$

$$2. R \subseteq P \circ (P^{-1} \circ R)$$

$$3. P \subseteq (P \circ Q) \circ Q^{-1}$$

$$4. R \subseteq (R \circ Q^{-1}) \circ Q$$

Proof :

$$(1) \because P \circ Q \subseteq (P^{-1})^{-1} \circ Q \text{ ---- (A)}$$

$$\because (5.26) \Longleftrightarrow (5.25)$$

$$\text{i.e., } (\vec{Q} \subseteq \vec{P}^{-1} \circ \vec{R}) \Longleftrightarrow (\vec{P} \circ \vec{Q} \subseteq \vec{R})$$

$$\text{let } P \circ Q = \vec{Q}, P^{-1} = \vec{P}, Q = \vec{R}$$

$$\therefore (A) \Longleftrightarrow P^{-1} (P \circ Q) \subseteq Q$$

$$(2) \because P^{-1} \circ R \subseteq P^{-1} \circ R \text{ ---- (B)}$$

$$\text{Let } P^{-1} = \vec{P}, Q = \vec{R}, P^{-1} \circ R = \vec{R}$$

$$\therefore (B) \Longleftrightarrow R \subseteq P \circ (P^{-1} \circ R)$$

$$(3) \text{ by (3.33), } [P^{-1} \circ (P \circ Q)]^{-1} \subseteq Q^{-1}$$

$$\Rightarrow (P \circ Q)^{-1} \circ P \subseteq Q^{-1} \text{ ---- (C)}$$

$$\text{Let } P \circ Q = \vec{P}, P = \vec{Q}, Q^{-1} = \vec{R}$$

$$\therefore (C) \Longleftrightarrow P \subseteq (P \circ Q)^{-1} \circ Q^{-1}$$

(4) follows (3)

DRAFT

Chapter 4

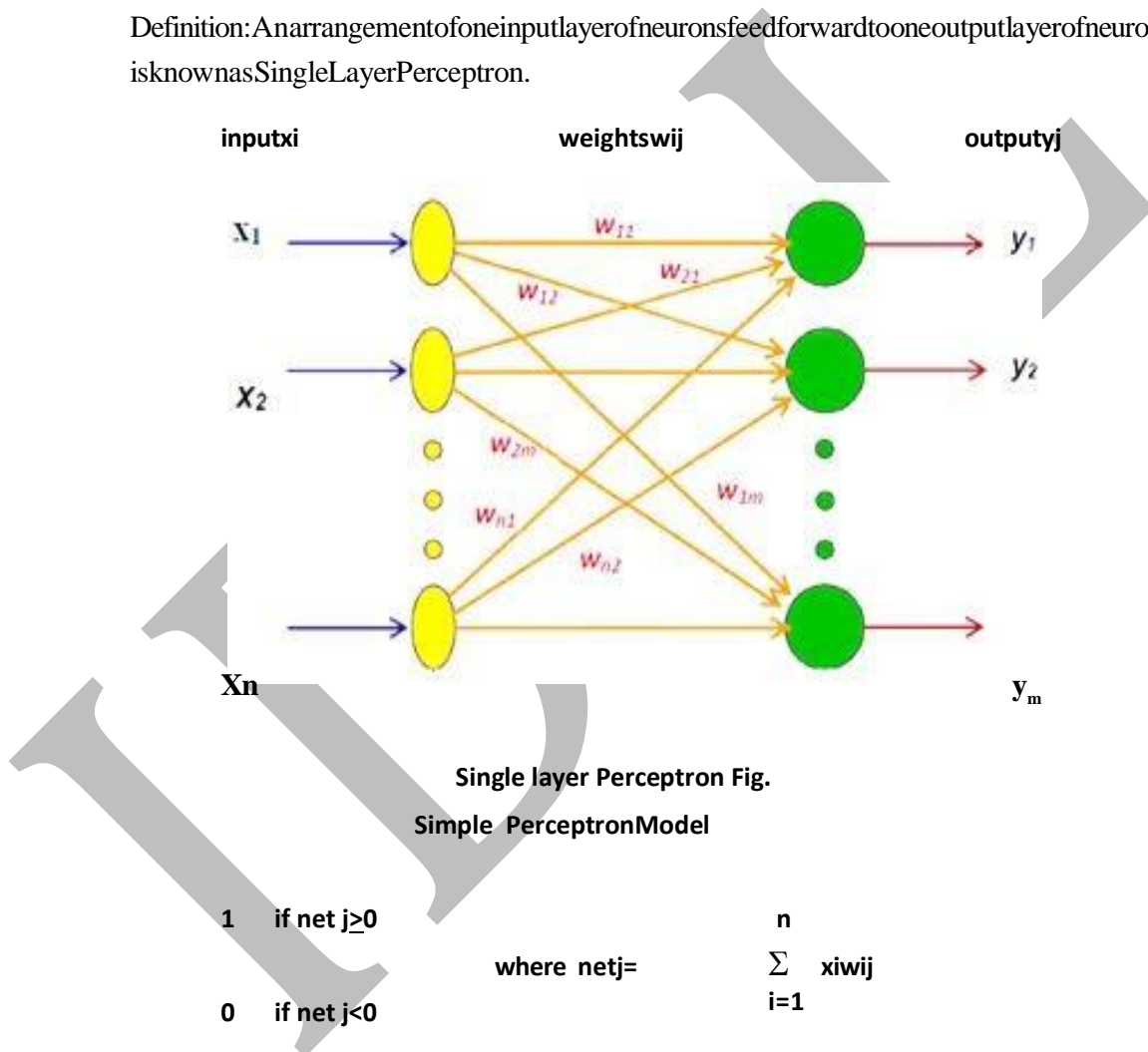
Single Layer Perceptron: Perceptron convergence theorem, Method of steepest descent-least mean square algorithms.

Single-Layer NN Systems

Here, a simple Perceptron Model and an ADALINE Network Model is presented.

Single layer Perceptron

Definition: An arrangement of one input layer of neurons feed forward to one output layer of neurons is known as Single Layer Perceptron.



- Learning Algorithm: Training Perceptron

The training of Perceptron is a supervised learning algorithm where weights are adjusted to minimize error whenever the output does not match the desired output.

- If the output is correct then no adjustment of weights is done.

i.e.

$$w_{ij}^{K+1} = w_{ij}^K$$

- If the output is **1** but should have been **0** then the weights are decreased on the active input link

i.e.
$$w_{ij}^{k+1} = w_{ij}^k - \alpha_i$$

- If the output is 0 but should have been 1 then the weights are increased on the active input link

i.e.
$$W_{ij}^{K+1} = W_{ij}^K + \alpha \cdot x_i$$

Where

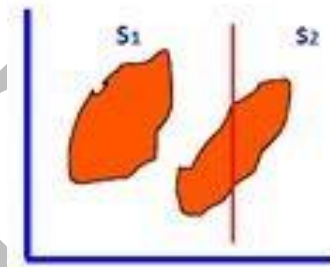
W_{ij}^{K+1} is the new adjusted weight, W_{ij}^K is the old weight
 x_i is the input and α is the learning rate parameter.
 α small leads to slow and α large leads to fast learning.

- Perceptron and Linearly Separable Task

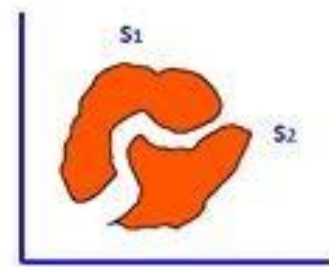
Perceptron cannot handle tasks which are not separable.

- Definition: Set of points in 2-D space are linearly separable if the sets can be separated by a straight line.
- Generalizing, a set of points in n-dimensional space are linearly separable if there is a hyperplane of (n-1) dimensions that separates the sets.

Example



(a) Linearly separable patterns



(b) Not Linearly separable patterns

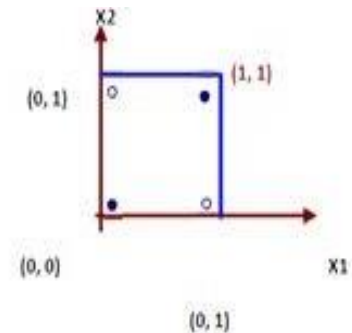
Note: Perceptron cannot find weights for classification problems that are not linearly separable.

- XOR Problem

:Exclusive

Input x1	Input x2	Output
0	0	0
1	1	0
0	1	1
1	0	1

XOR truth table



OR operation

Fig. Output of XOR in x_1, x_2 plane

Even parity is, even number of 1 bits in the input

Odd parity is, odd number of 1 bits in the input

- There is no way to draw a single straight line so that the circles are on one side of the line and the dots on the other side.
- Perceptron is unable to find a line separating even parity input patterns from odd parity input patterns.

- Perceptron Learning Algorithm

The algorithm is illustrated step-by-step.

* **Step 1:**

Create a perceptron with $(n+1)$ input neurons x_0, x_1, \dots, x_n , where $x_0 = 1$ is the bias input. Let O be the output neuron.

* **Step 2:**

Initialize weight $W = (w_0, w_1, \dots, w_n)$ to random weights.

* **Step 3:**

Iterate through the input patterns X_j of the training set using the weight set; i.e. compute the weighted sum of inputs $net_j = \sum_{i=1}^n x_i w_i$ for each input pattern j .

* **Step 4:**

Compute the output y_j using the step function

$$y_j = f(net_j) = \begin{cases} 1 & \text{if } net_j \geq 0 \\ 0 & \text{if } net_j < 0 \end{cases} \quad \text{where } net_j = \sum_{i=1}^n x_i w_i$$

* **Step 5:**

Compare the computed output y_j with the target output y_j
for each input pattern j . If all the input patterns have been classified correctly, then output (read) the weights and exit.

* **Step 6:**

Otherwise, update the weights as given below :

If the computed output y_j is 1 but should have been 0,

Then $w_i = w_i - \alpha x_i$, $i = 0, 1, 2, \dots, n$

If the computed output y_j is 0 but should have been 1,

Then $w_i = w_i + \alpha x_i$, $i = 0, 1, 2, \dots, n$

where α is the learning parameter and is constant.

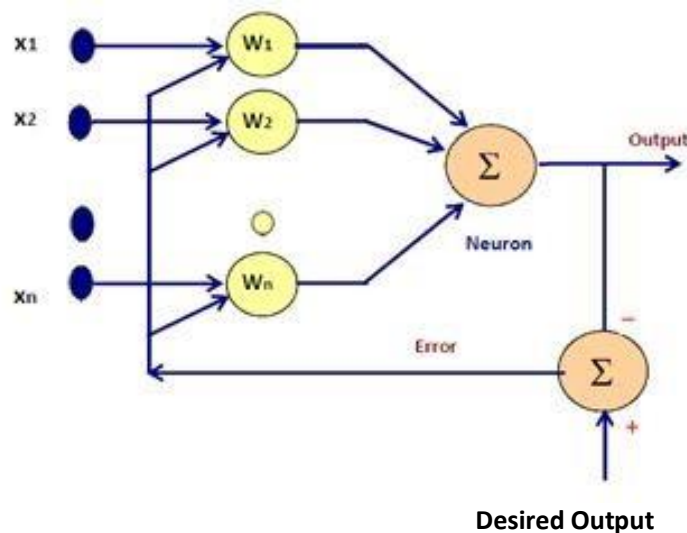
* **Step 7:**

goto step 3

* **END**

An ADALINE consists of a single neuron of the McCulloch-Pitts type, where its weights are determined by the normalized least mean square (LMS) training law. The LMS learning rule is also referred to as the delta rule. It is a well-established supervised training method that has been used over a wide range of diverse applications.

• **Architecture of a simple ADALINE**



The basic structure of an ADALINE is similar to a neuron with a linear activation function and a feedback loop. During the training phase of ADALINE, the input vector as well as the desired output are represented to the network.

[The complete training mechanism has been explained in the next slide.]

- **ADALINE Training Mechanism**

(Ref. Fig. in the previous slide - Architecture of a simple ADALINE)

- * The basic structure of an ADALINE is similar to a linear neuron with an extra feedback loop.
- * During the training phase of ADALINE, the input vector $X = [x_1, x_2, \dots, x_n]^T$ as well as desired output are presented to the network.
- * The weights are adaptively adjusted based on delta rule.
- * After the ADALINE is trained, an input vector presented to the network with fixed weights will result in a scalar output.
- * Thus, the network performs an n -dimensional mapping to a scalar value.
- * The activation function is not used during the training phase. Once the weights are properly adjusted, the response of the trained unit can be tested by applying various inputs, which are not in the training set. If the network produces consistent responses to a high degree with the test inputs, it is said that the network could generalize. The process of training and generalization are two important attributes of this network.

Usage of ADALINE :

In practice, an ADALINE is used to

- Make binary decisions; the output is sent through a binary threshold.
- Realizations of logic gates such as AND, NOT and OR.
- Realize only those logic functions that are linearly separable.

Applications of Neural Network

Neural Network Applications can be grouped in following categories:

- * **Clustering:**

A clustering algorithm explores the similarity between patterns and places similar patterns in a cluster. Best known applications include data compression and data mining.

- * **Classification/Pattern recognition:**

The task of pattern recognition is to assign an input pattern (like handwritten symbol) to one of many classes. This category includes algorithmic implementations such as associative memory.

- * **Function approximation:**

The task of function approximation is to find an estimate of the unknown function subject to noise. Various engineering and scientific disciplines require function approximation.

- * **Prediction Systems:**

The task is to forecast some future values of a time-sequenced data. Prediction has a significant impact on decision support systems. Prediction differs from function approximation by considering time factor. System may be dynamic and may produce different results for the same input data based on system state (time).

Multilayer Perceptron: Derivation of the back-propagation, Algorithm , Learning Factors

MULTI-LAYER PERCEPTRONS

In the previous section we showed that by adding an extra hidden unit, the XOR problem can be solved. For binary units, one can prove that this architecture is able to perform many transformations given the correct connections and weights. The most primitive is the next one. For a given transformation $y = d(x)$, we can divide the set of all possible input vectors into two classes: $X^+ = \{x | d(x) = 1\}$ and $X^- = \{x | d(x) = 0\}$.

Since there are N input units, the total number of possible input vectors x is 2^N . For every $x^p \in X^+$ a hidden unit can be reserved of which the activation y_h is 1 if and only if the specific pattern p is present at the input: we can choose its weights w_{ih} equal to the specific pattern x_i^p and the bias U_h equal to $1 - N$ such that

$$y_h^p = \text{sgn} \left[\sum_i w_{ih} x_i^p - N + 1 \right]$$

is equal to 1 for $x^p = w$ only. Similarly, the weights to the output neuron can be chosen such that the output is one as soon as one of the M predicate neurons is one:

$$y_o^p = \text{sgn} \left[\sum_{h=1}^M w_{oh} y_h^p - M + 1 \right]$$

This perceptron will give $y_o = 1$ only if $x \in X^+$; it performs the desired mapping. The problem is the large number of predicate units, which is equal to the number of patterns in X^+ , which is maximally 2^N . Of course we can do the same trick for X^- , and we will always take the minimal number of mask units, which is maximally 2^{N-1} . A more elegant proof is given by Minsky and Papert, but the point is that for complex transformations the number of required units in the hidden layer is exponential in N .

Back-Propagation

As we have seen in the previous chapter, a single-layer network has severe restrictions: the class of tasks that can be accomplished is very limited. In this chapter we will focus on feed-forward networks with layers of processing units.

Minsky and Papert showed in 1969 that a two-layer feed-forward network can overcome many restrictions, but did not present a solution to the problem of how to adjust the weights from input to hidden units. An answer to this question was presented by Rumelhart, Hinton and Williams in 1986, and similar solutions appeared to have been published earlier (Parker, 1985; Cun, 1985).

The central idea behind this solution is that the errors for the units of the hidden layer are determined by back-propagating the errors of the units of the output layer. For this reason the method is often called the back-propagation learning rule. Back-propagation can also be considered as a generalization of the delta rule for non-linear activation functions and multilayer networks.

IDOL

MULTI - LAYER FEED - FORWARD NETWORKS

A feed-forward network has a layered structure. Each layer consists of units, which receive their input from units from a layer directly below and send their output to units in a layer directly above the unit. There are no connections within a layer.

The N_i inputs are fed into the first layer of $N_{h,1}$ hidden units. The input units are merely 'fan-out' units; no processing takes place in these units. The activation of a hidden unit is a function F of the weighted inputs plus a bias, as given in eq. (10.4). The output of the hidden units is distributed over the next layer of $N_{h,2}$ hidden units, until the last layer of hidden units, of which the outputs are fed into a layer of N_o output units (see Fig. 12.1).

Although back-propagation can be applied to networks with any number of layers, just as for networks with binary units (section 11.7) it has been shown (Cybenko, 1989; Funahashi, 1989; Hornik, Stinchcombe, & White, 1989; Hartman, Keeler, & Kowalski, 1990) that only one layer of hidden units suffices to approximate any function with finitely many discontinuities to arbitrary precision, provided the activation functions of the hidden units are non-linear (the universal approximation theorem).

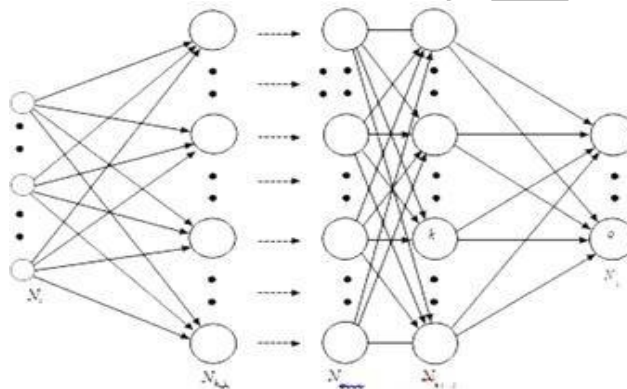


Fig. A multi-layer network with layers of units.

In most applications a feed-forward network with a single layer of hidden units is used with a sigmoid activation function for the units.

THE GENERALISED DELTA RULE

Since we are now using units with non-linear activation functions, we have to generalise the delta rule, which was presented in chapter 11 for linear functions to these sets of non-linear activation functions. The activation is a differentiable function of the total input, given by

$$y_k^p = F(S_k^p) \quad \dots\dots\dots (12.1)$$

in which

$$S_k^p = \sum_j w_{jk} y_j^p + \theta_k \quad \dots\dots(12.2)$$

To get the correct generalization of the delta rule as presented in the previous chapter, we must set

$$\Delta_p w_{jk} = -\gamma \frac{\partial E^p}{\partial w_{jk}} \quad \dots\dots(12.3)$$

The error E^p is defined as the total quadratic error for pattern p at the output units :

$$E^p = \frac{1}{2} \sum_{o=1}^{N_o} (d_o^p - y_o^p)^2 \quad \dots\dots\dots(12.4)$$

where d^p is the desired output for unit k when pattern p is clamped. We further set $E = \sum_p E^p$ as the summed squared error. We can write

$$\frac{\partial E^p}{\partial w_{jk}} = \frac{\partial E^p}{\partial S_k^p} \frac{\partial S_k^p}{\partial w_{jk}} \quad \dots(12.5)$$

By equation (12.2) we see that the second factor is

$$\frac{\partial S_k^p}{\partial w_{jk}} = y_{jk}^p \quad \dots(12.6)$$

When we define

$$\delta_k^p = \frac{\partial E^p}{\partial S_k^p} \quad \dots(12.7)$$

we will get an update rule which is equivalent to the delta rule as described in the previous chapter, resulting in a gradient descent on the error surface if we make the weight changes according to:

$$\Delta_p w_{jk} = \gamma \delta_k^p y_{jk}^p \quad \dots(12.8)$$

The trick is to figure out what δ_k^p should be for each unit k in the network. The interesting result, which we now derive, is that there is a simple recursive computation of these δ 's which can be implemented by propagating error signals backward through the network.

To compute δ^p we apply the chain rule to write this partial derivative as the product of two factors, one factor reflecting the change in error as a function of the output of the unit and one reflecting the change in the output as a function of changes in the input. Thus, we have

$$\delta_k^p = \frac{\partial E^p}{\partial S_k^p} \frac{\partial S_k^p}{\partial y_k^p} \quad \dots(12.9)$$

Let us compute the second factor. By equation (12.1) we see that

$$\frac{\partial y_k^p}{\partial S_k^p} = F'(S_k^p) \quad \dots(12.10)$$

which is simply the derivative of the squashing function F for the k th unit, evaluated at the act input S_k to that unit. To compute the first factor of equation (12.9), we consider two cases. First, assume that unit k is an output unit $k = o$ of the network. In this case, it follows from the definition of E^p that

$$\frac{\partial E^p}{\partial y_o^p} = -(d_o^p - y_o^p) \quad \dots(12.11)$$

which is the same result as we obtained with the standard delta rule. Substituting this and equation (12.10) in equation (12.9), we get

$$\delta_o^p = (d_o^p - y_o^p) F'(S_o^p) \quad \dots(12.12)$$

for any output unit o . Secondly, if k is not an output unit $k = h$, we do not readily know the contribution of the unit to the output error of the network. However, the error measure can be written as a function of the net inputs from hidden to output layer $E_o = E_p(x^p, x^p, \dots, s^p)$ and we use the chain rule to write.

$$\frac{\partial E^p}{\partial y_h^p} = \sum_{o=1}^{N_o} \frac{\partial E^p}{\partial S_o^p} \frac{\partial S_o^p}{\partial y_h^p} = \sum_{o=1}^{N_o} \frac{\partial E^p}{\partial S_o^p} \frac{\partial}{\partial y_h^p} \sum_{j=1}^{N_o} w_{hj} y_j^p = \sum_{j=1}^{N_o} \frac{\partial E^p}{\partial S_o^p} w_{hj} = \sum_{j=1}^{N_o} \delta_o^p w_{hj} \quad \dots(12.13)$$

Substituting this in equation (12.9) yields.

$$\delta_h^p = F'(S_h^p) \sum_{j=1}^{N_o} \delta_o^p w_{hj}$$

Equations (12.12) and (12.14) give a recursive procedure for computing the δ 's for all units in the network, which are then used to compute the weight changes according to equation (12.8). This procedure constitutes the generalized delta rule for a feed-forward network of non-linear units.

12.3.1 Understanding Back-Propagation

The equations derived in the previous section may be mathematically correct, but what do they actually mean?

Is there a way of understanding back-propagation other than reciting the necessary equations?

The answer is, of course, yes. In fact, the whole back-propagation process is intuitively very clear. What happens in the above equations is the following. When a learning pattern is clamped, the activation values are propagated to the output units, and the actual network output is compared with the desired output values, we usually end up with an error in each of the output units. Let's call this error e_o for a particular output unit o . We have to bring e_o to zero.

The simplest method to do this is the greedy method: we strive to change the connections in the neural network in such a way that, next time around, the error e_o will be zero for this particular pattern. We know from the delta rule that, in order to reduce an error, we have to adapt its incoming weights according to

$$\Delta w_{ho} = (d_o - y_o) y_h$$

That is step one. But it alone is not enough: when we only apply this rule, the weights from input to hidden units are never changed, and we do not have the full representational power of the feed-forward network as promised by the universal approximation theorem.

In order to adapt the weights from input to hidden units, we again want to apply the delta rule. In this case, however, we do not have a value for δ for the hidden units. This is solved by the chain rule which does the following: distribute the error of an output unit o to all the hidden units that it is connected to, weighted by this connection. Differently put, a hidden unit h receives a delta from each output unit o equal to the delta of that output unit weighted with (= multiplied by) the weight of the connection between those units.

In symbols: $\delta_h = \sum_o \delta_o w_{ho}$. Well, not exactly: we forgot the activation

0

function of the hidden unit; F' has to be applied to the delta, before the backpropagation process can continue.

WORKING WITH BACK-PROPAGATION

The application of the generalised delta rule thus involves two phases: During the first phase the input x is presented and propagated forward through the network to compute the output values y^p for each output unit. This output is compared with its desired value d_o , resulting in an error signal δ_o^p for each output unit.

The second phase involves a backward pass through the network during which the error signal is passed to each unit in the network and appropriate weight changes are recalculated.

12.4.1 Weight Adjustments with Sigmoid Activation Function

The results from the previous section can be summarised in three equations:

- * The weight of a connection is adjusted by an amount proportional to the product of an error signal δ_o^p on the unit receiving the input and the output of the unit sending this signal along the connection:

$$\Delta_p W_{kj} = \gamma \delta_o^p y_k^p \quad \dots(12.16)$$

- * If the unit is an output unit, the error signal is given by

$$\delta_o^p = -(d_o^p - y_o^p) F'(S_o^p) \quad \dots(12.17)$$

Take as the activation function F the 'sigmoid' function as defined in chapter 2 :

$$y^p = F(S^p) = \frac{1}{1 + e^{-t^p}} \quad \dots(12.18)$$

In this case the derivative is equal to

$$F'(S^p) = \frac{\partial}{\partial S^p} \frac{1}{1 + e^{-S^p}} = \frac{1}{(1 + e^{-S^p})^2} (-e^{-t^p}) = - \frac{1}{(1 + e^{-S^p})^2} (-e^{-S^p}) = y^p(1 - y^p) \quad \dots(12.19)$$

such that the error signal for an output unit can be written as :

$$\delta_o^p = -(d_o^p - y_o^p) y_o^p (1 - y_o^p) \quad \dots(12.20)$$

- * The error signal for a hidden unit is determined recursively in terms of error signals of the units to which it directly connects and the weights of those connections. For the sigmoid activation function:

$$\delta_k^p = F'(S_k^p) \sum_{j=1}^{N_o} d_o^p w_{ho} - y_k^p (1 - y_k^p) \sum_{j=1}^{N_o} d_o^p w_{ho} \quad \dots(12.21)$$

Learning Rate And Momentum

The learning procedure requires that the change in weight is proportional to

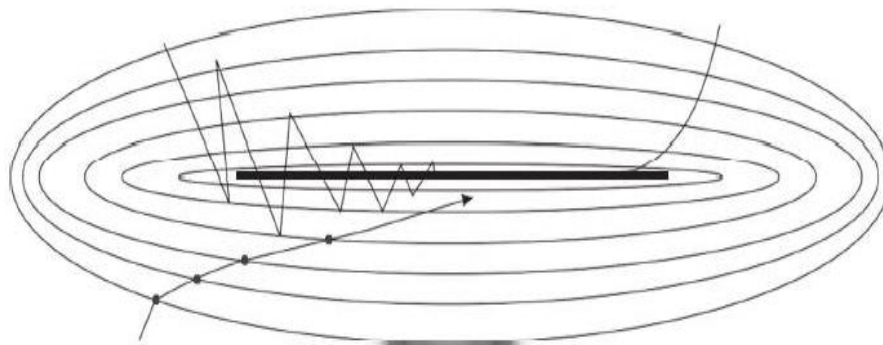
$$\frac{\partial E}{\partial w_{jk}}$$

requires that infinitesimal steps are taken. The constant of proportionality is the learning rate. For practical purposes we choose a learning rate that is as large as possible without leading to oscillation. One way to avoid oscillation at large, is to make the change in weight dependent of the past weight change by adding a momentum term:

$$\Delta w_{jk}(t+1) = \eta \frac{\partial E}{\partial w_{jk}} + a \Delta w_{jk}(t) \quad (12.22)$$

where t indexes the presentation number and a is a constant which determines the effect of the previous weight change.

The role of the momentum term is shown in Fig. 12.2. When no momentum term is used, it takes a long time before the minimum has been reached with a low learning rate, whereas for high learning rates the minimum is never reached because of the oscillations. When adding the momentum term, the minimum will be reached faster.



The descent in weight space. (a) for small learning rate; (b) for large learning rate: note the oscillations, and (c) with large learning rate and momentum term added.

Learning Per Pattern

Although, theoretically, the back-propagation algorithm performs gradient descent on the total error only if the weights are adjusted after the full set of learning patterns has been presented, more often than not the learning rule is applied to each pattern separately, i.e., a pattern is applied, E^p is calculated, and the weights are adapted ($p=1, 2, \dots, P$). There exist empirical indications that this results in faster convergence. Care has to be taken, however, with the order in which the patterns are taught. For example, when using the same sequence over and over again in the network may become focused on the first few patterns. This problem can be overcome by using a permuted training method.

Example 12.1: A feed-forward network can be used to approximate a function from examples. Suppose we have a system (for example a chemical processor or a financial market) of which we want to know the characteristics. The input of the system is given by the two-dimensional vector x and the output is given by the one-dimensional vector d . We want to estimate the relationship $d = f(x)$ from 80 examples $\{x^p, d^p\}$ as depicted in Fig. 12.3 (top left). A feed-forward network was programmed with two inputs, 10 hidden units with sigmoid activation function and an output unit with a linear activation function. Check for yourself how equation (4.20) should be adapted for the linear instead of sigmoid activation function. Then the network weights are initialized to small values and the network is trained for 5,000 learning iterations with the back-propagation training rule, described in the previous section. The relationship between x and d as represented by the network is shown in Fig. 12.3 (top right), while the function which generated the learning samples is given in Fig. 12.3 (bottom left). The approximation error is depicted in Fig. 12.3 (bottom right). We see that the error is higher at the edges of the region within which the learning samples were generated. The network is considerably better at interpolation than extrapolation.

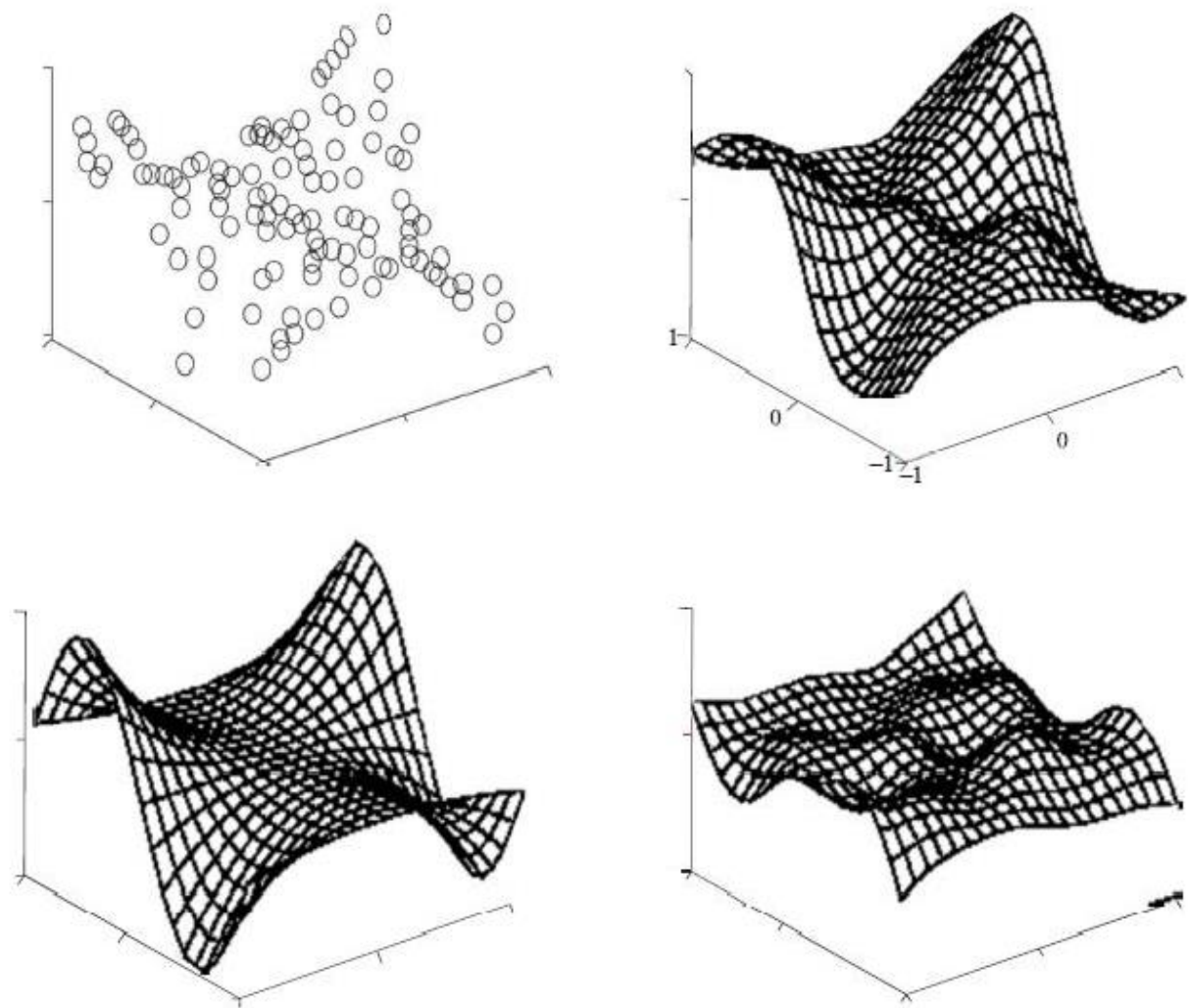


Fig. Example of function approximation with a feed forward network. Top left: The original learning samples; Top right: The approximation with the network; Bottom left: The function which generated the learning samples; Bottom right: The error in the approximation.

Exercise:

- Q1. What is mean by Single-Layer NN Systems.
- Q2. Explain Architecture of a simple ADALINE.
- Q3. What are the use of ADLINE.
- Q4. What are the Applications of Neural Network.
- Q5. What is mean by Learning Algorithm.
- Q6. Explain MULTI-LAYER PERCEPTRONS
- Q7. What is the process of Back Propagation.
- Q8. Write a short note on MULTI - LAYER FEED - FORWARD NETWORKS.
- Q 9. List out the GENERALISED DELTA RULE.
- Q10. How we can understand the Back Propagation.

Q11. Explain WORKINGWITH BACK-PROPAGATION.
Q12. What is mean by Learning Rate And Momentum.

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Chapter 5

Radial Basis and Recurrent Neural Networks: RBF network structure , theorem and the reparability of patterns, RBF learning strategies

Radial Basis and Recurrent Neural Networks: RBF network structure, theorem and the reparability of patterns RBF learning strategies, K-means and LMS algorithms, comparison of RBF and MLP networks: energy function, spurious states, error performance.

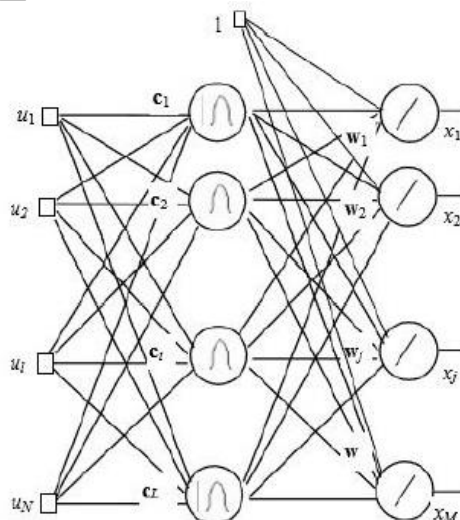
Radial basis function (RBF) networks are feed-forward networks trained using a supervised training algorithm. They are typically configured with a single hidden layer of units whose activation function is selected from a class of functions called basis functions. While similar to backpropagation in many respects, radial basis function networks have several advantages. They usually train much faster than backpropagation networks. They are less susceptible to problems with non-stationary inputs because of the behaviour of the radial basis function hidden units.

Popularized by Moody and Darken (1989), RBF networks have proved to be a useful neural network architecture. The major difference between RBF networks and backpropagation networks (that is, multilayer perceptron trained by Back Propagation algorithm) is the behaviour of the single hidden layer. Rather than using the sigmoidal or S-shaped activation function as in backpropagation, the hidden units in RBF networks use a Gaussian or some other basis kernel function. Each hidden unit acts as a locally tuned processor that computes a score for the match between the input vector and its connection weights or centres. In effect, the basis units are highly specialized pattern detectors. The weights connecting the basis units to the outputs are used to take a linear combination of the hidden unit outputs to produce the final classification or output. In this chapter first the structure of the network will be introduced and it will be explained how it can be used for function approximation and data interpolation. Then it will be explained how it can be trained.

The Structure of the RBF Networks

Radial Basis Functions are first introduced in the solution of the real multivariable interpolation problems. Broomhead and Lowe (1988), and Moody and Darken (1989) were the first to exploit the use of radial basis functions in the design of neural networks.

The structure of an RBF network in its most basic form involves three entirely different layers



Structure of the Standart RBF network

The input layer is made up of source nodes (sensory units) whose number is equal to the dimension of the input vector.

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Hidden layer

The second layer is the hidden layer which is composed of nonlinear units that are connected directly to all of the nodes in the input layer. It is of high enough dimension, which serves a different purpose from that in a multilayer perceptron.

Each hidden unit takes its input from all the nodes at the components at the input layer. As mentioned above the hidden units contain a basis function, which has the parameters center and width. The center of the basis function for a node i at the hidden layer is a vector c_i whose size is the same as the input vector u and there is normally a different center for each unit in the network.

First, the radial distance d_i , between the input vector u and the center of the basis function c_i is computed for each unit i in the hidden layer as

$$d_i = ||u - c_i|| \quad (5.1.1)$$

using the Euclidean distance.

The output h_i of each hidden unit i is then computed by applying the basis function G to this distance.

$$h_i = G(d_i, \sigma_i) \quad (5.1.2)$$

As it is shown in Figure 5.2, the basis function is a curve (typically a Gaussian function, the width corresponding to the variance, σ_i) which has a peak at zero distance and it decreases as the distance from the center increases.

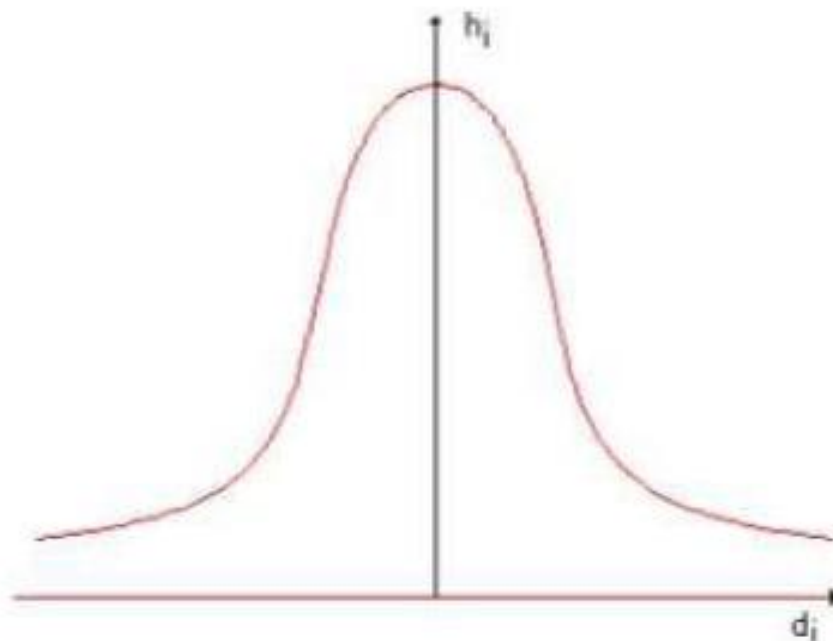


Figure 5.2. The response region of an RBF hidden node around its center as a function of the distance from this center.

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For an input space $u \in \mathbb{R}^2$, that is $M=2$, this corresponds to the two-dimensional Gaussian centered at the origin in the input space, where also $c \in \mathbb{R}^2$, as it is shown in Figure 5.3

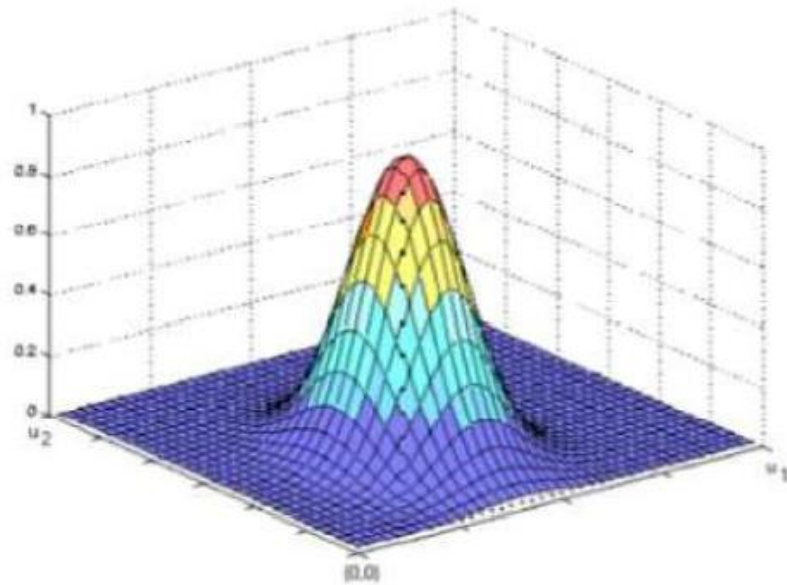


Figure 5.3 Response of a hidden unit on the input space for $u \in \mathbb{R}^2$

5.1.2 Output layer

The transformation from the input space to the hidden units space is nonlinear, whereas the transformation to the hidden units space to the output space is linear.

The j th output is computed as

$$x_j = f_j(u) = w_{oj} + \sum_{(w)} w_{yj} h_i \quad j = 1, 2, \dots, M \quad (5.1.3)$$

Mathematical model

In summary, the mathematical model of the RBF network can be expressed as;

$$x = f(u), f: \mathbb{R}^N \rightarrow \mathbb{R}^M \quad (5.1.4)$$

$$x_j = f_j(u) = w_{oj} + \sum_{|u|} w_{gj} G(|u - c_j|) \quad j = 1, 2, \dots, M \quad (5.1.5)$$

where $d(u, c_i)$ is the Euclidean distance between u and c_i

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5.2 Function approximation

Let $y=g(u)$ be a given function of u , $y \in \mathbb{R}$, $u \in \mathbb{R}$, $g: \mathbb{R} \rightarrow \mathbb{R}$, $G_i, i=1..L$, be a finite set of basis functions.

The function g can be written in terms of the given basis functions as

$$y = g(u) = \sum_{|u|}^t w_t G_t(u) + r(u) \quad (5.2.1)$$

where $r(u)$ is the residual.

The function y can be approximated as

$$y = g(u) \cong \sum_{|u|}^t w_t G_t(u) \quad (5.2.2)$$

The aim is to minimize the error by setting the parameters of G_i appropriately. A possible choice for the error definition is the L2 norm of the residual function $r(u)$ which is defined as

$$\|y(u)\|_{1.2} = \|r(u)\|^2 \quad (5.2.3)$$

5.2.1 Approximation by RBFNN

Now, consider the single input single output RBF network shown in Figure 5.4. Then x can be written as

$$x = f(u) = \sum_{|u|} w_t G_t(\|u - c_t\|) \quad (5.2.4)$$

By the use of such a network, y can be written as

$$y = \sum_{|u|} w_t G(\|u - c_t\|) + r(u) = f(u) + r(u) \quad (5.2.5)$$

where $f(u)$ is the output of the RBFNN given in Figure 5.4 and $r(u)$ is the residual.

By setting the center c_i , the variance σ_i , and the weight w_i the error appropriately, the error can be

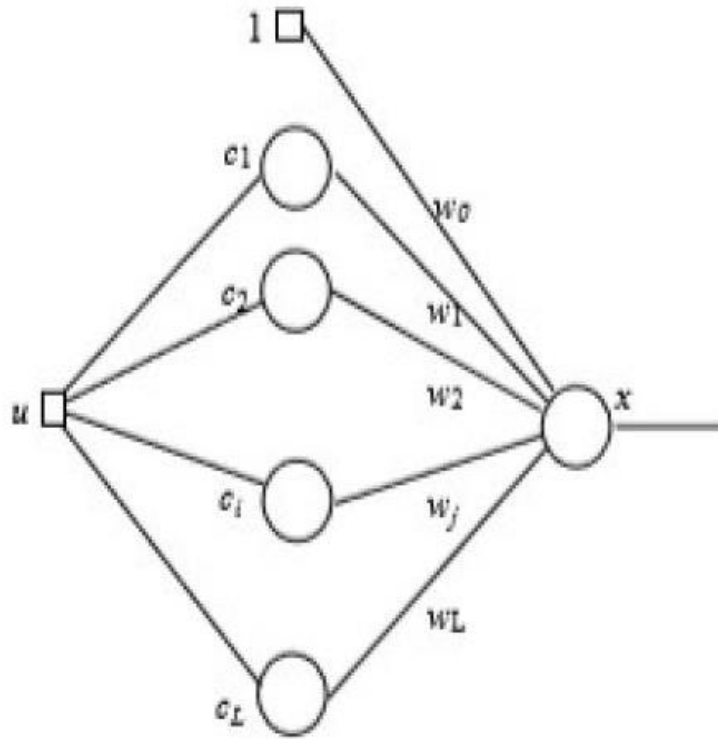


Figure 5.4 Single input, single output RBF network

Whatever we discussed here for $g: \mathbb{R} \rightarrow \mathbb{R}$, can be generalized to $g: \mathbb{R}^N \rightarrow \mathbb{R}^M$ easily by using an N input, M output RBFNN given in figure 5.1 previously.

5.2.2 Data Interpolation

Given input/output training patterns $(u^k, y^k)_{k=1,2,\dots,K}$, the aim of data interpolation is to approximate the function y from which the data is generated. Since the function is unknown, the problem can be stated as a minimization problem which takes only the sample points into consideration:

Choose \mathbf{w}_{ij} and \mathbf{c}_i , $i=1,2,\dots,L$, $j=1,2,\dots,M$ so as to minimize

$$J(\mathbf{w}, \mathbf{c}) = \sum_{k=1}^K ||y^k - f(u^k)||^2 \quad (5.2.6)$$

As an example, the output of an RBF network trained to fit the datapoints given in Table 5.1 is given in Figure 5.5.

TABLE I: 13 datapoints generated by using sum of three gaussians with $c_1=0.2000$

$c_2 = 0.6000$

$c_3=0.9000$ $w_1 = 0.2000$ $w_2 = 0.5000$ $w_3 = 0.3000$ $\sigma = 0.1000$

data no	1	2	3	4	5	6	7	9	10	11	12	13
x	0.0500	0.2000	0.2500	0.3000	0.4000	0.4300	0.4800	0.6000	0.7000	0.8000	0.9000	0.9500
f(x)	0.863	0.2662	0.2362	0.1687	0.1260	0.1756	0.3290	0.6694	0.4573	0.3320	0.4063	0.3535

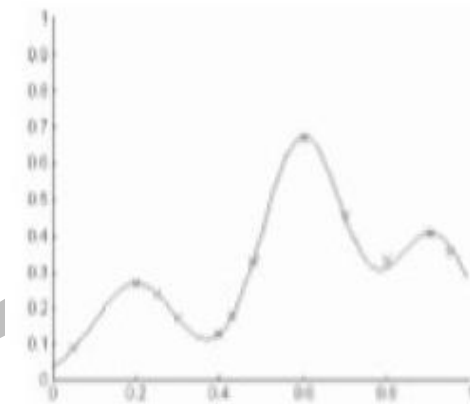


Figure 5.5 Output of the RBF network trained to fit the datapoints given in Table 5.1

5.3 Training RBF Networks

The training of an RBF network can be formulated as a nonlinear unconstrained optimization problem given below:

Given input/output training patterns (u^k, y^k) , $k=1, 2, \dots, K$, Choose w_{ij} and c_i , $i=1, 2, \dots, L$, $j=1, 2, \dots, M$ so as to minimize

Note that the training problem becomes quadratic once if c_i 's (radial basis function centers) are known.

5.3.1 *Adjusting the widths*

In its simplest form, all hidden units in the RBF network have the same width or degree of sensitivity to inputs. However, in portions of the input space where there are few patterns, it is sometimes desirable to have hidden units with a wide area of reception. Likewise, in portions of the input space, which are crowded, it might be desirable to have very highly tuned processors with narrow reception fields. Computing these individual widths increases the performance of the RBF network at the expense of a more complicated training process.

5.3.2 *Adjusting the centers*

Remember that in a backpropagation network, all weights in all of the layers are adjusted at the same time. In radial basis function networks, however, the weights into the hidden layer basis units are usually set before the second layer weights are adjusted. As the input moves away from the connection weights, the activation value falls off. This behavior leads to the use of the term "center" for the first-layer weights. These center weights can be computed using Kohonen feature maps, statistical methods such as K-Means clustering, or some other means. In any case, they are then used to set the areas of sensitivity for the RBF network's hidden units, which then remain fixed.

5.3.3 *Adjusting the weights*

Once the hidden layer weights are set, a second phase of training is used to adjust the output weights. This process typically uses the standard steepest descent algorithm. Note that the training problem becomes quadratic once if c_i 's (radial basis function centers) are known.

Exercise

- Q1. Write a short on RadialBasisandRecurrentNeuralNetworks
- Q2. Explain the Structure of the RBF Networks.
- Q3. What is mean by Hidden layer.
- Q4. What is mean by Function approximation.
- Q5. Write a short note on DataInterpolation.
- Q6. Write a short note on Adjusting the centers.
- Q7. Write a short note on Adjusting the widths.

Chapter 6

K-means and LMS algorithms, Comparison of RBF and MLP networks

k-means clustering algorithm

k-means is one of the simplest unsupervised learning algorithms that solve the well-known clustering problem. The procedure follows a simple and easy way to classify a given dataset through a certain number of clusters (assume k clusters) fixed a priori. The main idea is to define k centers, one for each cluster. These centers should be placed in a cunning way because of different locations cause different results. So, the better choice is to place them as much as possible far away from each other. The next step is to take each point belonging to a given dataset and associate it to the nearest center. When no point is pending, the first step is completed and a new groupage is done. At this point we need to recalculate k new centroids as barycenter of the clusters resulting from the previous step. After we have these k new centroids, a new binding has to be done between the same dataset points and the nearest new center. A loop has been generated. As a result of this loop we may notice that the k centers change their location step by step until no more changes are done or in other words centers do not move anymore. Finally, this algorithm aims at minimizing an objective function known as squared error function given by:

$$J(V) = \sum_{i=1}^c \sum_{j=1}^{c_i} (||x_i - v_j||)^2$$

where,

‘ $||x_i - v_j||$ ’ is the Euclidean distance between x_i and v_j .

‘ c_i ’ is the number of data points in i^{th} cluster.

‘ c ’ is the number of cluster centers.

Algorithmic steps for k-means clustering

Let $X = \{x_1, x_2, x_3, \dots, x_n\}$ be the set of data points and $V = \{v_1, v_2, \dots, v_c\}$ be the set of centers.

- 1) Randomly select ‘ c ’ cluster centers.
- 2) Calculate the distance between each data point and cluster centers.
- 3) Assign the data point to the cluster center whose distance from the cluster center is minimum of all the cluster centers.
- 4) Recalculate the new cluster center using:

$$V_i = \left(\frac{1}{c_i} \sum_{j=1}^{c_i} x_j \right)$$

where, ' c_i ' represents the number of data points in i^{th} cluster.

- 5) Recalculate the distance between each data point and new obtained cluster centers.

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Advantages

- 1) Fast, robust and easy to understand.
- 2) Relatively efficient: $O(knd)$, where n is #clusters, d is #dimension of each object, and t is #iterations.
Normally, $k, t, d \ll n$.
- 3) Gives best result when dataset are distinct or well separated from each other.

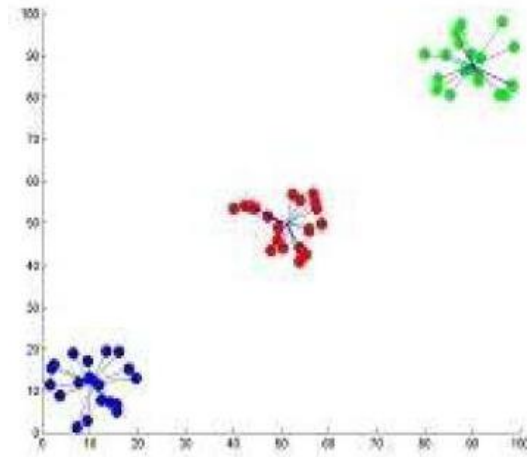


Fig. I : Showing the result of k-means for ‘N’ = 60 and ‘c’ = 3

Note : For more detailed figure for k-means algorithm please refer to [k-means figures](#) sub page.

Disadvantages

- 1) The learning algorithm requires prior specification of the number of cluster centers.
- 2) The use of Exclusive Assignment - If there are two highly overlapping data then k-means will not be able to resolve that there are two clusters.
- 3) The learning algorithm is not invariant to non-linear transformation i.e. with different representation of data we get different results (data represented in form of cartesian co-ordinates and polar co-ordinates will give different results).
- 4) Euclidean distance measures can unequally weight underlying factors.
- 5) The learning algorithm provides the local optima of the squared error function.

- 5) The learning algorithm provides the local optima of the squared error function.
- 6) Randomly choosing of the cluster center cannot lead to the fruitful result. Pl. refer Fig.
- 7) Applicable only when mean is defined i.e. fails for categorical data.
- 8) Unable to handle noisy data and outliers.
- 9) Algorithm fails for non-linear data set.

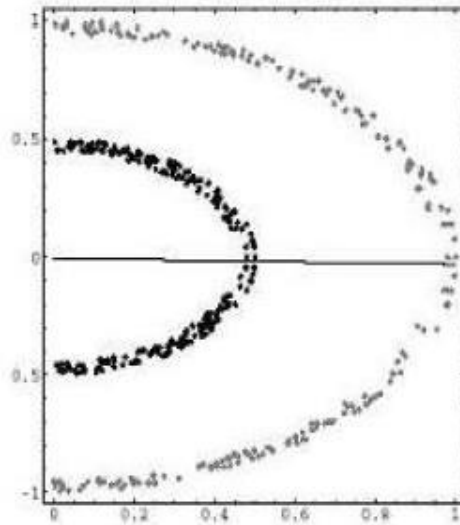


Fig II : Show the non-linear data set where k-means algorithm fails.

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6.4 Least-Mean-Square Algorithm

The least-mean-square (LMS) algorithm is based on the use of instantaneous estimates of the autocorrelation function $r_{tt}(j, k)$ and the cross-correlation function $r_{ti}(k)$. These estimates are deducted directly from the defining equations (6.8) and (6.7) as follows:

$$p_i(j, k, n) = x_i(m)x_i(n) \quad (6.20)$$

and

$$p_{tt}(k, n) = x_i(m)x_i(n)$$

The use of \hat{p}_{ti} and \hat{p}_{tt} is intended to signify that these quantities are “estimates.” The definitions introduced in Eqs. (6.20) and (6.21) have been generalized to include a nonstationary environment, in which case all the sensory signals and the desired response assume time-varying forms too. Thus, substituting $\hat{p}_i(j, k, n) = x_i(m)x_i(n)$ and $\hat{p}_{tt}(k, n)$ in place of $r_{ti}(j, k)$ and $r_{tt}(k)$ in Eq. (6.17). We get

$$\begin{aligned} w_y^t(n+1) &= w_y^t(n) + \eta [x(n)f(n) - \sum_{j=1}^t w_y(n)x_j(n)x(n)] \\ &= w_y^t(n) + \eta [d(n) - \sum_{j=1}^t w_y(n)x_j(n)]x(n) \\ &= w_y^t(n) + \eta [d(n) - y(n)]x(n), \quad k=1, 2, \dots, P. \end{aligned} \quad (6.22)$$

where $y(n)$ is the output of the spatial filter computed at iteration n in accordance with the LMS algorithm; that is,

$$d(n) = \sum_{j=1}^t w_y(n)x_j(n) \quad (6.23)$$

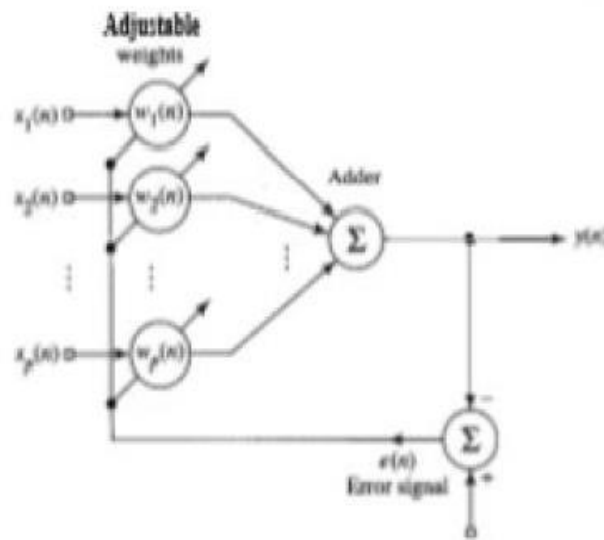
Note that in Eq. (6.22) we have used $\hat{w}_y(n)$ in place of $w_y(n)$ to emphasize the fact that Eq. (6.22) involves “estimates” of the weights of the spatial filter.

Figure 6.3 illustrates the operational environment of the LMS algorithm, which is completely described by Eqs. (6.22) and (6.23). A summary of the LMS algorithm is presented in Table 6.1, which clearly illustrates the simplicity of the algorithm. As indicated in this table, for the initialization of the algorithm, it is customary to set all the initial values of the weights of the filter equal to zero.

In the method of steepest descent applied to a “known” environment, the weight vector $w(n)$, made up of the weights $w_1(n), w_2(n), \dots, w_t(n)$, starts at some initial value $w(i)$, and then follows a precisely defined trajectory (along the errors surface) that eventually terminates on the optimum solution w , provided that the learning-rate parameter η is chosen properly. In contrast, in the LMS algorithm applied to an “unknown” environment, the weight vector $w(n)$, representing an “estimate” of w , follows a random trajectory. For this reason, the LMS algorithm is sometimes referred to as a “stochastic gradient algorithm.” As the number of iterations in the LMS algorithm approaches

infinity, $w(n)$ performs a random walk (Brownian motion) about the optimum solution w ; see Appendix D.

DDOL



Desired response

FIGURE 6.3 Adaptive spatial filter.

Another way of stating the basic difference between the method of steepest descent and the LMS algorithm is in terms of the error calculations involved. At any iteration n , the method of steepest descent minimizes the mean-squared error $j(n)$. This cost function involves ensemble averaging, the effect of which is to give the method of steepest descent an “exact” gradient vector that improves in pointing accuracy with increasing n . The LMS algorithm, on the other hand, minimizes an instantaneous estimate of the cost function $j(n)$. Consequently, the gradient vector in the LMS algorithm is “random,” and its pointing accuracy improves “on the average” with increasing n .

The basic difference between the method of steepest descent and the LMS algorithm may also be stated in terms of time-domain ideas, emphasizing other aspects of the adaptive filtering problem. The method of steepest descent minimizes the sum of error squares $\sum_{n=1}^N e^2(n)$, integrated over all previous iterations of the algorithm up to and including estimates of the autocorrelation function \hat{y} , and cross-correlation function \hat{y}_n . In contrast, the LMS algorithm simply minimizes the instantaneous error squared $e^2(n)$, defined as $(\frac{1}{2})e^2(n)$, thereby reducing the storage requirement to the minimum possible. In particular, it does not require storing any more information than is present in the weights of the filter.

TABLE 6.1 Summary of the LMS Algorithm

1. Initialization. Set

$$w_v(1) = 0 \quad \text{for } v=1, 2, \dots, p.$$

2. Filtering. For time $n=1, 2, \dots$, compute.

$$y(n) = \sum_{j=1}^p w_v(n) x_j(n)$$

$$e(n) = d(n) - y(n)$$

$$w_v^*(n+1) = w_v(n) + \eta y(n) x_v(n) \quad \text{for } v=1, 2, \dots, p$$

It is also important to recognize that the LMS algorithm can operate in a stationary or nonstationary environment. By a “nonstationary” environment we mean one in which the statistics vary with time. In such a situation the optimum solution assumes a time-varying form, and the LMS algorithm therefore has the task of not only seeking the minimum point of the error surface but also tracking it. In this context, the smaller we make the learning-rate parameter, the better will be the tracking behaviour of the algorithm. However, this improvement in performance is attained at the cost of a slow adaptation rate (Haykin, 1991; Widrow and Stearns, 1985).

Signal-Flow Graph Representation of the LMS Algorithm

Equation (6.22) provides a complete description of the time evolution of the weights in the LMS algorithm. Rewriting the second line of this equation in matrix form, we may express it as follows:

$$\mathbf{W}(n+1) = \mathbf{W}(n) + n[d(n) - \mathbf{x}^T(n)\mathbf{W}(n)]\mathbf{x}(n) \quad (6.24)$$

where

$$\mathbf{W}(n) = [w_1(n), w_2(n), \dots, w_t(n)]^T \quad (6.25)$$

and

$$\mathbf{x}(n) = [x_1(n), x_2(n), \dots, x_t(n)]^T \quad (6.26)$$

Rearranging terms in Eq. (6.24), we have

$$\mathbf{W}(n+1) = [\mathbf{I} - n\mathbf{x}(n)\mathbf{x}^T(n)]\mathbf{W}(n) + n\mathbf{x}(n)d(n) \quad (6.27)$$

where \mathbf{I} is the identity matrix. In using the LMS algorithm, we note that

$$\mathbf{W}(n) = z^{-1}[\mathbf{W}(n+1)] \quad (6.28)$$

where z^{-1} is the unit-delay operator implying storage. Using Eqs. (6.27) and (6.28), we may thus represent the LMS algorithm by the signal-flow graph depicted in Fig. 6.4.

The signal-flow graph of Fig. 6.4 reveals that the LMS algorithm is an example of a stochastic feedback system. The presence of feedback has a profound impact on the convergence behaviour of the LMS algorithm, as discussed next.

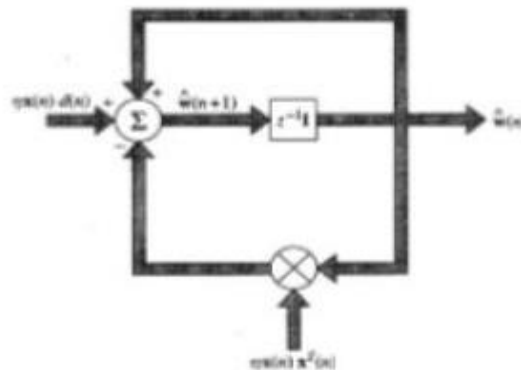


FIGURE 6.4 Signal-flow graph representation of the LMS algorithm.

- Q1. What is mean by k-means clustering algorithm.
Q2. What are the Algorithmic steps for k-means clustering.
Q3. Write a Advantages and Dis Advantage of k-means clustering.
Q4. Explain Least-Mean-Square Algorithm.
Q5. Explain Signal-Flow Graph Representation of the LMS Algorithm

DDOL

Chapter 7

Hopfield networks: energy function, spurious states, error performance.

Hopfield Network

Hopfield neural network was invented by Dr. John J. Hopfield in 1982. It consists of a single layer which contains one or more fully connected recurrent neurons. The Hopfield network is commonly used for auto-association and optimization tasks.

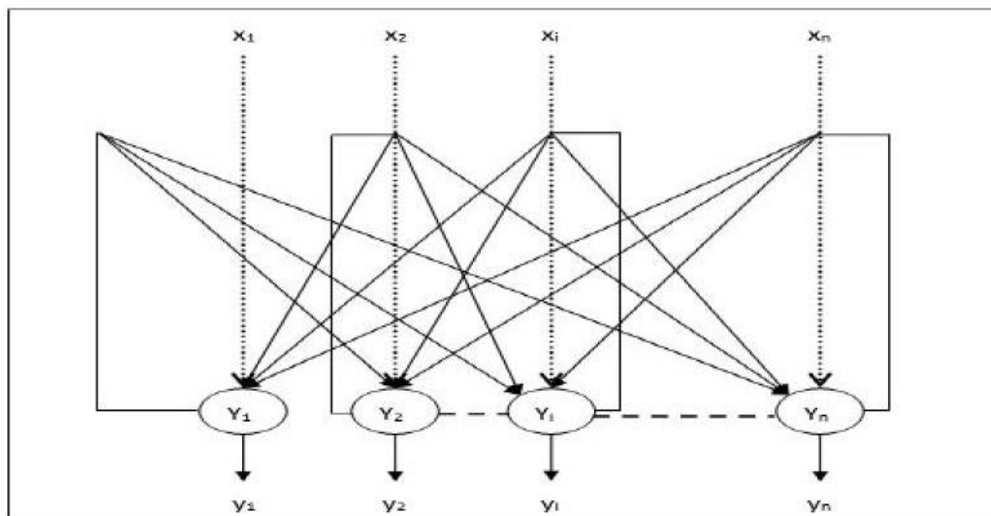
Discrete Hopfield Network

A Hopfield network which operates in a discrete line fashion or in other words, it can be said the input and output patterns are discrete vector, which can be either binary (0,1) or bipolar (+1,-1) in nature. The network has symmetrical weights with no self connections i.e., $w_{ij} = w_{ji}$ and $w_{ii} = 0$.

Architecture

Following are some important points to keep in mind about discrete Hopfield network -

- * This model consists of neurons with one inverting and one non-inverting output.
- * The output of each neuron should be the input of other neurons but not the input of self.
- * Weight/connection strength is represented by w_{ij} .
- * Connections can be excitatory as well as inhibitory. It would be excitatory, if the output of the neuron is same as the input, otherwise inhibitory.
- * Weights should be symmetrical, i.e. $w_{ij} = w_{ji}$.



The output from Y_1 going to Y_2 , Y_i and Y_n have the weights w_{12} , w_{1i} and w_{1n} respectively.

Similarly, other arcs have the weights on them.

DDOL

During training of discrete Hopfield network, weights will be updated. As we know that we can have the binary input vectors as well as bipolar input vectors. Hence, in both the cases, weight updates can be done with the following relation.

Case 1 - Binary input patterns

For a set of binary patterns $s(p)$, $p = 1$ to P

Here, $s(p) = s_1(p), s_2(p), \dots, s_i(p), \dots, s_n(p)$

Weight Matrix is given by

$$w_{ij} = \sum_{p=1}^P [2s_i(p) - 1][2s_j(p) - 1] \text{ for } i \neq j$$

Case 2 - Bipolar input patterns

For a set of binary patterns $s(p)$, $p = 1$ to P

Here, $s(p) = s_1(p), \dots, s_2(p), \dots, s_i(p), \dots, s_n(p)$

Weight Matrix is given by

$$w_{ij} = \sum_{p=1}^P [s_i(p)][s_j(p)] \text{ for } i \neq j$$

Step 1- Initialize the weights, which are obtained from training algorithm by using Hebbian principle.

Step 2- Perform steps 3-9, if the activation of the network is not consolidated.

Step 3- For each input vector x , perform steps 4-8.

Step 4- Make initial activation of the network equal to the external input vector x as follows-

$$y_i = x_i \text{ for } i = 1 \text{ to } n$$

Step 5- For each unit y_i , perform steps 6-9.

Step 6- Calculate the net input of the network as follows-

$$Y_{ini} = x_i + \sum_j y_j w_{ji}$$

Step 7- Apply the activation as follows over the net input to calculate the output-

$$y_i = \begin{cases} 1 & \text{if } Y_{ini} > \theta_i \\ Y_{ini} & \\ 0 & \text{if } Y_{ini} < \theta_i \end{cases}$$

Here θ_i is the threshold.

Step9- Test the network for conjunction.

Energy Function Evaluation

An energy function is defined as a function that is bounded and non-increasing function of the state of the system. Energy function also called Lyapunov function determines the stability of discrete Hopfield network, and is characterized as follows-

$$E_f = -\frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n y_i y_j w_{ij} - \sum_{i=1}^n x_i y_i + \sum_{i=1}^n \theta_i y_i$$

Condition : In a stable network, whenever the state of node changes, the above energy function will decrease. Suppose when node i has changed state from $y_i^{(k)}$ to $y_i^{(k+1)}$ then the Energy change ΔE_f is given by the following relation.

$$\begin{aligned} \Delta E_f &= E_f(y_i^{(k+1)}) - E_f(y_i^{(k)}) \\ &= - \left(\sum_{j=1}^n w_{ij} y_j^{(k)} + x_i - \theta_i \right) (y_i^{(k+1)} - y_i^{(k)}) \\ &= - (net)_i \Delta y_i \end{aligned}$$

Here $\Delta y_i = y_i^{(k+1)} - y_i^{(k)}$

The change in energy depends on the fact that only one unit can update its activation at a time.

Continuous Hopfield Network

In comparison with Discrete Hopfield network, continuous network has time as a continuous variable. It is also used in auto association and optimization problems such as travelling salesman problem.

Model: The model or architecture can be built up by adding electrical components such as amplifiers which can map the input voltage to the output voltage over a sigmoid activation function.

Energy Function Evaluation :

$$E_f = \frac{1}{2} \sum_{i=1}^n \sum_{j=1, j \neq i}^n y_i y_j w_{ij} - \sum_{i=1}^n x_i y_i + \frac{1}{\square} \sum_{i=1}^n \sum_{j=1, j \neq i}^n w_{ij} g_{ij} \int_{-\infty}^{y_i} a^{-1}(y) dy$$

Here \square is gain parameter and g_{ij} input conductance.

Some important points about Boltzmann Machine -

- * They use recurrent structure.
- * They consist of stochastic neurons, which have one of the two possible states, either 1 or 0.
- * Some of the neurons in this are adaptive (free state) and some are clamped (frozen state).
- * If we apply simulated annealing on discrete Hopfield network, then it would become Boltzmann Machine.

Reference :

https://www.tutorialspoint.com/artificial_neural_network/artificial_neural_network_hopfield.htm

spurious states

Spurious states are patterns $x \notin P$, where P is the set of patterns to be memorized. In other words, they correspond to local minima in the energy function that shouldn't be there. They can be composed of various combinations of the original patterns or simply the negation of any pattern in the original pattern set. These tend to become present when $\alpha = |P|/N$ (where N is the number of neurons) becomes too high for a certain learning rule.

It turns out that spurious states are important for deriving algorithms in Hopfield networks. Because we know that the dynamical update equations always reduce the energy of a system, spurious minima will trap the network and return incorrect or incomplete results. Typically these spurious minima have a higher energy and smaller basin than real patterns (though this is not guaranteed when $|P|$ is too large). This naturally leads to stochastic solution using a Monte Carlo type approach, where enough energy is given to the neurons so that they aren't stuck in the local minima but don't jump out of the closest correct pattern minimum (these are Boltzmann machines).

Here's some hand-

wavy intuition. The learning rules project the current configuration of the network into the subspace spanned by the pattern vectors and then calculate the pattern vector that lies closest to the projected configuration vector. But even if you had completely orthogonal patterns, you cannot specify more patterns than the number of neurons (because then either you duplicate a pattern or the next pattern you add isn't orthogonal).

The real problem is that most learning rules give $\alpha \ll N \alpha \ll N$ (e.g. the Hebb rule gives $\alpha \approx 0.138$ using mean-field derivations) because the projection into the subspace is not orthogonal. This is not an issue if the patterns

themselves are orthogonal (i.e. completely uncorrelated), but that is very rarely the case in practice.

There are ways to “unlearn” these spurious minima too. See [this question](#) for good references, especially check the Rojas book which is available for free online. Also, if you can get your hands on the Hertz book, look at Eq. (10.22) which is the mean field equation whose solutions give the possible states, including spurious ones (they also give an explanation for how to find them specifically).

Exercise:

- Q1. Write short note on Hopfield Network.
- Q2. Explain Discrete Hopfield Network.
- Q3. What are the important points of discrete Hopfield network.
- Q4. Write a note on Training Algorithm with its cases
- Q5. Explain Energy Function Evaluation.
- Q6. What is mean by Continuous Hopfield Network.

Chapter 8

Simulated Annealing: The Boltzmann machine, Boltzmann learning rule, Bidirectional Associative Memory.

Simulated Annealing

Statistical Mechanics and the Simulated Annealing

The starting point of statistical mechanics is an energy function. We consider a physical system with a set of probabilistic states $x = \{x\}$, each of which has energy $E(x)$. For a system at a temperature $T > 0$, its state x varies with time, and quantities such as E that depend on the state fluctuate. Although there must be some driving mechanism for these fluctuations, part of the idea of temperature involves treating them as random. When a system is first prepared, or after a change of parameters, the fluctuations have on average a defined direction such that the energy E decreases. However, sometimes later, any such trend ceases and the system just fluctuates around a constant average value. Then we say that the system is in thermal equilibrium.

A fundamental result from physics tells us that in thermal equilibrium each of the possible states x occurs with probability, determined according to Boltzmann-Gibbs distribution,

$$P(x) = \frac{1}{Z} e^{-\frac{E(x)}{T}}$$

where the normalizing factor

$$Z = \sum_x e^{-\frac{E(x)}{T}}$$

is called the partition function and it is independent of the state x but temperature.

The Boltzmann-Gibbs distribution is usually derived from very general assumptions about microscopic dynamics of materials. The coefficient T is related to absolute temperature T_a of the system as

$$T = k_B T_a$$

where coefficient k_B is Boltzmann's constant having value 1.38×10^{-16} erg/K. Interestingly enough, the same

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distribution can also be achieved in the viewpoint of information theory. Although the temperature T has no physical meaning in information theory, it is interpreted as a pseudotemperature in an abstract manner.

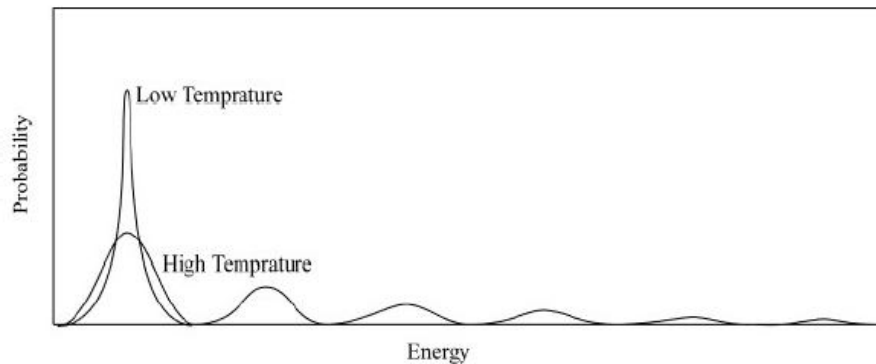
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Given a stated distribution function $f_d(x)$, let $P(x(k)=x_i)$ be the probability of the system being at state x_i at the present time k . Furthermore, let $P(x(k+1)=x_j | x(k)=x_i)$ represent the conditional probability of next state x_j given the present state is x_i . The notation $P(x_i)$ and $P(x_j | x_i)$ will be used simply to denote these probabilities respectively. In equilibrium the stated distribution and the state transition reaches a balance satisfying:

$$P(x^j | x^i) P(x^i) = P(x^i | x^j) P(x^j)$$

Therefore, in equilibrium the Boltzmann Gibbs distribution given by equation (8.1.1) results in :

$$P(x^j | x^i) = \frac{1}{1 + e^{\beta(E(x^j) - E(x^i))}}$$



where $\beta = \frac{1}{k_B T}$ and $E(x^i)$ is the energy of state x^i .

Figure 8.1 Relation between temperature and probability of the states [Kung 93]

The Metropolis algorithm provides a simple method for simulating the evolution of physical systems in a heat bath to thermal equilibrium [Metropolis et al.]. It is based on Monte Carlo simulation technique, which aims to approximate the expected value $\langle g(x) \rangle$ of some function $g(x)$ of a random vector x with a given density function $f_d(x)$. For this purpose several x vectors, say $x = X_k$ $k=1..K$, are randomly generated according to the density function $f_d(x)$ and then Y_k is found as $Y_k = g(X_k)$. By using the strong law of large numbers:

$$\lim_{K \rightarrow \infty} \frac{1}{K} \sum_{k=1}^K Y_k = \langle Y^k \rangle = \langle g(x) \rangle$$

the average of generated Y vectors can be used as an estimate of $\langle g(x) \rangle$ [Sheldon 1989].

In each step of the Metropolis algorithm, an atom (unit) of the system is subjected to a small random displacement, and the resulting change ΔE in the energy of the system is observed. If $\Delta E \leq 0$, then the displacement is accepted, and the new system configuration with the displaced atom is used as the starting point for the next step of

the algorithm. If, on the other hand, $\Delta E > 0$, then the algorithm proceeds in a probabilistic manner so that the configuration with the displaced atom is accepted with a probability given by:

$$P(\Delta E) = e^{-\Delta E / k_B T}$$

MDOL

Provided enough number of transitions in the Metropolis algorithm, the system reaches thermal equilibrium. Thus, by repeating the basic steps of Metropolis algorithm, we effectively simulate the motions of the atoms of a physical system at temperature T . Moreover, the choice of $P(\Delta E)$ defined in Eq. (8.1.8) ensures that thermal equilibrium is characterized by the Boltzmann-Gibbs distribution provided in Eq. (8.1.5).

Referring to Eq. (8.1.5), notice that if $P(x_i) > P(x_j)$ implies $E(x_i) < E(x_j)$, and vice versa. So maximizing the probability function is equivalent to minimizing the energy function. Furthermore, notice that this property is independent of the temperature, although the discrimination becomes more apparent as the temperature decreases (Figure 8.1).

Therefore, the temperature parameter T provides a new free parameter for steering the steps size toward the global optimum. With a high temperature, the equilibrium can be reached more rapidly. However, if the temperature is too high, all the states will have a similar level of probability. On the other hand, when $T \rightarrow 0$, the average state becomes very close to the global minimum. This idea, though very attractive at the first glance, cannot be implemented directly in practice. In fact, with a low temperature, it will take a very long time to reach equilibrium and, more seriously, the state is more easily trapped by local minima. Therefore, it is necessary to start at a high temperature and then decrease it gradually. Correspondingly, the probable states then gradually concentrate around the global minimum (Figure 8.2).

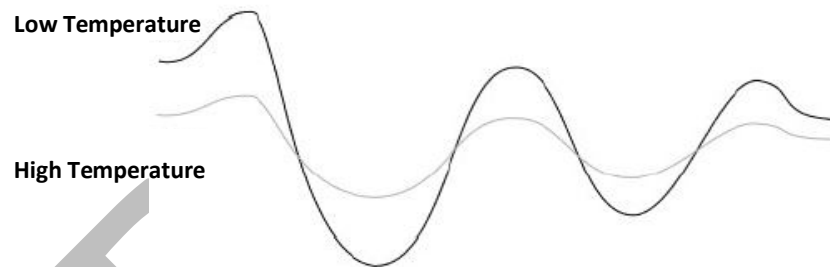


Figure 8.2 The energy levels adjusted for high and low temperature

This has an analogy with metallurgical annealing, in which a body of metal is heated near to its melting point and is then slowly cooled back down to room temperature. This process eliminates dislocations and other crystal lattice disruptions by thermal agitation at high temperature. Furthermore, it prevents the formation of new dislocations by cooling the metal very slowly. This provides necessary time to repair any dislocation that occurs as the temperature drops. The essence of this process is that global energy function of the metal will eventually reach an absolute minimum value.

If the material is cooled rapidly, its atoms are often captured in unfavorable locations in the lattice. Once the temperature has dropped far below the melting point, these defects survive forever, since any local rearrangement of atoms costs more energy than whatever is available in thermal fluctuations. The atomic lattice thus remains captured in a local energy minimum. In order to escape from local minima and to have the lattice in the global energy minimum, the thermal fluctuations can be enhanced by reheating the material until energy-consuming local rearrangements occur at a reasonable rate. The lattice imperfections then start to move and annihilate, until the atomic lattice is free of defects—except for those caused by thermal fluctuations. These can be gradually reduced if the temperature is decreased so slowly that thermal equilibrium is maintained at all times during the cooling process. How much time must be spent for the cooling process depends on the specific situation. A great deal of experience is required to perform the annealing in an optimal way. If the temperature is decreased quickly, some thermal fluctuations are frozen in. On the other hand, if one proceeds too slowly, the process never ends.

IDOL

The amazing thing about annealing is that the statistical process of thermal agitation leads to approximately the same final energy state. This result is independent of the initial condition of the metal and any of the details of the statistical annealing process. The mathematical concept of simulated annealing derives from an analogy with this physical behavior.

The simulated annealing algorithm, is a variant of the Metropolis algorithm in which the temperature is time dependent. In an analogy with metallurgical annealing, it starts with a high temperature and gradually decreases it. At each temperature, it applies several times the update rule given by Eq. (8.1.8). An annealing schedule specifies a finite sequence of temperature values and a finite number of transitions attempted at each value of the temperature. The annealing schedule developed by [Kirkpatrick et al 1983] is as follows.

The initial value T_0 of the temperature is chosen high enough to ensure that virtually all proposed transitions be accepted by the simulated annealing algorithm. Then the cooling is performed. At each temperature, enough transitions are attempted so that there is a predetermined number of transitions per experiment on the average. At the end, the system is frozen and annealing stops if the desired number of acceptances is not achieved at a predetermined number of successive temperatures. In the following, we provide the annealing procedure in more detail:

SIMULATED ANNEALING

- Step 1. Set initial values:** assign a high value to temperature as $T(0) = T_0$, decide on constants K_T , K_A and K_S , Typical values for which are $0.8 < K_T < 0.99$, $K_A = 10$, and $K_S = 3$.
- Step 2. Decrement the temperature:** $T(k) = K_T \cdot T(k-1)$ where K_T is a constant smaller but close to unity.
- Step 3. Attempt enough number of transitions** at each temperature, so that there are K_A accepted transitions per experiment on the average.
- Step 4. Stop** if the desired number of acceptances is not achieved at K_S successive temperatures else repeat steps 2 and 3.

A very important property of simulated annealing is its asymptotic convergence. It has been proved in [Geman and Geman 84] that if $T(k)$ at iteration k is chosen such that it satisfies

$$T(k) \geq \frac{T_0}{\log(1+k)}$$

provided the initial temperature T_0 is high enough, then the system will converge to the minimum energy configuration. The main drawback of simulated annealing is the large amount of computational time necessary for stochastic relaxation.

Many elementary transformations are performed at each temperature step in order to reach an equilibrium state.

IDOL

The main drawback of simulated annealing is the large amount of computational time necessary for stochastic relaxation. Many elementary transformations are performed at each temperature step in order to reach a near equilibrium state.

Exercise:

- Q1. What is meant by Simulated Annealing.
- Q2. Explain Statistical Mechanics.
- Q3. Explain Simulated Annealing.
- Q4. Write a short note on Boltzmann learning rule.
- Q5. Explain Bidirectional Associative Memory.

9.2 Boltzmann Machine

Boltzmann machine [Hinton et al 83] is a connectionist model having stochastic nature. The structure of the Boltzmann machine is similar to Hopfield network, but it adds some probabilistic component to the output function. It uses simulated annealing concepts, in spite of the deterministic nature in state transition of the Hopfield network [Hinton et al 83, Aarts et al 1986, Allwright and Carpenter 1989, Laarhoven and Aarts 1987].

A Boltzmann machine can be viewed as a recurrent neural network consisting of N two state units. Depending on the purpose, the states can be chosen from binary space, that is $x \in \{0, 1\}^N$ or from bipolar space $x \in \{-1, 1\}^N$. The energy function of the Boltzmann machine is:

$$E(x) = -\frac{1}{2} \sum_i \sum_j w_{ij} x_i x_j - \sum_i \theta_i x_i$$

The connections are symmetrical by definition, that is $w_{ij} = w_{ji}$. Furthermore in the bipolar case, the convergence of the machine requires $w_{ii} = 0$ (or equivalently $\theta_i = 0$). However in the binary case self-loops are allowed.

The objective of a Boltzmann machine is to reach the global minimum of its energy function, which is the state having minimum energy. Similar to simulated annealing algorithm, the state transition mechanism of Boltzmann Machine uses a stochastic acceptance criterion, thus allowing it to escape from its local minima. In a sequential Boltzmann machine, units change their states one by one, while they change state all together in a parallel Boltzmann machine.

Let X denote the state space of the machine, that is the set of all possible states. Among these, the state vectors differing only one bit are called neighboring states. The neighborhood $N_x \subset X$ is defined as the set of all neighboring states of x . Let x_j denote the neighboring state that is obtained from x by changing the state of neuron j . Hence, in binary case we have

$$x_j = \begin{cases} x_i & \text{if } i \neq j \\ 1 - x_i & \text{if } i = j \end{cases} \quad x \in (0, 1)^n, x_j \in N_x$$

In bipolar case, this becomes:

$$x_j = \begin{cases} x_i & \text{if } i \neq j \\ -x_i & \text{if } i = j \end{cases} \quad x \in (-1, 1)^n, x_j \in N_x$$

The difference in energy when the global state of the machine is changed from x to x^j is :

Note that the contribution of the connections $w_{km}, k \neq j, m \neq j$, to $E(x)$ and $E(x^j)$ is identical, furthermore $w_{ij} = w_{ji}$. For the binary case, by using equations (9.2.1) and (9.2.2), we obtain

$$\Delta E(x^j | x) = (2x_j - 1) \left(\sum_i w_{ij} x_i + \theta_j \right) \quad x \in \{0, 1\}^N$$

For the bipolar case it is

$$\Delta E(x^j | x) = (2x_j) \left(\sum_i w_{ij} x_i + \theta_j \right) \quad x \in \{-1, 1\}^N, w_{ii} = 0$$

Therefore, the change in the energy can be computed by considering only local information. In a sequential Boltzmann machine, a trial for a state transition is a two-step process. Given a state x , first a unit j is selected as a candidate to change state. This selection probability usually has uniform distribution over the units. Then a probabilistic function determines whether a state transition will occur or not. The state x^j is accepted with probability

$$P(x^j | x) = \frac{1}{1 + e^{\Delta E(x^j | x) / T}}$$

where T is a control parameter having analogy in temperature. Initially the temperature is set large enough to accept almost all state transitions with probability close to 0.5, and then T is decreased in time to zero (Figure 9.3). With a proper cooling schedule, these sequential Boltzmann machine converges asymptotically to a state having minimum energy.

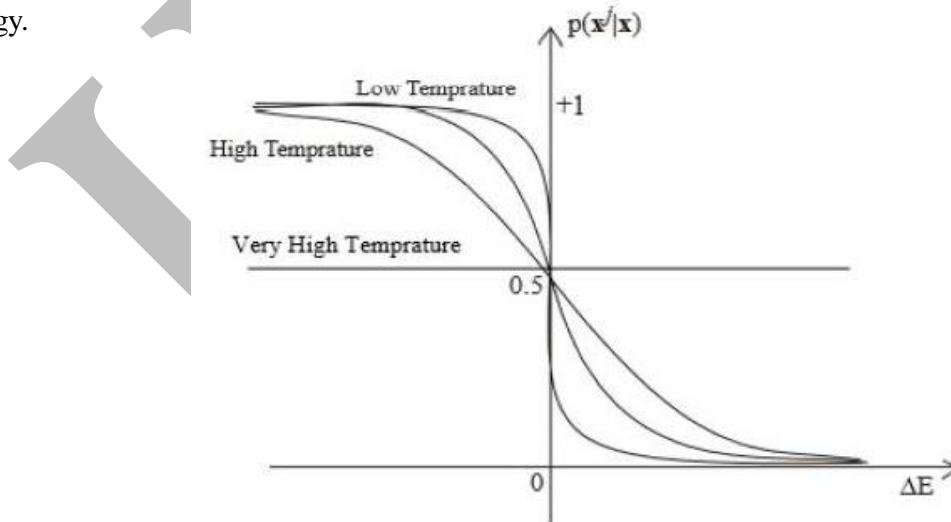


Figure 9.3. Acceptance probability in Boltzmann machine for different temperatures

A Boltzmann machine starts execution with a random initial configuration. Initially, the value of T is very large. A cooling schedule determines how and when to decrement the control parameter. As $T \rightarrow 0$, less and less state transitions occur. If no state transitions occur for a specified number of trials, it is decided that the Boltzmann machine has reached the final state.

A state $x^* \in X$ is called a locally minimal state, if

$$\square E(x^{*j} | x^*) \geq 0 \quad j = 1..N$$

DDOL

Note that, a local minimum is a state whose energy cannot be increased by a single state transition. Let the set of all local minima be denoted by X^* . While the Hopfield network is trapped mostly in one of these local minima, the Boltzmann machine can escape from the local minimum because of its probabilistic nature. Although the machine asymptotically converges to a global minimum, the finite-time approximation of the Boltzmann Machine prevents guaranteeing convergence to a state with minimum energy. However, still the final state of the machine will be a nearly minimum one among X^* .

Use of Boltzmann machine as a neural optimizer involves two phases as it is explained for the Hopfield network in Chapter 4. In the first phase, the connection weights are determined. For this purpose, an energy function for the given application is decided. In the non-constrained optimization applications, the energy function can be directly obtained by using the cost function. However, in the case of constrained optimization, the energy function must be derived using both the original cost function and the constraints. The next step is to determine the connection weights $\{w_{ij}\}$ by considering this energy function. Then in the second phase, the machine searches the global minimum through the annealing procedure.

Exercise:

- Q1. Write short note on Boltzmann Machine.
- Q2. Explain what is mean by Boltzmann learning.
- Q3. What are the rule of Boltzmann learning rule.
- Q4. Write a short note on Learning Per Pattern.
- Q5. Explain the Structure of the RBF Networks.

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Bidirectional Associative Memory

BIDIRECTIONAL ASSOCIATIVE MEMORY (BAM)

Several versions of the heteroassociative recurrent neural network, or bidirectional associative memory (BAM), developed by Kosko (1988, 1992a). - A bidirectional associative memory [Kosko, 1988] stores a set of pattern associations by summing bipolar correlation matrices (an n by m outer product matrix for each pattern to be stored). - The architecture of the net consists of two layers of neurons, connected by directional weighted connection paths. - The net iterates, sending signals back and forth between the two layers until all neurons reach equilibrium (i.e., until each neuron's activation remains constant for several steps). - Bidirectional associative memory neural nets can respond to input to either layer. - Because the weights are bidirectional and the algorithm alternates between updating the activations for each layer, we shall refer to the layers as the X-layer and the Y-layer (rather than the input and output layers).

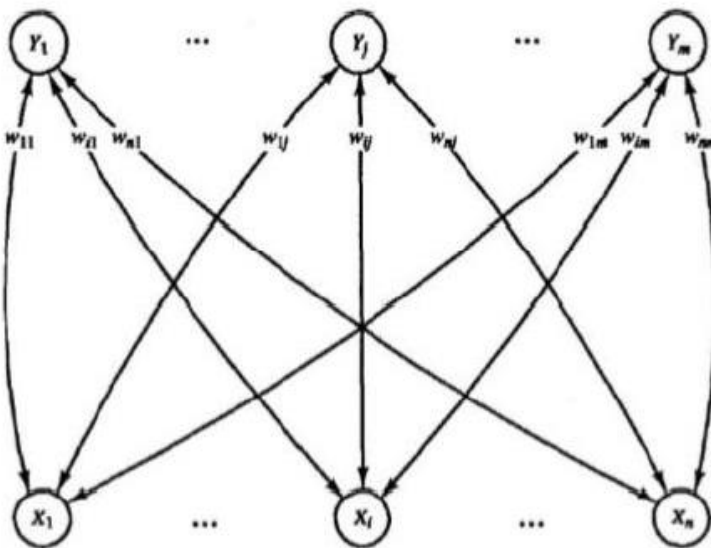


Figure : Bidirectional Associative Memory (BAM)

Architecture- The single-layer nonlinear feedback BAM network (with heteroassociative content addressable memory) has n units in the X-layer and m units in the Y-layer. The connections between the layers are bidirectional; i.e., if the weight matrix for signals sent from the X-layer to the Y-layer is W , the weight matrix for signals sent from the Y-layer to the X-layer is W^T .

Discrete BAM

The two bivalent (binary or bipolar) forms of BAM are closely related. In each, the weights are found from the sum of the outer products of the bipolar form of the training vector pairs. Also, the activation function is a step function, with the possibility of a non-zero threshold. The bipolar

vectors improve the performance of the net.

* The weight matrix to store a set of input and target vectors $s(p) : t(p)$, $p = 1, \dots, P$, where

$$s(p) = (s_1(p), \dots, s_i(p), \dots, s_n(p)) ;$$

$$t(p) = (t_1(p), \dots, t_j(p), \dots, t_m(p))$$

can be determined by the Hebb rule.

* The formulas for the entries depend on whether the training vectors are binary or bipolar.

For binary input vectors, the weight matrix

$W = \{w_{ij}\}$ is given by

$$w_{ij} = \sum_{p=1}^p (2s_i(p) - 1) (2t_j(p) - 1)$$

• For bipolar input vectors, the weight matrix $W = \{w_{ij}\}$ is given by

$$w_{ij} = \sum_{p=1}^p s_i(p) t_j(p)$$

Activation Function: The activation function for the discrete BAM is the appropriate step function, depending on whether binary or bipolar vectors are used.

♦ For binary input vectors, the activation function for the Y-layer is

$$y_j = \begin{cases} 1 & \text{if } y_in_j > 0 \\ y_j & \text{if } y_in_j = 0 \\ 0 & \text{if } y_in_j < 0 \end{cases}$$

and the activation function for the X-layer is

$$x_j = \begin{cases} 1 & \text{if } x_in_j > 0 \\ x_j & \text{if } x_in_j = 0 \\ 0 & \text{if } x_in_j < 0 \end{cases}$$

- ◆ For binary input vectors, the activation function for the Y-layer is

$$y_j = \begin{cases} 1 & \text{if } y_{in_j} > 0_j \\ y_j & \text{if } y_{in_j} = 0_j \\ -1 & \text{if } y_{in_j} < 0_j \end{cases}$$

and the activation function for the X-layer is

$$x_i = \begin{cases} 1 & \text{if } x_{in_i} > 0_i \\ x_i & \text{if } x_{in_i} = 0_i \\ -1 & \text{if } x_{in_i} < 0_i \end{cases}$$

Note that if the net input is exactly equal to the threshold value, the activation function “decides” to leave the activation of that unit at its previous value.

- The activations of all units are initialized to zero.
- The first signal to be sent from the X-layer to the Y-layer. However, if the input signal for the X-layer is the zero vector, the input signal to the Y-layer will be unchanged by the activation function, and the process will be the same as if the first piece of information had been sent from the Y-layer to the X-layer.
- * Signals are sent only from one layer to the other at any step of the process, not simultaneously in both directions.

Algorithm

1. Initialize the weights to a set of P vectors; initialize all activations to 0
2. For each testing input, do Steps 3-7.
- 3a. Present input pattern x to the X-layer, (i.e., set activations of X-layer to current input pattern). 3b. Present input pattern y to the Y-layer, (Either of the input patterns may be the zero vector.)
4. While activations are not converged, do Steps 5-7.
5. Update activations of units in Y-layer:

Compute net inputs :

$$y_{in_j} = \sum_i w_{ij} x_i$$

Compute activations :

$$y_j = f(y_{in_j})$$

Send signal to X-layer.

6. Update activations of units in X-layer:

Compute net inputs :

$$x_{in_i} = \sum_j w_{ij} y_j$$

Compute net inputs :

$$x_i = f(x_{in_i})$$

Send signal to Y-layer.

7. Test for convergence:

If the activation vectors x and y have reached equilibrium, then stop; otherwise, continue.

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Continuous BAM

A continuous bidirectional associative memory [Kosko, 1988] transforms inputs smoothly and continuously into output in the range [0, 1] using the logistic sigmoid function as the activation function for all units.

- For binary input vectors $(s(p), t(p))$, $p=1, 2, \dots, P$, the weights are determined by the formula

$$w_{ij} = \sum_{p=1}^p (2s_i(p) - 1) (2t_j(p) - 1)$$

- The activation function is the logistic sigmoid

$$f(y, in_j) = \frac{1}{1 + \exp(-y_{in_j})}$$

- where a bias is included in calculating the net input to any unit and corresponding formulas apply for the units in the X-layer.

$$y_{in_j} = b_j + \sum_i w_{ij} x_i$$

A number of other forms of BAM have been developed. In some, the activations change based on a differential equation known as Cohen-Grossberg activation dynamics (Cohen & Grossberg, 1983).

Application

Example-11 : A BAM net to associate letters with simple bipolar codes

Consider the possibility of using a (discrete) BAM network (with bipolar vectors) to map two simple letters (given by 5x3 patterns) to the following bipolar codes:



the target output vector t for letter A is $[-1 \ 1]$ and for the letter C is $[1 \ 1]$,

The weight matrices are :

$$\begin{array}{ccc}
 \text{(to store } A \rightarrow -1 \ 1) & \text{(} C \rightarrow 1 \ 1) & \text{(} w, \text{ to store both)} \\
 \begin{pmatrix} 1 & -1 \\ -1 & 1 \\ 1 & -1 \\ -1 & 1 \\ 1 & -1 \\ -1 & 1 \\ -1 & 1 \\ -1 & 1 \\ -1 & 1 \\ 1 & -1 \\ -1 & 1 \\ -1 & 1 \\ 1 & -1 \\ -1 & 1 \end{pmatrix} & \begin{pmatrix} -1 & -1 \\ 1 & 1 \\ 1 & 1 \\ 1 & 1 \\ -1 & -1 \\ -1 & -1 \\ 1 & 1 \\ -1 & -1 \\ -1 & -1 \\ -1 & -1 \\ 1 & 1 \\ -1 & -1 \\ -1 & -1 \\ 1 & 1 \end{pmatrix} & \begin{pmatrix} 0 & -2 \\ 0 & 2 \\ 2 & 0 \\ 0 & 2 \\ 0 & -2 \\ -2 & 0 \\ 0 & 2 \\ -2 & 0 \\ -2 & 0 \\ 0 & 2 \\ 0 & -2 \\ -2 & 0 \\ -2 & 0 \\ 2 & 0 \\ 0 & 2 \end{pmatrix}
 \end{array}$$

To illustrate the use of a BAM, we first demonstrate that the net gives the correct vector when presented with the vector for either the pattern A or the pattern C:

INPUT PATTERN A

$$[-12 \ -1 \quad 1 \ -11 \quad 111 \quad 1 \ -11 \quad 1 \ -1 \ 1] \ w = [-14 \ 16] \rightarrow [-11]$$

INPUT PATTERN C

$$[-111 \quad 1 \ -1 \ -1 \quad 1 \ -1 \ -1 \quad 1 \ -1 \ -1 \quad -1 \ 1 \ 1] \ w = [-14 \ 16] \rightarrow [11]$$

To see the bidirectional nature of the net, observe that the Y vectors can also be used as input. For signals sent from the Y-layer to the X-layer, the weight matrix is the transpose of the matrix W, i.e. w^T .

For the input vector associated with pattern A. namely. $(-1, 1)$, we have

$$\begin{aligned}
 [-1 \ 1] \ w^T &= [-2 \ 2 \ -2 \ 2 \ -2 \ 2 \ 2 \ 2 \ 2 \ 2 \ -2 \ 2 \ 2 \ 2 \ -2 \ 2] \\
 &\rightarrow [-11 \ -1 \quad 1 \ -11 \quad 111 \quad 1 \ -11 \quad 1 \ -11]
 \end{aligned}$$

This is pattern A.

Similarly, if we input the vector associated with pattern C, namely, $(1,$

$1)$. We obtain $[1 \ 1] \ w^T = [-111 \quad 1 \ -1 \ -1 \quad 1 \ -1 \ -1 \quad -1 \ 1 \ 1]$

Exercise:

- Q1. Write a short on Bidirectional Associative Memory.
- Q2. What is mean by Discrete BAM.
- Q3. Explain Algorithm of BAM.
- Q4. Explain Continuous BAM.
- Q5. What are the Application of BAM.

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Fuzzy Set, Properties, Operations on Fuzzy sets, Fuzzy relations **Introduction**

Fuzzy Logic

Fuzzy set theory was developed by Lotfi A. Zadeh [Zadeh, 1965], professor for computer science at the University of California in Berkeley, to provide a mathematical tool for dealing with the concepts used in natural language (linguistic variables). Fuzzy Logic is basically a multivalued logic that allows intermediate values to be defined between conventional evaluations.

However, the story of fuzzy logic started much more earlier. To devise a concise theory of logic, and later mathematics, Aristotle posited the so-called "Laws of Thought". One of these, the "Law of the Excluded Middle," states that every proposition must either be True (T) or False (F). Even when Parmenides proposed the first version of this law (around 400 Before Christ) there were strong and immediate objections: for example, Heraclitus proposed that things could be simultaneously True and not True. It was Plato who laid the foundation for what would become fuzzy logic, indicating that there was a third region (beyond T and F) where these opposites "tumbled about." A systematic alternative to the bi-valued logic of Aristotle was first proposed by Łukasiewicz around 1920, when he described a three-valued logic, along with the mathematics to accompany it. The third value he proposed can best be translated as the term "possible," and he assigned it a numeric value between T and F. Eventually, he proposed an entire notation and axiomatic system from which he hoped to derive modern mathematics.

Later, he explored four-valued logics, five-valued logics, and then declared that in principle there was nothing to prevent the derivation of an infinite-valued logic. Łukasiewicz felt that three- and infinite-valued logics were the most intriguing, but he ultimately settled on a fourvalued logic because it seemed to be the most easily adaptable to Aristotelian logic. It should be noted that Knuth also proposed a threevalued logic similar to Łukasiewicz's, from which he speculated that mathematics would become even more elegant than in traditional bi-valued logic. The notion of an infinite-valued logic was introduced in Zadeh's seminal work "Fuzzy Sets" where he described the mathematics of fuzzy set theory, and by extension fuzzy logic. This theory proposed making the membership function (or the values F and T) operate over the range of real numbers $[0, 1]$.

New operations for the calculus of logic were proposed, and showed to be in principle at least a generalization of classic logic. Fuzzy logic provides an inference morphology that enables approximate human reasoning capabilities to be applied to knowledge-based systems. The theory of fuzzy logic provides a mathematical strength to capture the uncertainties associated with human cognitive processes, such as thinking and reasoning. The conventional approaches

to knowledge representation lack the means for representating the meaning of fuzzy concepts. As a consequence, the approaches based on first order logic and classical probability theory do not provide an appropriate conceptual framework for dealing with the representation of commonsense knowledge, since such knowledge is by its nature both lexically imprecise and noncategorical. The development of fuzzy logic was motivated in large measure by the need for a conceptual framework which can address the issue of uncertainty and lexical imprecision. Some of the essential characteristics of fuzzy logic relate to the following (Zadeh, 1992): In fuzzy logic, exact reasoning is viewed as a limiting case of approximate reasoning. In fuzzy logic, everything is a matter of degree. In fuzzy logic, knowledge is interpreted a collection of elastic or, equivalently, fuzzy constraint on a collection of variables. Inference is viewed as a process of propagation of elastic constraints. Any logical system can be fuzzified. There are two main characteristics of fuzzy systems that give them better performance for specific applications. Fuzzy systems are suitable for uncertain or approximate reasoning, especially for the system with a mathematical model that is difficult to derive. Fuzzy logic allows decision making with estimated values under incomplete or uncertain information.

Theory has been attacked several times during its existence. For example, in 1972 Zadeh's colleague R. E. Kalman (the inventor of Kalman filter) commented on the importance of fuzzy logic: "...Zadeh's proposal could be severely, ferociously, even brutally criticized from a technical point of view. This would be out of place here. But a blunt question remains: Is Zadeh presenting important ideas or is he indulging in wishful thinking?..."

The heaviest critique has been presented by probability theoreticians and that is the reason why many fuzzy logic authors (Kosko, Zadeh and Klir) have included the comparison between probability and fuzzy logic in their publications. Fuzzy researchers try to separate fuzzy logic from probability theory, whereas some probability theoreticians consider fuzzy logic a probability in disguise.

Fuzzy Set

Since set theory forms a base for logic, we begin with fuzzy set theory in order to pave the way to fuzzy logic and fuzzy logic systems.

6.1.1 Set - Theoretical Operations and Basic Definitions

In classical set theory the membership of element x of a set A (A is a crisp subset of universe X) is defined by

$$A(x) = \begin{cases} 0, & \text{if } x \notin A, \\ 1, & \text{if } x \in A \end{cases}$$

6.1

The element either belongs to the set or not. In fuzzy set theory, the element can belong to the set partially with a degree and the set does not have crisp boundaries. That leads to the following definition.

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Definition 6.1.1 (fuzzy set, membership function) Let X be a nonempty set, for example $X = \mathbb{R}^n$, and be called the universe of discourse. A fuzzy $A \subset X$ is characterized by the membership function

$$\mu_A : X \rightarrow [0, 1]$$

6.2

Where $\mu_A(x)$ is a grade (degree) of membership of x in set A .

From the definition we can see that the fuzzy set theory is a generalized set theory that includes the classical set theory as a special case. Since $\{0,1\} \times [0,1]$, crisp sets are fuzzy sets. Membership function (2.2) can also be viewed as a distribution of truth of a variable. In literature fuzzy set A is often presented as a set of ordered pairs :

$$A = \{(x, \mu_A(x)) \mid x \in X\}$$

6.3

where the first part determines the element and the second part determines the grade of membership. Another way to describe fuzzy set has been presented by Zadeh, Dubois and Prade. If X is infinite, the fuzzy set A can be expressed by

$$A = \int \mu_A(x) / x$$

6.4

If X is finite, A can be expressed in the form

$$A = \mu_A(x_1)/x_1 + \dots + \mu_A(x_n)/x_n = \sum_{i=1}^n \mu_A(x_i)/x_i$$

6.5

Note : Symbol \int in (2.4) has nothing to do with integral (it denotes an uncountable enumeration) and $/$ denotes a tuple. The plus sign represents the union. Also note that fuzzy sets are membership functions. Nevertheless, we may still use the set theoretic notations like $A \cup B$.

This is the name of a fuzzy set given by $\mu_{A \cup B}$.

Example 6.1.1

Discrete case: $\mu_A = 0.1/x_1 + 0.4/x_2 + 0.8/x_3 + 1.0/x_4 + 0.8/x_5 + 0.4/x_6 + 0.1/x_7$

$$\text{Continuous case : } \mu_A(x) = \begin{cases} 1 - \frac{x-c}{h}, & x \in [c-h, h] \\ \frac{x-c}{h}, & x \in [c, c+h] \end{cases}$$

0 , otherwise

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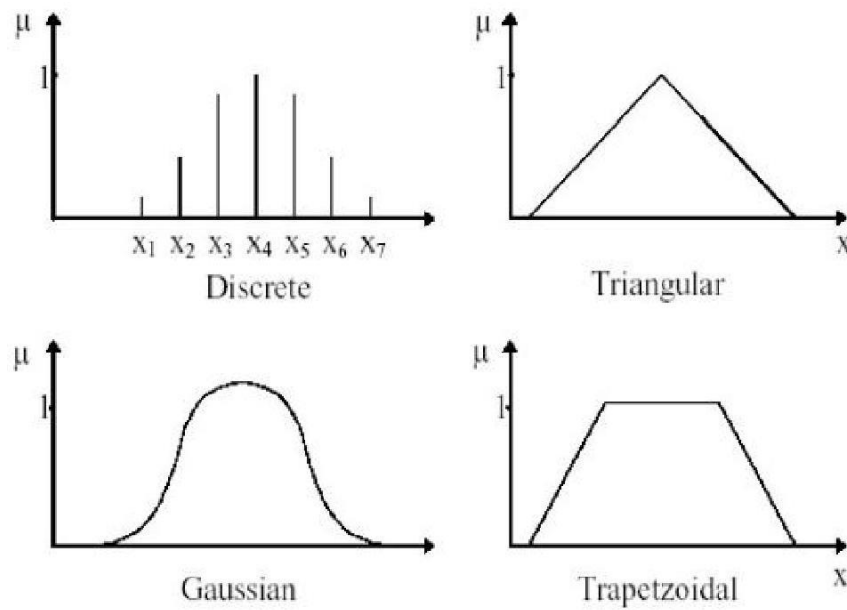


Figure 6.1 A discrete and continuous membership functions

Definition 6.1.2 (support) The support of a fuzzy set A is the crisp set that contains all elements of A with non-zero membership grade :

$$\text{supp}(A) = \{x \in X \mid \mu_A(x) > 0\}$$

6.6

If the support is finite, it is called compact support. If the support of fuzzy set A consists of only one point, it is called a fuzzy singleton. If the membership grade of this fuzzy singleton is one, A is called a crisp singleton 'Zimmermann, 1985'.

Definition 6.1.3 (core) The core (nucleus, center) of a fuzzy set A is defined by

$$\text{core}(A) = \{x \in X \mid \mu_A(x) = 1\}$$

Definition 6.1.4 (height) The height of a fuzzy set A on X is defined by

$$\text{hgt}(A) = \sup_{x \in X} \mu_A(x)$$

and A is called normal if $\text{hgt}(A) = 1$, and subnormal if $\text{hgt}(A) < 1$.

Note : Non-empty fuzzy set can be normalized by dividing $\mu_A(x)$ by $\sup_{x \in X} \mu_A(x)$. Normalizing of A can be regarded as a mapping from fuzzy sets to possibility distributions :

Normal (A) is defined by $\mu_A / \sup_{x \in X} \mu_A$ [Joslyn, 1994].

The relation between fuzzy set membership function , possibility distribution and probability distribution

p :The definition $p_A(x) \div \mu_A(x)$ could hold, if μ is additively normal. Additively normal means here that the stochastic normalization

$$\int \mu_A(x) dx = 1$$

would have to be satisfied. So it can be concluded that any given fuzzy set could define either a probability distribution or a possibility distribution, depending on the properties of μ . Both probability distributions and possibility distributions are special cases of fuzzy sets. All general distributions are in fact fuzzy sets [Joslyn, 1994].

Definition 6.1.5 (convex fuzzy set) A fuzzy set A is convex if

$$\forall x, y \in X \quad \text{and} \quad \forall Z \in [0,1] \\ \mu_A(Zx + (1-Z)y) \geq \min(\mu_A(x), \mu_A(y))$$

Definition 6.1.6 (width of a convex fuzzy set) The width of a convex fuzzy set A is defined by

$$\text{width}(A) = \sup(\text{supp}(A)) - \inf(\text{supp}(A))$$

Definition 6.1.7 (α -cut) The α -cut of a fuzzy set A is defined by

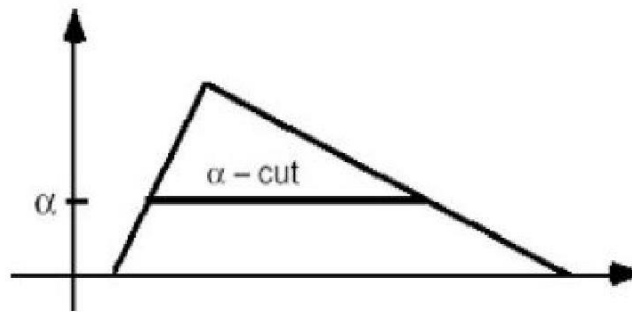


Figure 2.2 An α -cut of a triangular fuzzy number

Definition 6.1.8 (fuzzy partition) A set of fuzzy sets is called fuzzy partition if

$$\forall x, y \in X$$

$$\sum_{i=1}^{N_A} \mu_{A_i} = 1$$

provided A_j are nonempty and subsets of x .

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Definition 6.1.9 (fuzzy number) A fuzzy set (subset of a real line \mathbf{R}) is a fuzzy number, if the fuzzy set is convex, normal, membership function is piecewise continuous and the core consists of G. one value only. The family of fuzzy numbers is \mathcal{F} . In many situations people are only able to characterize numeric information imprecisely. For example, people use terms such as, about 5000, near zero, or essentially bigger than 5000. These are examples of what are called *fuzzy numbers*. Using the theory of fuzzy subsets we can represent these fuzzy numbers as fuzzy subsets of the set of real numbers.

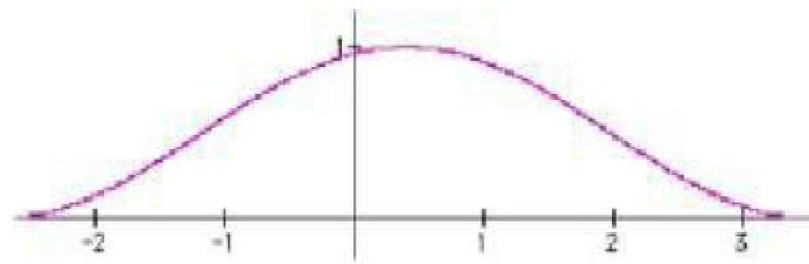


Figure 6.3 Fuzzy number

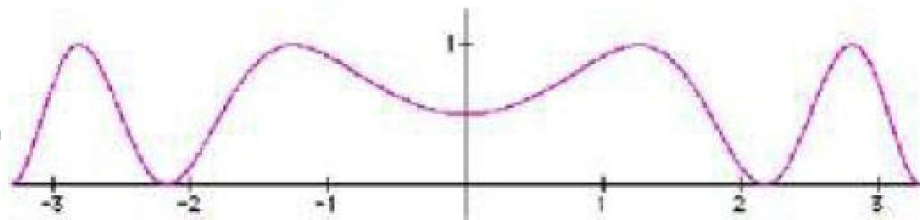


Figure 6.4 Non-fuzzy number

Note : Fuzzy number is always a fuzzy set, but a fuzzy set is not always a fuzzy number.

An example of fuzzy number is 'about 1' that is defined by $\mu(x) = \exp(-j(x-1)^2)$

It is also a quasi fuzzy number because $\lim_{x \rightarrow \pm\infty} \mu_A(x) = 0$.

$$x \rightarrow \pm\infty$$

Definition 6.1.10 (fuzzy interval) A fuzzy interval is a fuzzy set with the same restrictions as in definition 6.1.9, but the core does not need to be a one point only.

Fuzzy intervals are direct generalizations of crisp intervals $[a, b] \subset \mathbf{R}$.

Definition 6.1.11 (LR-representation of fuzzy numbers) Any fuzzy number can be described by

$$\mu_A(x) = \begin{cases} L((a-x)/\alpha) & , x \in [a-\alpha, a] \\ 1 & , x \in [a, b] \\ R((x-b)/\beta) & , x \in [b, b+\beta] \\ 0 & , \text{otherwise} \end{cases}$$

Where $[a, b]$ is the core of A , and $L:[0,1] \rightarrow [0,1]$, $R:[0,1] \rightarrow [0,1]$ are shape functions (called briefly-shape functions) that are continuous and nonincreasing such that $L(0)=R(0)=1, L(1)=R(1)=0$, where L stands for left-hand side and R stands for right-hand side of membership function [Zimmermann, 1993].

Definition 6.1.12 (*LR-representation of quasi fuzzy numbers*) Any quasi fuzzy number can be described by

$$\mu_A(x) = \begin{cases} L((a-x)/\alpha) & , x \in [a-\alpha, a] \\ 1 & , x \in [a, b] \\ R((x-b)/\beta) & , x \in [b, b+\beta] \end{cases}$$

Where $[a, b]$ is the core of A , and $L:[0,\infty) \rightarrow [0,1], R:[0,\infty) \rightarrow [0,1]$

are shape functions that are continuous and non-increasing such that $L(0)=R(0)=1$

and they approach zero: $\lim_{x \rightarrow \infty} L(x) = 0, \lim_{x \rightarrow \infty} R(x) = 0$.

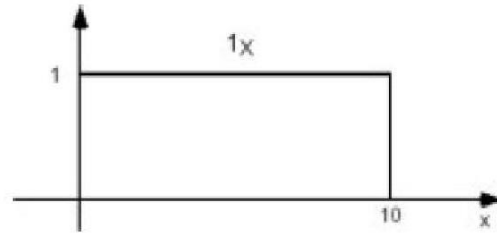
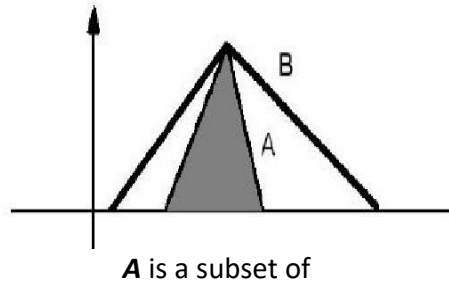
For example, $f(x) = e^{-x}$, $g(x) = e^{-x^2}$ and $f(x) = \max(0, 1-x)$ are such shape functions. In the following the classical set theoretic operations are extended to fuzzy sets.

Definition 6.1.13 (*set theoretic operations*)

$\mu_{\emptyset} \doteq 0$		(empty set)
$\mu_X \doteq 1$		(basic set, universe)
$A = B \Leftrightarrow \mu_A(x) = \mu_B(x) \quad \forall x \in X$		(identity)
$A \subseteq B \Leftrightarrow \mu_A(x) \leq \mu_B(x) \quad \forall x \in X$		(subsethood)
$\forall x \in X: \mu_{A \cup B}(x) = \max(\mu_A(x), \mu_B(x))$		(union)
$\forall x \in X: \mu_{A \cap B}(x) = \min(\mu_A(x), \mu_B(x))$		(intersection)
$\forall x \in X: \mu_{A^c}(x) = 1 - \mu_A(x)$		(complement)

Union could also be represented by $A \cup B = \{(x, \max(\mu_A(x), \mu_B(x))) \mid x \in X\}$

Same notation could also be used with intersection and complement.



The graph of the universal fuzzy subset in $X = [0, 10]$.

Figure 6.5 Fuzzy sets

Theorem 2.1.1 The following properties of set theory are valid

$A = A$	(involution)
$A \cap B = B \cap A$	(commutativity)
$A \cup B = B \cup A$	(commutativity)
$(A \cap B) \cap C = A \cap (B \cap C)$	(associativity)
$(A \cup B) \cup C = A \cup (B \cup C)$	(associativity)
$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$	(distributivity)
$A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$	(distributivity)
$A \cap A = A$	(idempotence)
$A \cup A = A$	(idempotence)
$A \cap (A \cup B) = A$	(absorption)
$A \cup (A \cap B) = A$	(absorption)
$(A \cap B)^c = A^c \cup B^c$	(De Morgan's laws)
$(A \cup B)^c = A^c \cap B^c$	(De Morgan's laws)

Proof : Above properties can be proved by simple direct calculations. For example,

$$A = 1 - (1 - A) = A.$$

Fuzzy Relations

Definition 2.1.17 (fuzzy relation) Fuzzy relation is characterized by a function

$$\mu_R : X_1 \times \dots \times X_m \rightarrow [0, 1] \text{ where } X_i \text{ are the universes of discourse and } X_1 \times \dots \times X_m$$

is the product space. If we have two finite universes, the fuzzy relation can be presented as a matrix (fuzzy matrix) whose elements are the intensities of the relation and **R** has the membership function

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$\mu_R(u, w)$, where $u \in X_1$, $w \in X_2$. Two fuzzy relations are combined by a so called sup-* or max-min composition, which will be given in the definition 2.1.19.

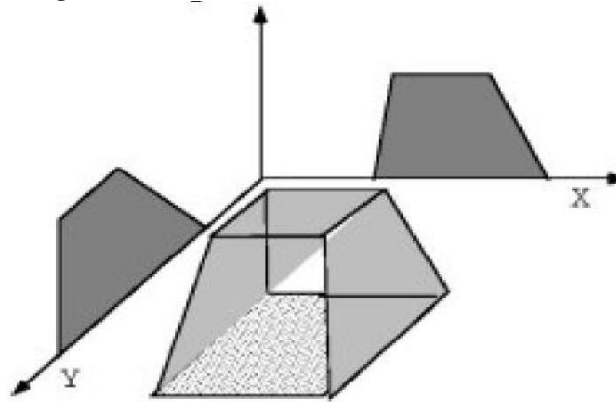


Figure 2.7 Shadows of a fuzzy relation

Note: Fuzzy relations are fuzzy sets, and so the operations of fuzzy sets (union, intersection, etc.) can be applied to them.

Example 2.1.2 Let the fuzzy relation R = “approximately equal” correspond to the equality of two numbers. The intensity of cell $R(u, w)$ of the following matrix can be interpreted as the degree of membership of the ordered pair in R . The numbers to be compared are $\{1, 2, 3, 4\}$ and $\{3, 4, 5, 6\}$.

$u \setminus$	1	2	3	4
3	.6	.8	1	.8
4	.4	.6	.8	1
5	.2	.4	.6	.8
6	.1	.2	.4	.6

The matrix shows, that the pair (4, 4) is approximately equal with intensity 1 and the pair (1, 6) is approximately equal with intensity 0.1.

Definition 2.1.18 (*Cartesian product*) Let $A_1 \times X_1$ be fuzzy sets. Then the Cartesian product is defined

$$\mu_{A_1 \times \dots \times A_n}(x_1, \dots, x_n) = \min\{\mu_{A_i}(x_i)\}$$

Note: min can be replaced by a more general t-norm.

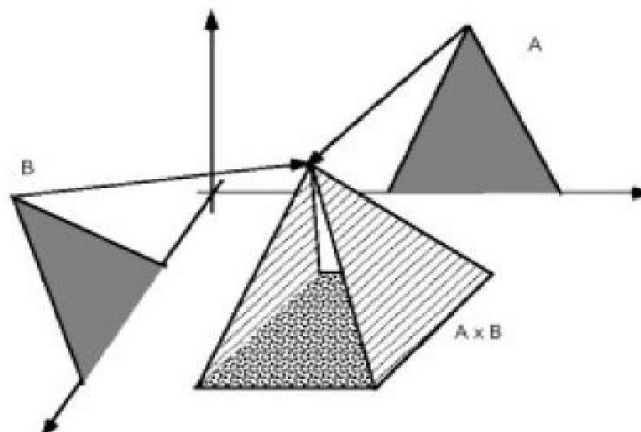


Figure 2.8 Cartesian product of two fuzzy sets is a fuzzy relation in $X \times Y$

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Definition 2.1.19 (*sup- \ast -composition, max-min composition*) The composition of two relations R o S is defined as a membership function in $X \times Y$

$$\mu_{ROS}(x) = \sup_{\varpi} T(\mu_R(x, \varpi), \mu_S(\varpi, y))$$

where R is relation in $X \times V$ and S is relation in $V \times Y$, $x \in X$, $y \in Y$ and T is a t-norm sup is

replaced by the sum. If S is just a fuzzy set (not a relation) in V , then (2.15) becomes

$$\mu_{ROS}(x) = \sup_{\varpi} T(\mu_R(x, \varpi), \mu_S(\varpi))$$

Example 2.1.3 Let $X = \{1, 2, 3, 4\}$, fuzzy set $A = \text{small} = \{(1, 1), (2, 0.6), (3, 0.2), (4, 0)\}$ and fuzzy relation $R = \text{"approximately equal"}$.

R	1	.5	0	0
	.5	1	.5	0
	0	.5	1	.5
	0	0	.5	1

$$\begin{aligned} A \circ B &= \max_x \min \{ \mu_A(x), \mu_R(x, y) \} \\ &= \{(1, 1), (2, .6), (3, .5), (4, .2)\} \end{aligned}$$

The interpretation of example: x is small. If x and y are approximately equal, y is more or less small.

Exercise:

Q1. What is mean by Fuzzy Logic.

Q2. Explain Fuzzy Set.

Q3. Explain Fuzzy Relation.

Q4. Write a definition of Fuzzy Number.

Q5. Write a definition of LR-representation of fuzzy numbers.

Chapter 12

The extension principle, fuzzy measures, Membership function's

The Extension Principle

The extension principle is said to be one of the most important tools in fuzzy logic. It gives means to generalize non-fuzzy concepts, e.g., mathematical operations, to fuzzy sets. Any fuzzifying generalization must be consistent with the crisp cases.

Definition 2.1.20 (extension principle) Let A_1, \dots, A_n be fuzzy sets, defined on $X_1 \dots X_n$ and let f be a function $f: X_1 \times \dots \times X_n \rightarrow Y$. The extension of f operating on A_1, \dots, A_n gives a membership function (fuzzy set F)

$$\mu_F(y) = \sup_{u_1 \dots u_n : f(u_1, \dots, u_n) = y} \min(\mu_{A_1}(u_1), \dots, \mu_{A_n}(u_n))$$

when the inverse of f exists. Otherwise define $\mu_F(y) = 0$. Function f is called inducing mapping.

If the domain is either discrete or compact, sup-min can be replaced by max-min. On continuous domains sup-operation and the operation that satisfies criterion

$$S_W(x, y) = \begin{cases} x, & y = 0 \\ y, & x = 0 \\ 1, & \text{otherwise} \end{cases}$$

should be used [Driankov et. al., 1993].

Fuzzy Rules

Fuzzy logic was originally meant to be a technique for modeling the human thinking and reasoning, which is done by fuzzy rules. This idea has been replaced by the thought that the fuzzy rules form an interface between humans and computers [Brown & Harris, 1994]. Humans explain their actions and knowledge using linguistic rules and fuzzy logic is used to represent this knowledge on computers. There are three principal ways to obtain these rules :

1. human experts provide rules
2. data driven: rules are formed by training methods
3. combination of 1. and 2.

The first way is the ideal case for fuzzy systems. Although the rules are not precise, they contain important information about the system. In practice human experts may not provide a sufficient number of rules and especially in the case of complex systems the amount of knowledge may be very small or even non-existent. Thus the second way must be used instead of the first one (provided the data is available). The third way is suited for the cases when some knowledge exists and sufficient amount of data for training is available. In this case fuzzy rules got from experts roughly

approximate the behavior of the system, and by applying training this approximation is made more precise. Rules provided by the expert's form an initial point for the training and thus exclude the necessity of random initialization and diminish the risk of getting stuck in a local minimum (provided the expert knowledge is good enough).

It has been shown in [Mouzouris, 1996] that linguistic information is important in the absence of sufficient numerical data but it becomes less important as more data become available.

Fuzzy rules define the connection between input and output fuzzy linguistic variables and they can be seen to act as associative memories. Resembling inputs are converted to resembling outputs. Rules have a structure of the form :

IF (antecedent) **THEN** (consequent)

In more detail, the structure is

IF (x is A_i AND AND x is A_d THEN (y is B)

where A_{ij} and B_i are fuzzy sets (they define complete fuzzy partitions) in $U \subset R^V$, respectively. Linguistic variable x is a vector of dimension d in $U \subset R^d$ and linguistic variable $V \in \{1, \dots, V\}$. Vector x is an input to the fuzzy system and y is an output of the fuzzy system. Note that B_i can also be a singleton (consequence part becomes: ...THEN (y is z)). Further, if fuzzy system is used as a classifier, the consequence part becomes: ...THEN class is c .

A fuzzy rule base consists of a collection of rules $\{R^1, R^2, \dots, R^M\}$, where each rule R can be considered to be of the form (2.34). This does not cause a loss of generality, since multi-input-multioutput (MIMO) fuzzy logic system can always be decomposed into a group of multi-input-singleoutput (MISO) fuzzy logic systems. Furthermore (2.34) includes the following types of rules as a special case [Wang, 1994] :

- IF (x_1 is A^i_1 AND ... AND x_i is A^i_k) THEN (y is B^i) where $k > d$
- IF (x_1 is A^i_1 AND ... AND x_i is A^i_k) OR (x_{k+1} is A^i_{k+1} AND ... AND x_d is A^i_d) THEN (y is B^i)
- (y is B^i)
- (y is B^i) UNLESS (x_i is A^i_1) AND ... AND x_d is A^i_d)
- non-fuzzy rules

In control systems the production rules are of the form

IF <process state> **THEN** <control action>

where the <process state> part contains a description of the process output at the k th sampling instant. Usually this description contains values of error and change-of-error. The <control action> part describes the control output (change-in-control) which should be produced given the particular <process state>. For example, a fuzzified PI-controller is of the type (2.35). Important properties for a set of rules are completeness, consistency and continuity. These are defined in the following.

Definition 6.1.27 (completeness) A rule base is said to be complete if any combination of input values results in an appropriate output value:

$$\sum_{x \in X} \text{hgt}(\text{out}(x)) > 0$$

Definition 6.1.28 (inconsistency) A rule base is inconsistent if there are two rules with the same rule antecedent but different rule consequences.

This means that two rules that have the same antecedent map to two non-overlapping fuzzy output sets. When the output sets are non-overlapping, then there is something wrong with the output variables or the rules are inconsistent or discontinuous.

Definition 6.1.29 (continuity) A rule base is continuous if the neighboring rules do not have fuzzy output sets that have

empty intersection.

Exercise:

- Q1. What is meant by The Extension Principle.
Q2. What are the Fuzzy Rule.
Q3. Write a definition of completeness.
Q4. Write a Definition of inconsistency.

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Fuzzifier and Defuzzifier

The *fuzzifier* maps a crisp point x into a fuzzy set

$$\mu_{A'}(x) = \begin{cases} 1, & \text{if } x = x' \\ 0, & \text{otherwise} \end{cases}$$

Where x' is the input. Fuzzifier of the form (2.43) is called a singleton fuzzifier. If the input contains noise it can be modeled by using fuzzy number. Such fuzzifier could be called a nonsingleton fuzzifier. Because of simplicity, singleton fuzzifier is preferred to nonsingleton fuzzifier.

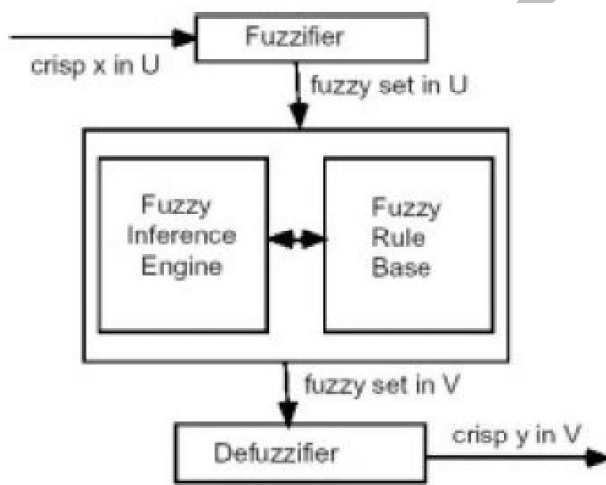


Figure 6.15 Fuzzy logic controller

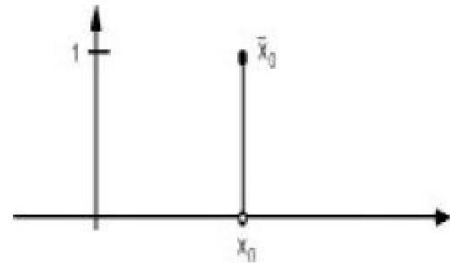


Figure 2.16 Fuzzy singleton as fuzzifier

The *defuzzifier* maps fuzzy sets to a crisp point. Several defuzzification methods have been suggested. The following five are the most common :

- **Center of Gravity (CoG)** : In the case of 1-dimensional fuzzy sets it is often called the Center of Area (CoA) method. Some authors (for example, [Driankov et. al., 1993]) regard CoG and CoA as a same method, when other (for example, [Jager, 1995]) give them different forms. If CoA is calculated by dividing the area of combination of output membership functions by two and then taking from the left so much that we get an area equal to the right one, then it is clearly a distinct method. CoG determines center of gravity of the mass, which is formed as a combination of the clipped or scaled output fuzzy membership functions. The intersection part of these membership functions can be taken once or twice into calculation. Driankov [1993] separates the Center of Gravity methods such that, if the

intersection is calculated once the method is Center of Area and if it is calculated twice the method is called Center of Sums (CoS). In Fig. 2.17 defuzzified value obtained by CoS is slightly smaller than obtained by the CoA - method.

- **Height method (HM)** : Can be considered as a special case of CoG, whose output membership functions are singletons. If symmetric output sets are used in CoG, they have the same centroid no matter how wide the set is and CoG reduces to HM. HM calculates a normalized weighted sum of the clipped or scaled singletons. HM is sometimes called Center average defuzzifier or fuzzy-mean defuzzifier.
- **Middle of Maxima (MoM)** : Calculates the center of maximum in clipped membership function. The rest of the distribution is considered unimportant.
- **First of Maxima (FoM)**: As MoM, but takes the leftmost value instead of center value.

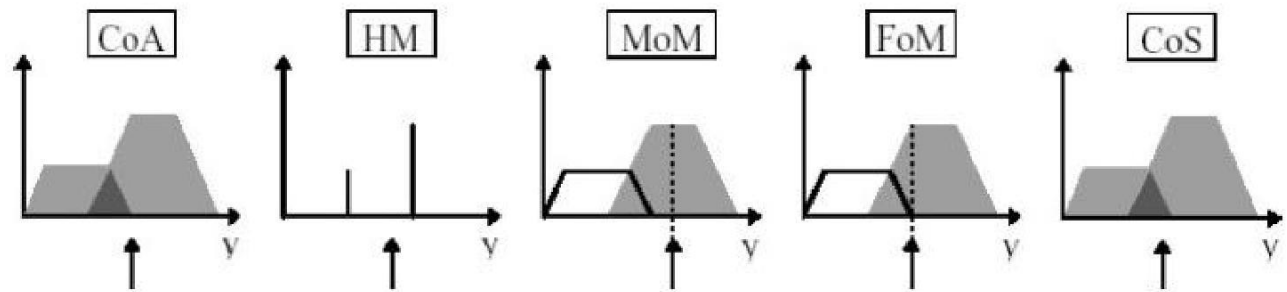


Figure 6.17 Defuzzification methods

Defuzzification methods can be compared by some criteria, which might be the continuity of the output and the computational complexity. HM, CoA and CoS produce continuous outputs. The simplest and quickest method of these is the HM method and for large problems it is the best choice. The fuzzy systems using it have a close relation to some well-known interpolation methods (we will return to this relation later).

The maximum methods (MoM, FoM) have been widely used. The underlying idea of MoM (with max-min inference) can be explained as follows. Each input variable is divided into a number of intervals, which means that the whole input space is divided into a large number of d-dimensional boxes. If a new input point is given, the corresponding value for y is determined by finding which box the point falls in and then returning the average value of the corresponding y -interval associated with that input box. Because of the piecewise constant output, MoM is inefficient for approximating nonlinear continuous functions. Kosko has shown [Kosko, 1997] that if there are many rules that fire simultaneously, the maximum function tends to approach a constant function. This may cause problems especially in control.

The property of CoS is that the shape of final membership function used as a basis for defuzzification will resemble more and more normal density function when the number of functions used in summation grows. Systems of this type that sum up the output membership functions to form final output set are called *additive fuzzy systems*.

Fuzzy logic is widely used in machine control. The term “fuzzy” refers to the fact that the logic involved can deal with concepts that cannot be expressed as the “true” or “false” but rather as “partially true”. Although alternative approaches such as genetic algorithms and neural networks can perform just as well as fuzzy logic in many cases, fuzzy logic has the advantage that the solution to the problem can be cast in terms that human operators can understand, so that their experience can be used in the design of the controller. This makes it easier to mechanize tasks that are already successfully performed by humans.

Fuzzy controllers are very simple conceptually. They consist of an input stage, a processing stage, and an output stage. The input stage maps sensor or other inputs, such as switches, thumbwheels, and

so on, to the appropriate membership functions and truth values. The processing stage invokes each appropriate rule and generates a result for each, then combines the results of the rules. Finally, the output stage converts the combined result back into a specific control output value.

The most common shape of membership functions is triangular, although trapezoidal and bell curves are also used, but the shape is generally less important than the number of curves and their placement. From three to seven

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curves are generally appropriate to cover the required range of an input value, or the “universe of discourse” in fuzzy jargon.

As discussed earlier, the processing stage is based on a collection of logic rules in the form of IF THEN statements, where the IF part is called the “antecedent” and the THEN part is called the “consequent”. Typical fuzzy control systems have dozens of rules.

Consider a rule for a thermostat :

IF (temperature is “cold”) THEN turn (heater is “high”)

This rule uses the truth value of the “temperature” input, which is some truth value of “cold”, to generate a result in the fuzzy set for the “heater” output, which is some value of “high”. This result is used with the results of other rules to finally generate the crisp composite output. Obviously, the greater the truth value of “cold”, the higher the truth value of “high”, though this does not necessarily mean that the output itself will be set to “high” since this is only one rule among many. In some cases, the membership functions can be modified by “hedges” that are equivalent to adverbs. Common hedges include “about”, “near”, “close to”, “approximately”, “very”, “slightly”, “too”, “extremely”, and “somewhat”. These operations may have precise definitions, though the definitions can vary considerably between different implementations. “Very”, for one example, squares membership functions; since the membership values are always less than 1, this narrows the membership function. “Extremely” cubes the values to give greater narrowing, while “somewhat” broadens the function by taking the square root.

In practice, the fuzzy rule sets usually have several antecedents that are combined using fuzzy operators, such as AND, OR, and NOT, though again the definitions tend to vary: AND, in one popular definition, simply uses the minimum weight of all the antecedents, while OR uses the maximum value. There is also a NOT operator that subtracts a membership function from 1 to give the “complementary” function.

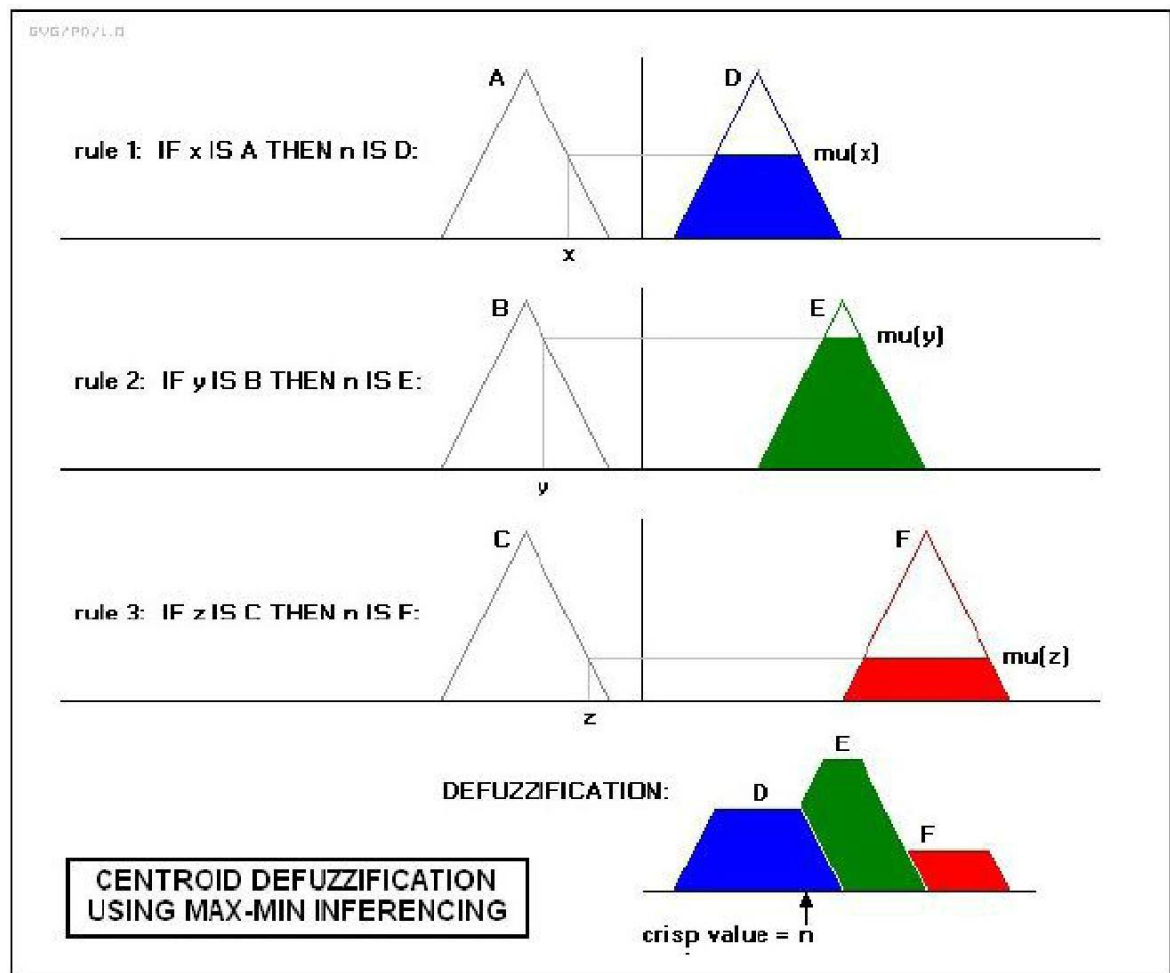
There are several ways to define the result of a rule, but one of the most common and simplest is the “max-min” inference method, in which the output membership function is given the truth value generated by the premise.

Rules can be solved in parallel in hardware, or sequentially in software. The results of all the rules that have fired are “defuzzified” to a crisp value by one of several methods. There are dozens, in theory, each with various advantages or drawbacks.

The “centroid” method is very popular, in which the “center of mass” of the result provides the crisp value. Another approach is the “height” method, which takes the value of the biggest contributor. The centroid method favors the rule with the output of greatest area, while the height method obviously favors the rule with the greatest output value.

The diagram below demonstrates max-min inferencing and centroid defuzzification for a system with input variables “x”, “y”, and “z” and an output variable “n”. Note that “mu” is standard fuzzylogic nomenclature for “truth value”:

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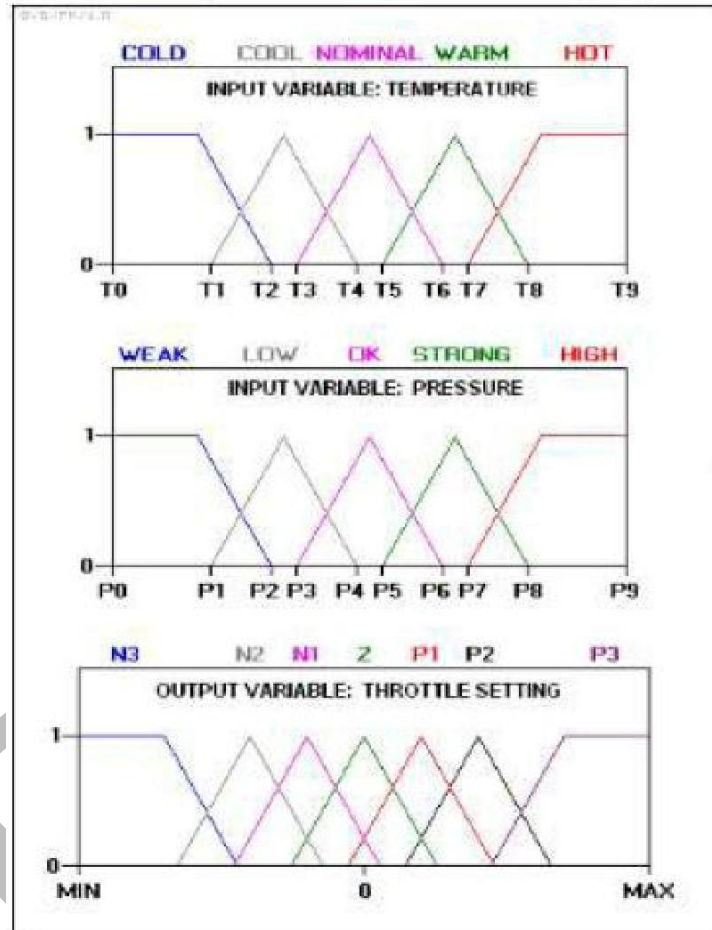


Notice how each rule provides a result as a truth value of a particular membership function for the output variable. In centroid defuzzification the values are OR'd, that is, the maximum value is used and values are not added, and the results are then combined using a centroid calculation. Fuzzy control system design is based on empirical methods, basically a methodical approach to trial-and-error. The general process is as follows :

- Document the system's operational specifications and inputs and outputs.
- Document the fuzzy sets for the inputs.
- Document the rule set.
- Determine the defuzzification method.
- Run through test suite to validate system, adjust details as required.
- Complete document and release to production.

As a general example, consider the design of a fuzzy controller for a steam turbine. The block diagram of this control system appears as follows :

The input and output variables map into the following fuzzy set :



where:

- N3 : Large negative.
- N2 : Medium negative.
- N1 : Small negative.
- Z : Zero.
- P1 : Small positive.
- P2 : Medium positive.
- P3 : Large positive.

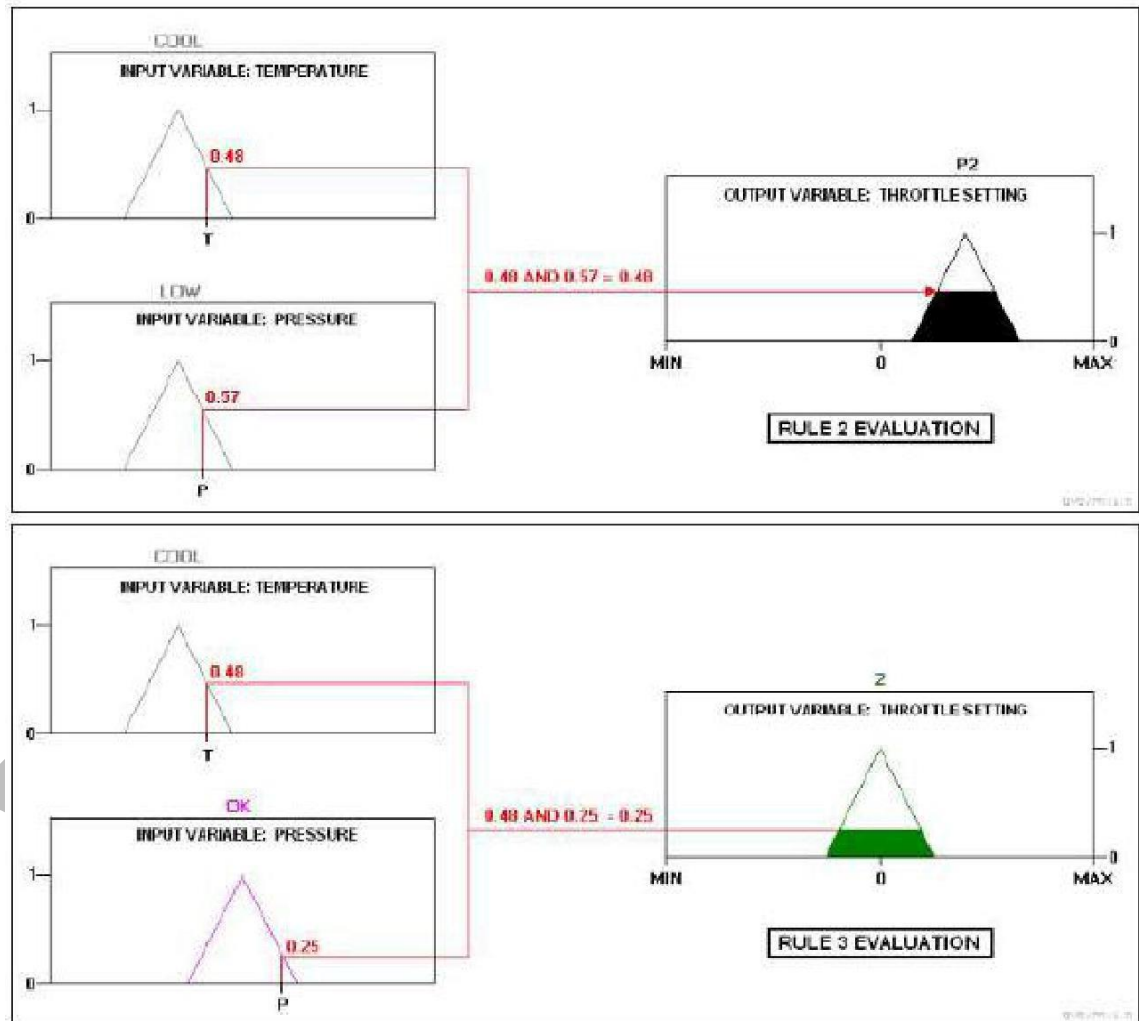
The rule set includes such rules as:

- rule 1 : IF temperature IS cool AND pressure IS weak, THEN throttle is P3.
- rule 2 : IF temperature IS cool AND pressure IS low, THEN throttle is P2.
- rule 3 : IF temperature IS cool AND pressure IS ok, THEN throttle is Z.

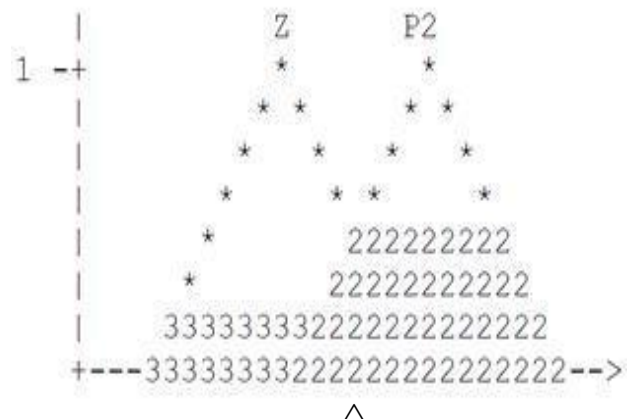
rule 4 : IF temperature IS cool AND pressure IS strong, THEN throttle is N2.

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In practice, the controller accepts the inputs and maps them into their membership functions and truth values. These mappings are then fed into the rules. If the rule specifies an AND relationship between the mappings of the two input variables, as the examples above do, the minimum of the two is used as the combined truth value; if an OR is specified, the maximum is used. The appropriate output state is selected and assigned a membership value at the truth level of the premise. The truth values are then defuzzified. For an example, assume the temperature is in the “cool” state, and the pressure is in the “low” and “ok” states. The pressure values ensure that only rules 2 and 3 fire:



The two outputs are then defuzzified through centroid defuzzication :



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The output value will adjust the throttle and then the control cycle will begin again to generate the next value .

Building a fuzzy controller[edit]

Consider implementing with a microcontroller chip a simple feedback controller :



A fuzzy set is defined for the input error variable “e”, and the derived change in error, “delta”, as well as the “output”, as follows:

- LP : large positive
- SP : small positive
- ZE : zero
- SN : small negative
- LN : large negative

If the error ranges from -1 to +1, with the analog-to-digital converter used having a resolution of 0.25, then the input variable’s fuzzy set (which, in this case, also applies to the output variable) can be described very simply as a table, with the error / delta / output values in the top row and the truth values for each membership function arranged in rows beneath :

	-1	-0.75	-0.5	-0.25	0	0.25	0.5	0.75	1
mu (LP)	0	0	0	0	0	0	0.3	0.7	1
mu (SP)	0	0	0	0	0.3	0.7	1	0.7	0.3
mu (ZE)	0	0	0.3	0.7	1	0.7	0.3	0	0
mu (SN)	0.3	0.7	1	0.7	0.3	0	0	0	0
mu (LN)	1	0.7	0.3	0	0	0	0	0	0

or, in

graphical form (where each “X” has a value of 0.1):

	LN	SN	ZE	SP	LP	
-1.0	XXXXXXXXXX	XXX	:	:	:	+
-0.75	XXXXXXX	XXXXXXX	:	:	:	
-0.5	XXX	XXXXXXXXXX	XXX	:	:	
-0.25	:	XXXXXXX	XXXXXXX	:	:	
0.0	:	XXX	XXXXXXXXXX	XXX	:	
0.25	:	:	XXXXXXX	XXXXXXX	:	
0.5	:	:	XXX	XXXXXXXXXX	XXX	
0.75	:	:	:	XXXXXXX	XXXXXXX	
1.0	:	:	:	XXX	XXXXXXXXXX	

13

Unedited Version: Neural Network and Fuzzy System

Suppose this fuzzy system has the following rule base :

rule 1 : IF e = ZE AND delta = ZE THEN output = ZE
 rule 2 : IF e = ZE AND delta = SP THEN output = SN
 rule 3 : IF e = SN AND delta = SN THEN output = LP
 rule 4 : IF e = LP OR delta = LP THEN output = LN

These rules are typical for control applications in that the antecedents consist of the logical combination of the error and error-delta signals, while the consequent is a control command output. The rule outputs can be defuzzified using a discrete centroid computation :

$$\text{SUM} (I = 1 \text{ TO } 4 \text{ OF } (\mu(I) * \text{output}(I))) / \text{SUM}(I = 1 \text{ TO } 4 \text{ OF } \mu(I))$$

Now, suppose that at a given time we have:

$$e = 0.25$$

$$\text{delta} = 0.5$$

Then this gives :

	e	delta
mu(LP)	0	0.3
mu(SP)	0.7	1
mu(ZE)	0.7	0.3
mu(SN)	0	0
mu(LN)	0	0

Plugging this into rule 1 gives :

rule 1 : IF e = ZE AND delta = ZE THEN output = ZE
 $\mu(1) = \text{MIN} (0.7, 0.3) = 0.3$
 $\text{output}(1) = 0$

-- where :

- $\mu(1)$: Truth value of the result membership function for rule 1. In terms of a centroid calculation, this is the "mass" of this result for this discrete case.
- $\text{output}(1)$: Value (for rule 1) where the result membership function (ZE) is maximum over the output variable fuzzy set range. That is, in terms of a centroid calculation, the location of the "center of mass" for this individual result. This value is independent of the value of "mu". It simply identifies the location of ZE along the output range.

The other rules give :

rule 2 : IF e = ZE AND delta = SP THEN output = SN
 $\mu(2) = \text{MIN} (0.7, 1) = 0.7$
 $\text{output}(2) = -0.5$
 rule 3 : IF e = SN AND delta = SN THEN output = LP
 $\mu(3) = \text{MIN} (0.0, 0.0) = 0$
 $\text{output}(3) = 1$

rule 4 : IF e = LP OR delta = LP THEN output = LN
mu (4) = MAX (0.0, 0.3) = 0.3
output (4) = -1

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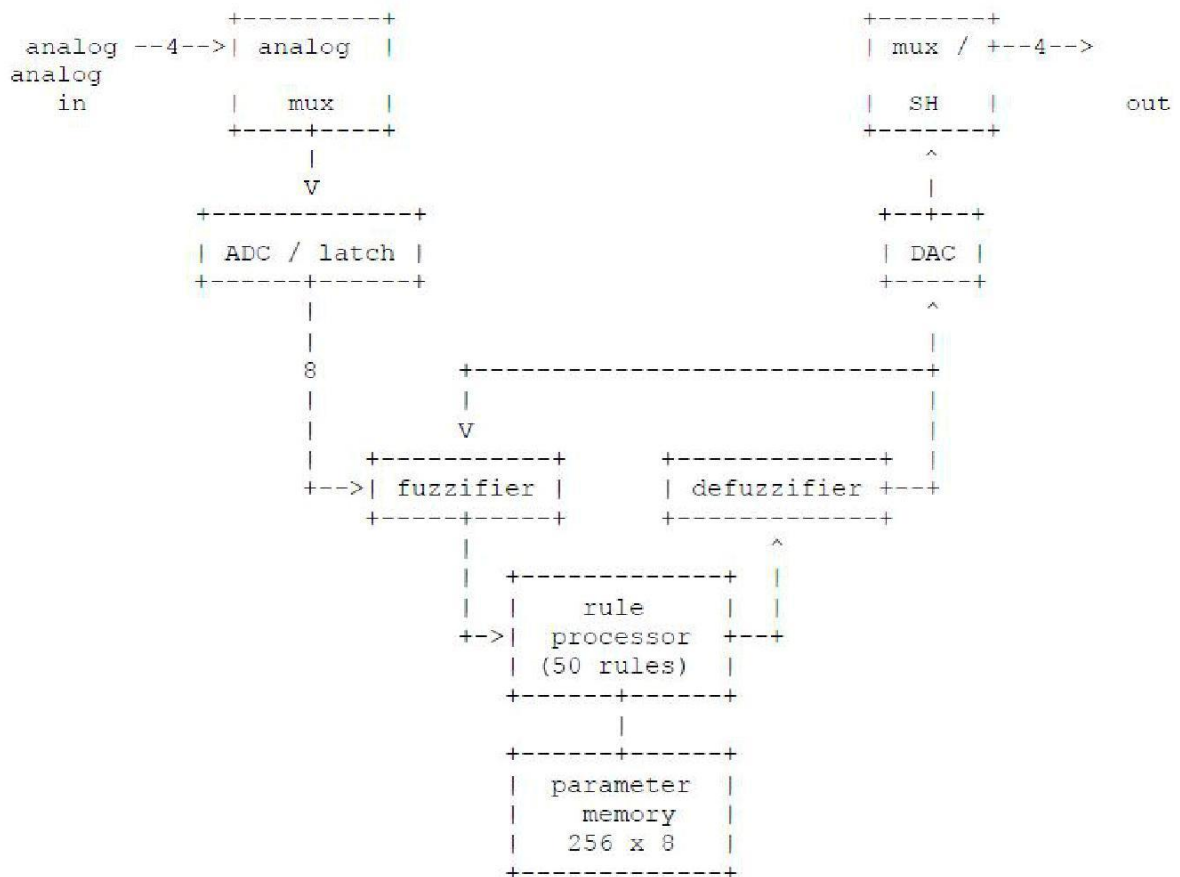
The centroid computation yields :

-for the final control output. Simple. Of course the hard part is figuring out what rules actually work correctly in practice.

If you have problems figuring out the centroid equation, remember that a centroid is defined by summing all the moments (location times mass) around the center of gravity and equating the sum to zero. So if is the center of gravity, is the location of each mass, and is each mass, this gives :

In our example, the values of μ correspond to the masses, and the values of X to location of the masses (μ , however, only 'corresponds to the masses' if the initial 'mass' of the output functions are all the same/equivalent. If they are not the same, i.e. some are narrow triangles, while others maybe wide trapizoids or shouldered triangles, then the mass or area of the output function must be known or calculated. It is this mass that is then scaled by μ and multiplied by its location X_i).

This system can be implemented on a standard microprocessor, but dedicated fuzzy chips are now available. For example, Adaptive Logic INC of San Jose, California, sells a "fuzzy chip", the AL220, that can accept four analog inputs and generate four analog outputs. A block diagram of the chip is shown below :



ADC : analog-to-digital converter

DAC : digital-to-analog converter

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:
sampl
e/hold

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Antilock brakes

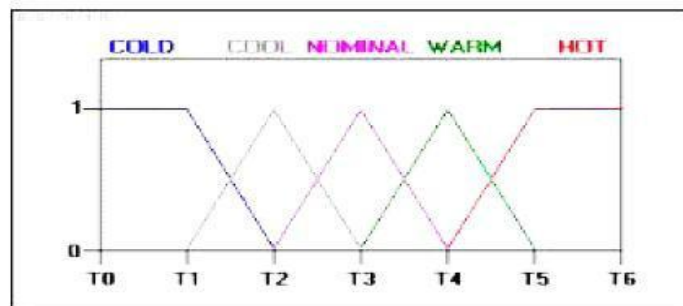
As a first example, consider an anti-lock braking system, directed by a microcontroller chip. The microcontroller has to make decisions based on brake temperature, speed, and other variables in the system.

The variable “temperature” in this system can be subdivided into a range of “states”: “cold”, “cool”, “moderate”, “warm”, “hot”, “very hot”. The transition from one state to the next is hard to define.

An arbitrary static threshold might be set to divide “warm” from “hot”. For example, at exactly 90 degrees, warm ends and hot begins. But this would result in a discontinuous change when the input value passed over that threshold. The transition wouldn’t be smooth, as would be required in braking situations.

The way around this is to make the states *fuzzy*. That is, allow them to change gradually from one state to the next. In order to do this there must be a dynamic relationship established between different factors.

We start by defining the input temperature states using “membership functions”:



With this scheme, the input variable’s state no longer jumps abruptly from one state to the next. Instead, as the temperature changes, it loses value in one membership function while gaining value in the next. In other words, its ranking in the category of cold decreases as it becomes more highly ranked in the warmer category.

At any sampled timeframe, the “truth value” of the brake temperature will almost always be in some degree part of two membership functions: i.e.: ‘0.6 nominal and 0.4 warm’, or ‘0.7 nominal and 0.3 cool’, and so on.

The above example demonstrates a simple application, using the abstraction of values from multiple values. This only represents one kind of data, however, in this case, temperature.

Adding additional sophistication to this braking system, could be done by additional factors such as traction, speed, inertia, set up in dynamic functions, according to the designed fuzzy system.[9]

Logical interpretation of fuzzy control

In spite of the appearance there are several difficulties to give a rigorous logical interpretation of the IF-THEN rules. As an example, interpret a rule as IF (*temperature is "cold"*) THEN (*heater is "high"*) by the first order formula $Cold(x) \rightarrow High(y)$ and assume that r is an input such that $Cold(r)$ is false. Then the formula $Cold(r) \rightarrow High(t)$ is true for any t and therefore any t gives a correct control given r . A rigorous logical justification of fuzzy control is given in Hajek's book (see Chapter 7) where fuzzy control is represented as a theory of Hajek's basic logic.[2] Also in Gerla 2005 [10] another logical approach to fuzzy control is proposed based on fuzzy logic programming. Indeed, denote by f the fuzzy function arising of an IF-THEN systems of rules. Then we can translate this system into a fuzzy program P containing a series of rules whose head is " $Good(x,y)$ ". The interpretation of this predicate in the least fuzzy Herbrand model of P coincides with f . This gives further useful tools to fuzzy control.

Exercise:

- Q1. Explain Fuzzifier and Defuzzifier.
- Q2. Write a short note on Fuzzification.
- Q3. Explain the concept of defuzzification methods.
- Q4. What is the process of Building a fuzzy controller.
- Q5. Write a short note on centroid computation yields.
- Q6. Write a short note on Antilock brakes.